

LINES IN THE PLANE AND DECOMPOSITIONS OF GRAPHS

1.

Incidence axioms (actually just one of them)

- Let A, B be distinct points, then there is exactly one line containing both of them.

Metric axioms

(Metric) notation of distance $d: \mathbb{X} \times \mathbb{X} \rightarrow \mathbb{R}$

$$d(x, y) \geq 0 \quad (= 0 \text{ iff } x = y)$$

$$d(x, y) = d(y, x)$$

$$d(x, z) \leq d(x, y) + d(y, z)$$

- triangle inequality

Order axioms (Tarski's axioms)

Notation : "betweenness"

$Bxyz \equiv y$ is "between" x and z , y is a point on the line segment xz

"congruence"

$wx \equiv yz$ means $d(w, x) = d(y, z)$

Theorem 1: In any configuration of n points in the plane, not all on a line, there is a line which contains exactly two of the points.

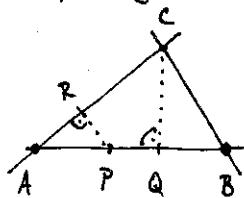
// Introduction (some cool pics and so on...) | \mathbb{P} - set of points, \mathbb{L} - set of lines
Let $\mathbb{A} = \{(P, \ell) \in \mathbb{P} \times \mathbb{L} \mid P \notin \ell\}$

Choose (P_0, ℓ_0) such that $\forall P \in \mathbb{P}, \ell \in \mathbb{L} \nexists (P \in \ell) : d(P_0, \ell_0) \leq d(P, \ell)$

Claim: ℓ_0 is the line we are looking for.

// Show the proof

// When proving that "created" distance is less or equal, use similarity of triangles



$Q, R \notin P : P \in AQ \wedge R \in AC \Rightarrow |\angle CAR| = |\angle RAP|$

from def of $d(P, \ell)$: $|\angle AQC| = |\angle ARP| = 90^\circ$

Hence $\triangle AQC \sim \triangle ARP$

Since $AR \leq AC \wedge AP \leq AQ \Rightarrow PR \leq CR \quad \square \text{QED}$

$\square \text{QED}$

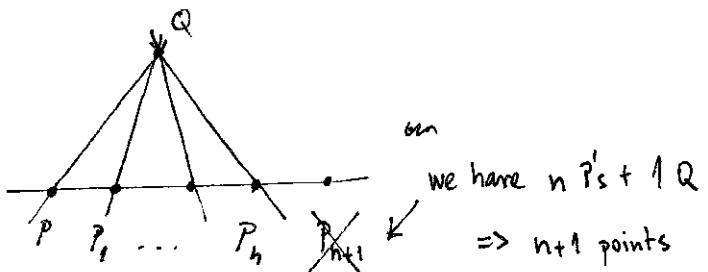
THEOREM 2: Let P be a set of $n \geq 3$ points in the plane, not all on a line. Then the set L of lines passing through at least two points contains at least n points.

|| Very easy - proof by induction

|| There is possibly a mistake in the book

:

□ QED



THEOREM 3: Let X be a set of $n \geq 3$ elements, and let A_1, \dots, A_m be proper subsets of X , such that every pair of elements of X is contained in precisely one set A_i . Then $m \geq n$ holds.

NOTATION: for $x \in X$: r_x be the number of sets A_i containing x .

PROOF 3A: Assumption: $2 \leq r_x \leq m$

If $x \notin A_i \Rightarrow r_x \geq |A_i|$

$$\text{Let } m < n \rightarrow m|A_i| < r_x n \quad / -mn$$

$$m(|A_i|-n) < n(r_x-m) \quad / .(-1)$$

$$m(n-|A_i|) > n(m-r_x)$$

$$\circ 1 = \sum_{x \in X} \frac{1}{n} = \sum_{x \in X} \sum_{A_i: x \notin A_i} \frac{1}{n(m-r_x)} = A$$

$$\circ 1 = \sum_{A_i} \frac{1}{m} = \sum_{A_i} \sum_{x: x \in A_i} \frac{1}{m(n-|A_i|)} = B$$

$$m(n-|A_i|) > n(m-r_x) \Rightarrow A > B \Rightarrow \underline{1 > 1} \text{ contradiction}$$

□ QED

PROOF 3.2 $B = \text{incidence matrix } (x; A_1, \dots, A_m)$

$$\text{where } B_{xi} = \begin{cases} 1 & \leftarrow x \in A_i \\ 0 & \leftarrow x \notin A_i \end{cases}$$

// Don't forget to show examples!

$$BB^T = \begin{pmatrix} r_{x_1} & 1 & \cdots & 1 \\ 1 & r_{x_2} & \cdots & 1 \\ \vdots & \vdots & \ddots & \vdots \\ 1 & 1 & \cdots & r_{x_n} \end{pmatrix} = \underbrace{\begin{pmatrix} r_{x_1}-1 & 0 & \cdots & 0 \\ 0 & r_{x_2}-1 & \cdots & 0 \\ \vdots & \ddots & \ddots & \vdots \\ 0 & 0 & \cdots & r_{x_n}-1 \end{pmatrix}}_{\text{positive definite: } \lambda \geq 0} + \underbrace{\begin{pmatrix} 1 & 1 & \cdots & 1 \\ 1 & 1 & \cdots & 1 \\ \vdots & \vdots & \ddots & \vdots \\ 1 & 1 & \cdots & 1 \end{pmatrix}}_{\text{semi-definite: } \lambda \in \{n, 0\}}$$

