Real-Time Scheduling

Resource Access Control

[Some parts of this lecture are based on a real-time systems course of Colin Perkins

http://csperkins.org/teaching/rtes/index.html]

Current Assumptions

- Single processor
- Individual jobs
  (that possibly belong to periodic/aperiodic/sporadic tasks)
  - Jobs can be preempted at any time and never suspend themselves
- Jobs are scheduled using a priority-driven algorithm
  i.e., jobs are assigned priorities, scheduler executes jobs according to these priorities
- n resources $R_1, \ldots, R_n$ of distinct types
  - used in non-preemptable and mutually exclusive manner; *serially reusable*

Motivation & Notation

Resources may represent:

- Hardware devices such as sensors and actuators
- Disk or memory capacity, buffer space
- Software resources: locks, queues, mutexes etc.

Assume a lock-based concurrency control mechanism

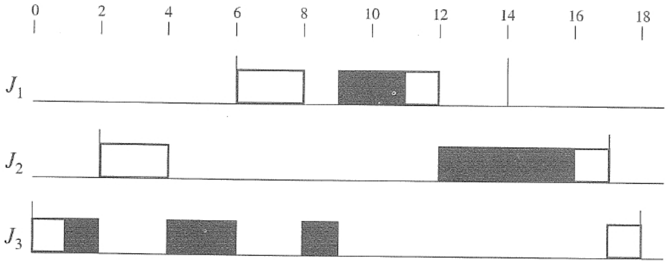
- A job wanting to use a resource R_k executes $L(R_k)$ to lock the resource R_k
- When the job is finished with the resource R_k, unlocks this resource by executing $U(R_k)$
- If lock request fails, the requesting job is **blocked** and has to wait, when the requested resource becomes available, it is unblocked

  In particular, a job holding a lock cannot be preempted by a higher priority job needing that lock

The segment of a job that begins at a lock and ends at a matching unlock is a *critical section* (CS)

- CS must be properly nested if a job needs multiple resources

Example



J_1, J_2, J_3 scheduled according to EDF.

- At 0, J_3 is ready and executes
- At 1, J_3 executes $L(R)$ and is granted R
- J_2 is released at 2, preempts J_3 and begins to execute
- At 4, J_2 executes $L(R)$, becomes blocked, J_3 executes
- At 6, J_1 becomes ready, preempts J_3 and begins to execute
- At 8, J_1 executes $L(R)$, becomes blocked, and J_3 executes
- At 9, J_3 executes $U(R)$ and both J_1 and J_2 are unblocked. J_1 has higher priority than J_2 and executes
- At 11, J_1 executes $U(R)$ and continues executing

Priority Inversion

Definition 1
Priority inversion occurs when

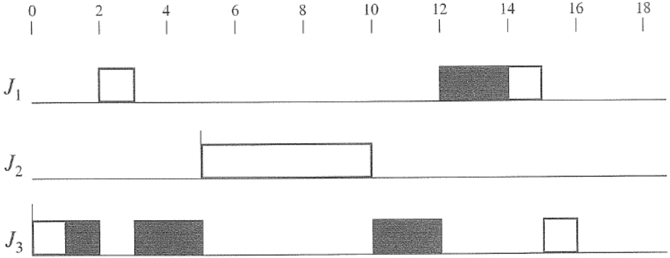
- a high priority job
- is blocked by a low priority job
- which is subsequently preempted by a medium priority job

Then effectively the medium priority job executes with higher priority than the high priority job even though they do not contend for resources

There may be arbitrarily many medium priority jobs that preempt the low priority job $\Rightarrow$ uncontrolled priority inversion

Priority Inversion – Example

Uncontrolled priority inversion:



High priority job (J_1) can be blocked by low priority job (J_3) for unknown amount of time depending on middle priority jobs (J_2)

Deadlock

Definition 2 (suitable for resource access control)

A deadlock occurs when there is a set of jobs $\mathcal{D}$ such that each job of $\mathcal{D}$ is waiting for a resource previously allocated by another job of $\mathcal{D}$.

