

# IA010: Principles of Programming Languages

## Constraints

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# Declarative programming

Describe **what** you want to compute, not **how**  
(no side-effects, no state)

## Advantages

- easier to reason about
- write separately and compose

## Logic programming

write set of constraints and search for solution

# Single-assignment variables

$\langle expr \rangle ::= \dots \mid \mathbf{let} \langle id \rangle ; \langle expr \rangle$

```
let x;  
let y;  
x := 1;  
x := 1; // ok  
x := 2; // error  
y := x+1;  
  
let add(x,y,z) {  
    z := x+y;  
};  
let u;  
add(1,2,u);
```

```
let reverse(lst, ret) {  
  let iter(lst, acc, ret) {  
    case lst  
    | []      => ret := acc  
    | [x|xs] => iter(xs, [x|acc], ret)  
  };  
  
  iter(lst, [], ret)  
};
```

# Unification

$\langle expr \rangle ::= \dots \mid \langle expr \rangle ::= \langle expr \rangle$

$1 ::= x$	$x := 1$
$x ::= y$	identifies $x$ and $y$
$[x, 2] ::= [1, y]$	$x := 1$ and $y := 2$

# Unification algorithm

solve  $u := v$

- If  $u$  is an uninitialised variable, set it to  $v$ .
- If  $v$  is an uninitialised variable, set it to  $u$ .
- If  $u = m$  and  $v = n$  are numbers, check that  $m = n$ .
- If  $u = c(s_0, \dots, s_{m-1})$  and  $v = d(t_0, \dots, t_{n-1})$  are constructors, check that  $c = d$ ,  $m = n$ , and  $s_i := t_i$ , for all  $i$ .
- If  $u = [l_0 = s_0, \dots, l_{m-1} = s_{m-1}]$  and  $v = [k_0 = t_0, \dots, k_{n-1} = t_{n-1}]$  are records, find bijection  $\varphi : m \rightarrow n$  such that  $l_i = k_{\varphi(i)}$  and  $s_i := t_{\varphi(i)}$ , for all  $i$ .
- In all other cases, fail.

(In particular, we cannot unify function values.)

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## Notes

- two kinds of uninitialised values: unknown value, equal to other variable
- need to prevent infinite loops

# Backtracking

$\langle expr \rangle ::= \dots | \mathbf{choose} | \langle expr \rangle \dots | \langle expr \rangle | \mathbf{fail}$

```
let is_one_or_two(x) {  
  choose  
  | x := 1  
  | x := 2  
};  
  
is_one_or_two(1); // ok  
is_one_or_two(3); // fail
```



# Primitive operations

## **checkpoint k**

- stores the current continuation and machine state

## **rewind**

- fetches the continuation associated with the last checkpoint,
- restores the machine state to its previous state (deleting the last checkpoint),
- and calls the fetched continuation.



# Implementation

- store stack of checkpoints
- each checkpoint contains: continuation, list of modified variables
- `checkpoint k` puts `k` on the stack
- when we set a variable `x`, we add `x` to the top list
- `rewind` pops the stack, unsets all variables in the top list, and calls the stored continuation

# Example

```
edge(a,b).  
edge(b,c).  
trans(X,Y) :- edge(X,Y).  
trans(X,Y) :- edge(X,Z), trans(Z,Y).
```

```
let edge(x,y) {  
  choose  
  | { x := a; y := b; }  
  | { x := b; y := c; }  
}  
let trans(x,y) {  
  choose  
  | edge(x,y)  
  | { let z; edge(x,z); trans(z,y); }  
}
```