BB^T is positive definite $\Rightarrow \text{rank}(BB^T) = n$

Hence $\text{rank}(B)$ is at least $n \Rightarrow n \leq m$ $\square \text{QED}$

NOTES :

Eigen vectors and eigenvalues for A (square matrix)

$$A \cdot v = \lambda v \quad \text{"characteristic value"}$$

$$A = \begin{pmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{pmatrix} \rightarrow \det \begin{pmatrix} 1-\lambda & 1 & 1 \\ 1 & 1-\lambda & 1 \\ 1 & 1 & 1-\lambda \end{pmatrix} = 0 \quad \det(A_{n \times n} - I_n \cdot \lambda) = 0$$

$$(1-\lambda)^3 + 2 - 3(1-\lambda) = 0 \quad \text{only for square}$$

$$1 - 3\lambda + 3\lambda^2 - \lambda^3 + 2 - 3 + 3\lambda = 0$$

$$\lambda^3 - 3\lambda^2 = 0$$

$$\lambda^2(\lambda - 3) = 0 \quad \boxed{\lambda = 0} \quad \boxed{\lambda = 3 = n}$$

$$A = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \rightarrow \det \begin{pmatrix} 1-\lambda & 0 & 0 \\ 0 & 1-\lambda & 0 \\ 0 & 0 & 1-\lambda \end{pmatrix} = 0$$

$$(1-\lambda)^3 = 0 \rightarrow \boxed{\lambda = 1 \geq 0} : \lambda = n-2$$

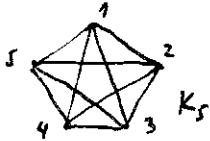
$$A = \begin{pmatrix} 2 & 1 & 1 \\ 1 & 2 & 1 \\ 1 & 1 & 2 \end{pmatrix} \rightarrow \det \begin{pmatrix} 2-\lambda & 1 & 1 \\ 1 & 2-\lambda & 1 \\ 1 & 1 & 2-\lambda \end{pmatrix} = 0$$

$$(2-\lambda)^3 + 2 - 3(2-\lambda) = 0$$

$$\lambda^3 - 6\lambda^2 + 9\lambda - 4 = 0 \quad \boxed{\lambda_1 = 4 \geq 0} \quad \boxed{\lambda_{2,3} = 1 \geq 0}$$

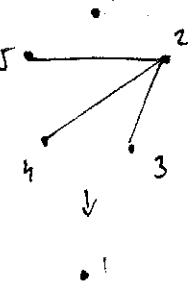
THEOREM 3 : If we decompose a complete graph K_n into m cliques different from K_n , such that every edge is in a unique clique, then $m \geq n$

PROOF 3.3 : $X \sim K_n \vee A_i \sim \text{cliques} \quad [K_n = (V, E)]$



How to decompose graph:

- number the vertices $\{1 \dots n\}$
- take the complete bipartite graph joining 1 to all other vertices $\{2 \dots n\}$
→ we obtain graph $K_{1,n-1}$ (star)
- repeat joining for $j \in \{2 \dots n\}$ with $\{j \dots n\}$
→ we obtain graphs $K_{1,n-2} \dots K_{1,1}$



This decomposition uses $n-1$ complete bipartite graphs

$$\text{Hence } m = \sum_{i=1}^{n-1} i = \frac{1}{2}n(n-1) \geq n \quad (n \geq 3) \quad \square \text{QED}$$

Proof of ~~the~~ claim above : (induction)

- Base : $n=3$: $3 \geq 3$ holds
- 1. Hypothesis : $\frac{1}{2}n(n-1) \geq n$
- 1. Step : $m=n+1 : \frac{1}{2}m(m-1) = \frac{1}{2}n(n+1) = \frac{1}{2}(n^2+n) = \frac{1}{2}(n^2-n+2n)$
 $= \underline{\frac{1}{2}(n-1)n + n \geq n+1 = m} \quad (n \geq 3) \quad \square \text{QED}$



THEOREM 4: If K_n is decomposed into complete bipartite ~~graphs~~ subgraphs H_1, \dots, H_m ,
then $m \geq n+1$

V.

PROOF 4: $K_n = (V, E) \quad V = \{1, \dots, n\}$

Let R_j, L_j be the defining vertex sets of complete bipartite graph $H_j \quad j \in \{1, \dots, m\}$

To every vertex i we assoc. var. x_i .

H_1, \dots, H_m decompose K_n :

$$(1) \quad \sum_{i < j \leq n} x_i \cdot x_j = \sum_{k=1}^m \left(\underbrace{\sum_{a \in L_k} x_a}_{0} \cdot \underbrace{\sum_{b \in R_k} x_b}_{0} \right)$$

Suppose that $m < n+1$

Then $x_1 + \dots + x_n = 0$

$$\sum_{a \in L_k} x_a = 0 \quad (k=1, \dots, m)$$

has fewer equations than variables and hence there exist
a non-trivial solution $c = (c_1, \dots, c_n)$

$$\sum_{i < j \leq n} c_i \cdot c_j = 0 \quad *$$

$$0 = (c_1 + \dots + c_n)^2 = \sum_{i=1}^n c_i^2 + 2 \sum_{i < j \leq n} c_i \cdot c_j = \sum_{i=1}^n c_i^2 > 0 \quad (\text{contradiction})$$

$$*\sum_{i < j \leq n} c_i \cdot c_j = \sum_{k=1}^m \left(\underbrace{\sum_{a \in L_k} c_a}_{0} \cdot \underbrace{\sum_{b \in R_k} c_b}_{0} \right) = 0$$

□ QED