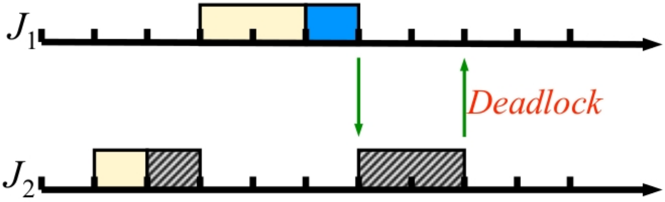
Deadlocks can be

- *detected:* regularly check for deadlock, e.g. search for cycles in a resource allocation graph regularly
- *avoided:* postpone unsafe requests for resources even though they are available (banker's algorithm, priority-ceiling protocol)
- *prevented:* many methods invalidating sufficient conditions for deadlock (e.g., impose locking order on resources, force jobs to release all resources before locking a new one, ...)

See your operating systems course for more information ....

Deadlock – Example

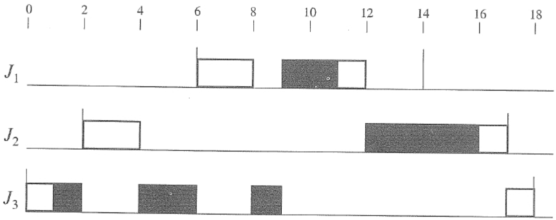
Deadlock can result from piecemeal acquisition of resources: classic example of two jobs J_1 and J_2 both needing both resources R and R'



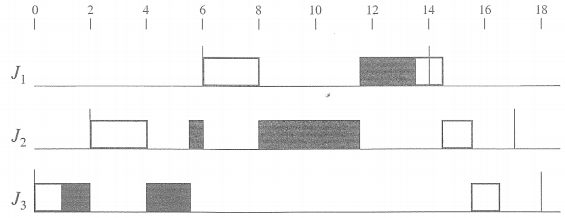
- J_2 locks R' and J_1 locks R
- J_1 tries to get R' and is blocked
- J_2 tries to get R and is blocked

Timing Anomalies due to Resources

Previous example, the critical section of J_3 has length 4



... the critical section of J_3 shortened to 2.5



... but response of J_1 becomes longer!

Controlling Timing Anomalies

Contention for resources causes timing anomalies, priority inversion and deadlock

Several protocols exist to control the anomalies

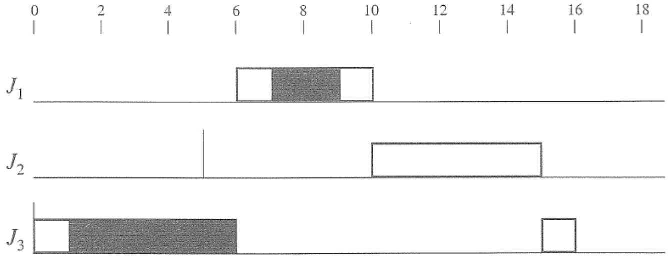
- Non-preemptive CS
- Priority inheritance protocol
- Priority ceiling protocol
- ....

Non-preemptive Critical Sections

The protocol: when a job locks a resource, it is scheduled with priority higher than all other jobs (i.e., is non-preemptive)

Example 3

Jobs J_1, J_2, J_3 with release times $2, 5, 0$, resp., and with execution times $4, 5, 7$, resp.



Non-preemptive Critical Sections

Features:

- no deadlock because no job holding a resource is ever preempted
- no priority inversion:
  - a job J_h can be blocked (by a lower priority job) at release time
  - if J_h is executing, then it cannot be blocked by a lower priority job (because all resources are free)
  - if a job J_h is blocked, then once the blocking critical section completes, no lower priority job can block J_h (since all resources are free)
  - ... thus any job can be blocked only once, at release time, blocking time is bounded by duration of one critical section of a lower priority job

Advantage: very simple; easy to implement both in fixed and dynamic priority; no prior knowledge of resource demands of jobs needed

