

THE PIGEON HOLE PRINCIPLE

If n objects are placed in r boxes, where $n < r$, then at least one of the boxes contains more than one object.

□ □ □ □ □ $n=6$
 $r=5$ "obvious observation"

N, R $|N|=n > r=|R|$

$f: N \rightarrow R$ $(\exists a \in R)(|f^{-1}(a)| \geq 2)$

stronger: $|f^{-1}(a)| \geq \lceil \frac{n}{r} \rceil$

Contr. $\forall a \in R: |f^{-1}(a)| < \frac{n}{r}$

$n < n \frac{n}{r}$

$n < n$ contradiction

① Numbers

Claim: Consider numbers $1, 2, \dots, 2n$ and take any $n+1$ of them.

Then there are two among those $n+1$ numbers which are coprime.

Proof: Obviously, there are some of the numbers which are consecutive.

(If it wasn't true, we'd need $\geq 2n+2$ numbers) ~~which~~

$(\forall k \in \mathbb{N})(\text{GCD}(k, k+1) = 1)$

claim transformation: ... such that one divides the other.

$A \subseteq \{1, 2, \dots, 2n\}$ $|A| = n+1$

$a \in A$ $a = 2^k \underline{m}$ \underline{n} different odd parts

□ QED

② Sequences

Claim: In any sequence a_1, a_2, \dots, a_{m+1} of $(m+1)$ distinct \mathbb{R} numbers,

there exists an increasing subsequence $a_{i_1} < a_{i_2} < \dots < a_{i_{m+1}}$ ($i_1 < i_2 < \dots < i_{m+1}$)

of length $m+1$ or a decreasing subsequence $a_{j_1} > a_{j_2} > \dots > a_{j_{m+1}}$ ($j_1 < j_2 < \dots < j_{m+1}$)

of length $m+1$ or both.

Proof: $a_i \mapsto t_i$ t_i is the length of the longest inc. subsequence starting at a_i

(i) $t_i \geq m+1$

(ii) $\forall i: t_i \leq m$

$$f: \overbrace{\{a_1, a_2, \dots, a_{m+1}\}}^P \rightarrow \overbrace{\{1, 2, \dots, m\}}^Q$$

$\exists s \in Q$ $f(a_i) = s$ for $\frac{nm}{m} + 1 = \underline{n+1}$

$a_{j_1}, a_{j_2}, \dots, a_{j_{n+1}}$ ($j_1 < j_2 < \dots < j_{n+1}$)

Cons. $a_{j_i}, a_{j_{i+1}} : \underline{a_j < a_{j+1}}$

This is a contradiction. Or it (at least) leads to some.

□ QED

③ SUMS

Claim: Suppose we are given n integers a_1, \dots, a_n , which needn't be distinct. Then there is always a set of consecutive numbers $a_{k+1}, a_{k+2}, \dots, a_k$ whose sum $\sum_{i=k+1}^k a_i$ is a multiple of n .

Proof:

$N = \{0, 1, \dots, n\}$

$R = \{0, 1, 2, \dots, n-1\}$

~~*~~ $f: N \rightarrow R$

$f(m)$ is the remainder of $a_1 + a_2 + \dots + a_m$ over n .

Since $|N| = n+1 > n = |R|$

$a_1 + \dots + a_k, a_1 + \dots + a_k + \dots + a_l$ ($k < l$)

have the same remainder.

$$\sum_{i=1}^l a_i - \sum_{i=1}^k a_i = \sum_{i=k+1}^l a_i$$

this is a multiple

this is too

(we are considering an empty sum ($k=0$) too!)

□ QED

ALGORITHM

Problem: Given n points P_1, P_2, \dots, P_n on a line, find a pair of them with the ^{maximal} distance such that ~~there~~ there is no other point lying between them. (Input is not in sorted order)

Distances: $\langle \pi_1, \dots, \pi_n \rangle$ Reordering of input seq. in sorted order.

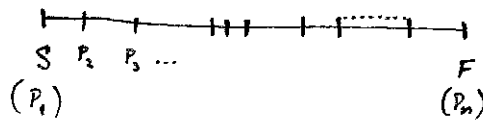
• Sorting can be done in $O(n \log n)$, but this problem can be solved in $O(n)$

Observation: max. dist. should be ^{at least} ~~at least~~ $\frac{\pi_n - \pi_1}{n}$ (obviously)

- can be proved by contradiction

$$\pi_n - \pi_1 = \sum_{i=1}^{n-1} \pi_{i+1} - \pi_i < \sum_{i=1}^{n-1} \dots \dots \dots \text{Meh, } \square \text{ QED}$$

$$\frac{\pi_n - \pi_1}{n-1} \leq \max$$



$\langle a_1, a_2 \rangle, \langle a_2, a_3 \rangle, \dots, \langle a_{n-1}, a_n \rangle$ distances

Our algorithm will remember left-most and right-most points ~~for~~.

left [1...n-1]

right [1...n-1]

(too many lines in the picture using 2 bit colors, I've got only 1 bit ~~!!~~)

right-most point of an interval

↑ length

left-most p. of the next non-empty interval

can be done in $O(n)$