PA197 Secure network design

Message security and key management

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Laboratory

- Protection of packets
 - Encrypt and MAC (AES), Speed? Latency? Replay?
- µTesla: Authenticated Broadcast
 - How it works
 - Implement verification on nodes
 - Try to attack protocol
 - Considerations

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Message confidentiality

- Encryption/MAC available in radio module vs. software
 - E.g., TelosB' CC2420 radio module has native AES support
 - JeeNode's RF12 radio module has no native encryption support
 - packet needs to be encrypted before passed to rf12_sendNow
 - rf12_encrypt(key): XXTEA, 128bits code, sequence num. up to 30 bits
 - http://jeelabs.org/2010/02/23/secure-transmissions/
- Software encryption with AES algorithm (IS \rightarrow encryptAES.zip)
 - Task: Measure approximate speed of encryption (ms/block)
 - Note: Implementation optimized for memory, not for speed
- Which mode to use?
 - CBC used often (be aware of trim attack)
 - CTR mode (possibility for precomputation => lower latency)
- Conclusion: even on simple nodes, encryption is possible
 - < 1ms / block

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Integrity protection (MAC)

- What algorithm to use for integrity checksum?
 - HMAC-SHA3 is great (use it in ordinary programs), but:
 - Requires additional code (for SHA-3)
 - Requires relatively large internal state (RAM, 64B)
 - If AES is already used (encryption), CBC-MAC can be utilized
 - Use it properly! http://blog.cryptographyengineering.com/2013/02/why-i-hate-cbc-mac.h tml
- What is proper length of MAC tag?
 - Longer packets => more energy consumed & higher chance of collision
 - 16B? 4B? Birthday paradox? Number of messages?
 - How can attacker misuse two messages with same MAC?
- Task: Estimate latency introduced by encryption and MAC
 - Scenarios: node-to-node enc, end-to-end enc
 - MAC verification: every hop, end node only
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Problem of packet replay

- How to deal with packet replay?
 - ... in noisy environment with natural packet loss
- MAC-chaining
 - MAC value from previous packet used as IV for next one
 - lost packet => impossible verification of next one
- Incremental counter
 - Possibility to detect and recover even when packet lost
- Think about advantages and disadvantages

µTesla: Authenticated Broadcast

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µTesla: Authenticated Broadcast

- 1. BS: generate new MACKeyi for next periodi
- 2. BS: compute MAC on message Mi to be send
- 3. Send message Mi via flood (code from last week)
 - Or directly by single hop (strong transmission)
- 4. Node: receive and store message Mi on nodes
- 5. BS: After end of periodi, flood MACKeyi to nodes
- 6. Node: Verify validity of MACKeyi against stored root
- 7. Node: Verify MAC on message Mi

Lab work: µTesla (client-only)

- 1. BS: generate new MACKeyi for next periodi
- 2. BS: compute MAC on message Mi to be send
- 3. Send message Mi via flood (code from last week)
- 4. Node: receive and store message Mi on nodes
- 5. BS: After end of periodi, flood MACKeyi to nodes
- 6. Node: Verify validity of MACKeyi against root
- 7. Node: Verify MAC on message Mi
- •. Implement node's reception and verification
 - Display received message and result of verification
 - Notify if something is wrong with MAC

µTesla details

- BS uses: rf12_initialize(1, RF12_868MHZ, 123);
- Format of authenticated broadcast message
 "msg"[3B]epoch[2B]message[11B]MAC[4B]
- Root for hash chain is:
 - '99534f2ec8332239741261aaa803a2f73a018e76'
 - seed \rightarrow root is 1001xSHA-1 => 1000 epochs possible
- MAC key is broadcasted after 15 seconds
 Format: "key"[3B]epoch[2B]key[20B]
- MAC computed simply as AES-ECB(msg[16B])
 First 4 bytes used as MAC value
- Hash chain computation function is pre-prepared
 - computeHashes()

µTesla client – how to start

- Capture message emitted by base station
- Parse content
 - Current (claimed) epoch (use memcpy(&epoch, packet+3, 2))
 - Data message
 - MAC value
 - MAC key (coming later)
- Verify MAC key against root
- Verify MAC (AES-ECB)
- Try to verify proper epochs etc... (considerations)

µTesla client – what to output

- Print on serial port received data message
 Print parsed values (epoch, data, MAC)
- Print received key message
 Print parsed values (epoch, MAC key)
- Print result of verification of MAC key against root
 MAC key is genuine
- Print result of verification of MAC on data message
 MAC is valid
- Print warning if older than expected epoch is received

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Attacks against µTesla

- Try to attack µTesla protocol for others
 - Be BS, other students may have incomplete implementation
 - Change data message to 'evilYourName' [11B] when trying
 - You win if others will accept your messages
- 1. Replay older intercepted messages
 - When will nodes accept it? What is successful attack?
- Wait for MACKeyi and create/flood forged message
 When will nodes accept it? How to prevent timing issues?
- 3. Impersonate BS and broadcast own messages including MACKey broadcast
 - How to prevent?
 - Pre-distributed root of hash chain inside every node

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Considerations of µTesla

- How should base station keep stable epoch size?
 - Timers need to be used (otherwise epochs will drift)
 - (not covered in this labs)
- How should node behave after reboot?
 - How to know current time (=> period)?
 - Real clocks necessary (millis() just return time from start)
 - How to permanently store last epoch seen?
 - https://www.arduino.cc/en/Reference/EEPROM
- Alternative: broadcast warning on incorrect detected epoch
 - Rebooted node will just listen for several epochs
 - If no warning from other (still running) nodes received, then current epoch is synchronized

No homework 🕲

- Homework Attack against routing still running
- Submit before: 8.5. 23:59am (full number of points)
 - Every additional started day (24h) means 1.5 points penalization