Disadvantage: every job can be blocked by every lower-priority job with a critical section, even if there is no resource conflict

Priority-Inheritance Protocol

Idea: adjust the scheduling priorities of jobs during resource access, to reduce the duration of timing anomalies

(As opposed to non-preemptive CS protocol, this time the priority is not always increased to maximum)

Notation:

- *assigned priority* = priority assigned to a job according to a standard scheduling algorithm
- At any time t, each ready job J_k is scheduled and executes at its *current priority* $\pi_k(t)$ which may differ from its assigned priority and may vary with time
  - The current priority $\pi_k(t)$ of a job J_k may be raised to the higher priority $\pi_h(t)$ of another job J_h
  - In such a situation, the lower-priority job J_k is said to *inherit* the priority of the higher-priority job J_h, and J_k executes at its inherited priority $\pi_h(t)$

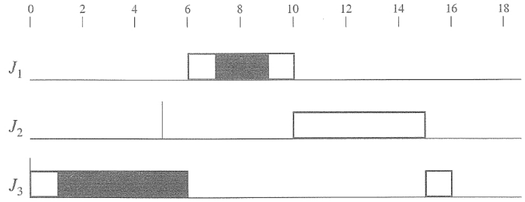
Priority-Inheritance Protocol

- **Scheduling rules**:
  - Jobs are scheduled in a preemptable priority-driven manner *according to their current priorities*
  - At release time, the current priority of a job is equal to its assigned priority
  - The current priority remains equal to the assigned priority, except when the priority-inheritance rule is invoked
- **Priority-inheritance rule**:
  - When a job J_h becomes blocked on a resource R, the job J_k which blocks J_h inherits the current priority $\pi_h(t)$ of J_h; all priorities are updated transitively
  - J_k executes at its inherited priority until it releases R; at that time, the priority of J_k returns to its priority $\pi_k(t')$ at the time t' when it acquired the resource R

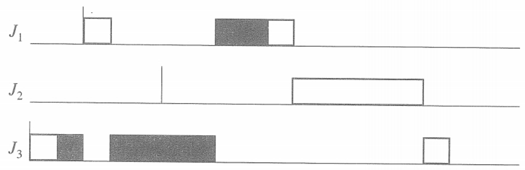
    (note that $\pi_k(t')$ may still be an inherited priority)
- **Resource allocation**: when a job J requests a resource R at t:
  - If R is free, R is allocated to J until J releases it
  - If R is not free, the request is denied and J is blocked
  - J is only denied R if the resource is held by another job

Priority-Inheritance Simple Example

non-preemptive CS:

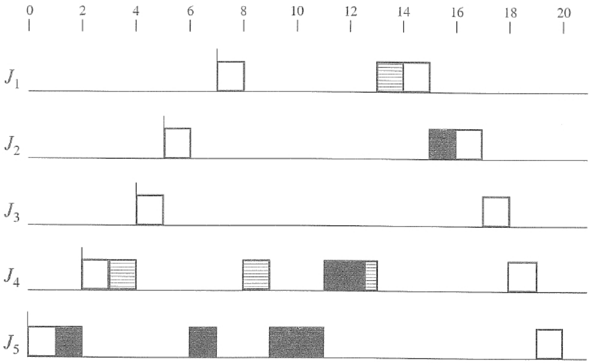


priority-inheritance:



- At 3, J_1 is blocked by J_3, J_3 inherits priority of J_1
- At 5, J_2 is released but cannot preempt J_3 since the inherited priority of J_3 is higher than the (assigned) priority of J_2

Priority-Inheritance Example



- At 0, J_5 starts executing at priority 5, at 1 it executes $L(Black)$
- At 2, J_4 preempts J_5 and executes
- At 3, J_4 executes $L(Shaded)$, J_4 continues to execute
- At 4, J_3 preempts J_4; at 5, J_2 preempts J_3
- At 6, J_2 executes $L(Black)$ and is blocked by J_5. Thus J_5 inherits the priority 2 of J_2 and executes
- At 8, J executes $L(Shaded)$ and is blocked by J. Thus J inherits the

Properties of Priority-Inheritance Protocol

- Simple to implement, does not require prior knowledge of resource requirements
- Jobs exhibit two types of blocking
  - Direct blocking due to resource locks

    i.e., a job J_ℓ locks a resource R, J_h executes $L(R)$ is directly blocked by J_ℓ on R
  - Priority-inheritance blocking

    i.e., a job J_h is preempted by a lower-priority job that inherited a higher priority
- Jobs may exhibit transitive blocking

  In the previous example, at 9, J_5 blocks J_4 and J_4 blocks J_1, hence J_5 inherits the priority of J_1
- Deadlock is *not* prevented

  In the previous example, let J_5 request *shaded* at 6.5, then J_4 and J_5 become deadlocked
- Can reduce blocking time (see next slide) compared to non-preemptable CS but does not guarantee to minimize blocking

Priority-Inheritance – Blocking Time

$z_{\ell,k}$ = the k-th critical section of J_ℓ

A job J_h is blocked by $z_{\ell,k}$ if J_h has higher assigned priority than J_ℓ but has to wait for J_ℓ to exit $z_{\ell,k}$ in order to continue

$\beta^*_{h,\ell}$ = the set of all maximal critical sections $z_{\ell,k}$ that may block J_h
(recall that CS are properly nested, maximal CS which may block J_h is the one which is not contained within any other CS which may block J_h)

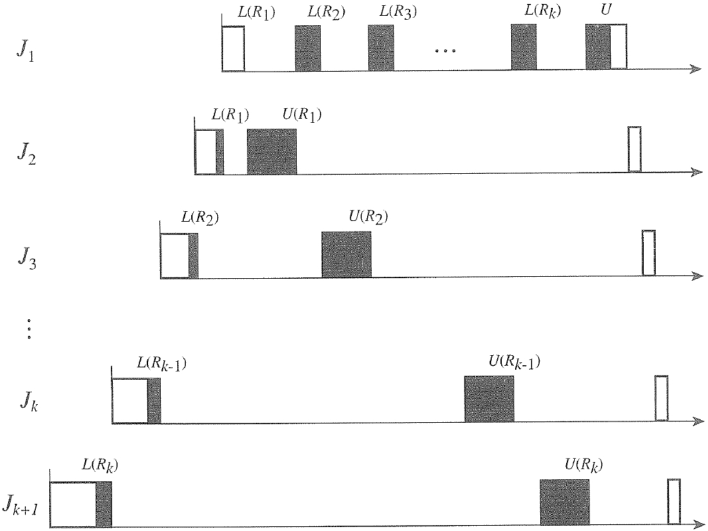
Theorem 4
Let J_h be a job and let $J_{h+1}, \ldots, J_{h+m}$ be jobs with lower priority than J_h. Then J_h can be blocked for at most the duration of one critical section in each *of $\beta^*_{h,\ell}$ where $\ell \in \{h+1, \ldots, h+m\}$.*

Lemma 5

J_h can be blocked by J_ℓ only if J_ℓ is executing within a critical section $z_{\ell,k}$ of $\beta^*_{h,\ell}$ when J_h is released

- Assume that J_h is released at t and J_ℓ is in no CS of $\beta^*_{h,\ell}$ at t. We show that J_ℓ never executes between t and completion of J_h:
  - If J_ℓ is not in any CS at t, then its current priority at t is equal to its assigned priority and cannot increase. Thus, J_ℓ has to wait for completion of J_h as the current priority of J_h is always higher than the assigned priority of J_ℓ.
  - If J_ℓ is still in a CS at t, then this CS does not belong to $\beta^*_{h,\ell}$ and thus cannot block J_h before completion and cannot execute before completion of J_h.
- Assume that J_ℓ leaves $z_{\ell,k} \in \beta^*_{h,\ell}$ at time t. We show that J_ℓ never executes between t and completion of J_h:
  - If J_ℓ is not in any CS at t, then, as above, J_ℓ never executes before completion of J_h and cannot block J_h.
  - If J_ℓ is still in a CS at t, then this CS does not belong to $\beta^*_{h,\ell}$ because otherwise $z_{\ell,k}$ would not be maximal. Thus J_ℓ cannot block J_h, and thus J_ℓ is never executed before completion of J_h.

Priority-Inheritance – The Worst Case



Priority-Ceiling Protocol

The goal: to furhter reduce blocking times due to resource contention and to prevent deadlock

- in its basic form priority-ceiling protocol works under the assumption that the priorities of jobs and resources required by all jobs are known apriori

  can be extended to dynamic priority (job-level fixed priority), see later

Notation:

- The *priority ceiling* of any resource R_k is the highest priority of all the jobs that require R_k and is denoted by $\Pi(R_k)$
- At any time t, the current priority ceiling $\Pi(t)$ of the system is equal to the highest priority ceiling of the resources that are in use at the time
- If all resources are free, $\Pi(t)$ is equal to Ω, a newly introduced priority level that is lower than the lowest priority level of all jobs

Priority-Ceiling Protocol

The scheduling and priority-inheritance rules are the same as for priority-inheritance protocol

- **Scheduling rules**:
  - Jobs are scheduled in a preemptable priority-driven manner *according to their current priorities*
  - At release time, the current priority of a job is equal to its assigned priority
  - The current priority remains equal to the assigned priority, except when the priority-inheritance rule is invoked
- **Priority-inheritance rule**:
  - When job J_h becomes blocked on a resource R, the job J_k which blocks J_h inherits the current priority $\pi_h(t)$ of J_h; also priorities are inherited transitively
  - J_k executes at its inherited priority until it releases R; at that time, the priority of J_k returns to its priority $\pi_k(t')$ at the time t' when it acquired the resource R
    (note that $\pi_k(t')$ may still be an inherited priority)

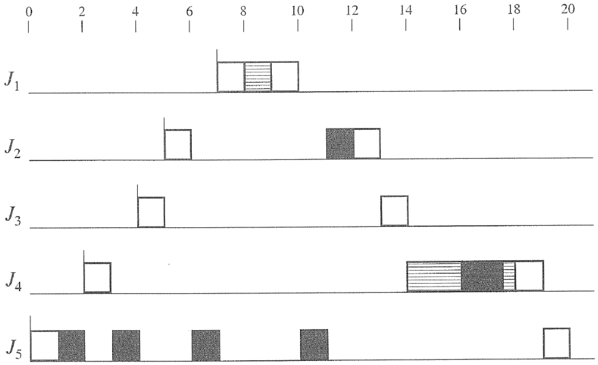
Priority-Ceiling Protocol

Resource allocation rules:

- When a job J requests a resource R held by another job, the request fails and the requesting job blocks
- When a job J requests a resource R at time t, and that resource is free:
  - If J's priority $\pi(t)$ is higher than current priority ceiling $\Pi(t)$, R is allocated to J
  - If J's priority $\pi(t)$ is not higher than $\Pi(t)$, R is allocated to J only if J is the job holding the resource(s) whose priority ceiling is equal to $\Pi(t)$, otherwise J is blocked
    (Note that only one job may hold the resources whose priority ceiling is equal to $\Pi(t)$)

Note that unlike priority-inheritance protocol, the priority-ceiling protocol can deny access to an available resource

Priority-Ceiling Protocol



- At 1, $\Pi(t) = \Omega$, J_5 executes $L(Black)$, continues executing
- At 3, $\Pi(t) = 2$, J_4 executes $L(Shaded)$; because the ceiling of the system $\Pi(t)$ is higher than the current priority of J_4, job J_4 is blocked, J_5 inherits J_4's priority and executes at priority 4
- At 4, J_3 preempts J_5; at 5, J_2 preempts J_3. At 6, J_2 requests *Black* and is directly blocked by J_5. Consequently, J_5 inherits priority 2 and executes until preempted by J_1

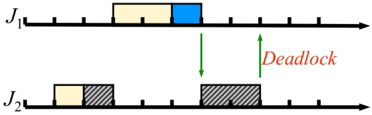
Priority-Ceiling Protocol

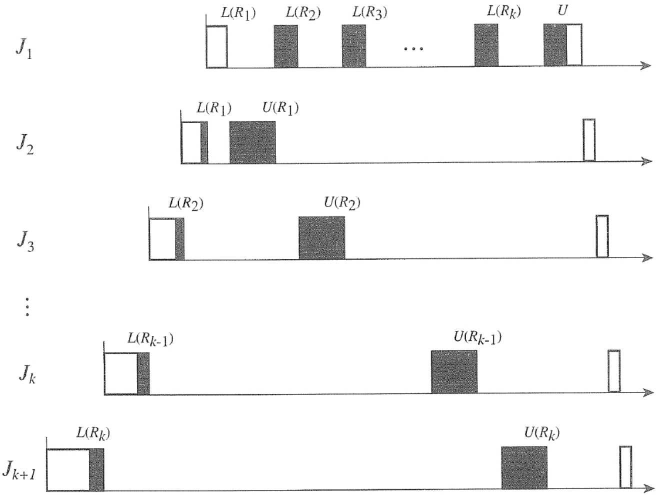
Theorem 6

Assume a system of preemptable jobs with fixed assigned priorities. Then

- *deadlock may never occur,*
- *a job can be blocked for at most the duration of one critical section.*

These situations cannot occur with priority ceiling protocol:





Differences between the priority-inheritance and priority-ceiling

- Priority-inheritance is greedy, while priority ceiling is not

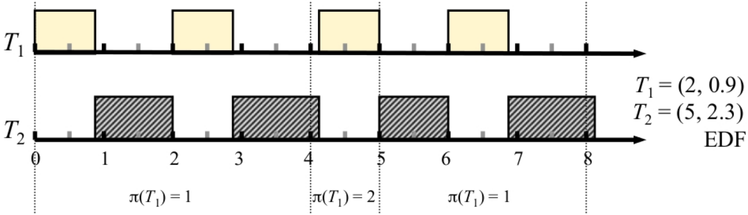
  The priority-ceiling protocol may withhold access to a free resource, i.e., a job can be blocked by a lower-priority job which does not hold the requested resource – *avoidance blocking*
- The priority ceiling protocol forces a fixed order onto resource accesses thus eliminating deadlock

Resources in Dynamic Priority Systems

The priority ceiling protocol assumes fixed and known priorities

In a dynamic priority system, the priorities of the periodic tasks change over time, while the set of resources is required by each task remains constant

- As a consequence, the priority ceiling of each resource changes over time



What happens if T_1 uses resource X, but T_2 does not?

- Priority ceiling of X is 1 for $0 \leq t \leq 4$, becomes 2 for $4 \leq t \leq 5$, etc. even though the set of resources is required by the tasks remains unchanged

Resources in Dynamic Priority Systems

- If a system is job-level fixed priority, but task-level dynamic priority, a priority ceiling protocol can still be applied
  - Each job in a task has a fixed priority once it is scheduled, but may be scheduled at different priority to other jobs in the task (e.g. EDF)
  - Update the priority ceilings of all resources each time a new job is introduced; use until updated on next job release
- Has been proven to prevent deadlocks and no job is ever blocked for longer than the length of one critical section
  - But: very inefficient, since priority ceilings updated frequently
  - May be better to use priority inheritance, accept longer blocking

Schedulability Tests with Resources

How to adjust schedulability tests?

Add the blocking times to execution times of jobs; then run the test as normal

The blocking time b_i of a job J_i can be determined for all three protocols:

- non-preemptable CS $\Rightarrow$ b_i is bounded by the maximum length of a critical section in lower priority jobs
- priority-inheritance $\Rightarrow$ b_i is bounded by the total length of the m longest critical sections where m is the number of jobs that may block J_i

  (For a more precise formulation see Theorem 2.)
- priority-ceiling $\Rightarrow$ b_i is bounded by the maximum length of a critical section