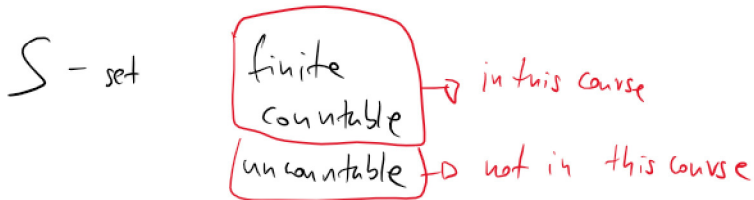


# 1.) A CRASH COURSE IN PROBABILITY THEORY

## 2.) RANDOMIZED QSORT

1.) **PROBABILITY SPACE** ~ A set of all possible outcomes of a random experiment.



**EXAMPLE** - a set of all  $n$ -bit strings

2.) **EVENTS** -  $E \subseteq S$

**Example** - a string with exactly 3 symbols '1'

## 3.) **PROBABILITY FUNCTION**

$$P: S \rightarrow [0,1]$$

$$\sum_{i \in S} P(i) = 1$$

**Example:** - Uniform distribution  $\forall x \in \{0,1\}^n$   $P(x) = \frac{1}{2^n}$

$$P(E) = \sum_{i \in E} P(i)$$

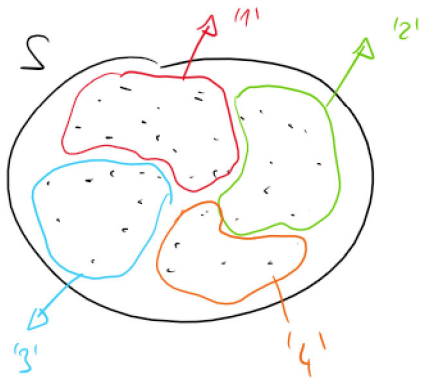
What is the probability to obtain a 5-bit string with exactly 3 symbols '1' (call it event  $E$ )

$$P(E) = \binom{5}{3} = 10 \cdot \frac{1}{2^5} = \frac{10}{32} \checkmark$$

## **RANDOM VARIABLES**

$$(X, Y, Z)$$

$$X: S \rightarrow \mathbb{R}$$



Essentially  $X$  is a division of probability space into mutually exclusive and collectively exhaustive set of events.

**EXAMPLE**  $X$  is the number of symbols '1' in  $n$ -bit string

**Q** For  $n=4$ , what is the distribution of  $X$

$$\Pr[X=0] = \frac{1}{16}$$

$$\Pr[X > 1] = \frac{11}{16}$$

$$\Pr[X=1] = \frac{4}{16}$$

$$\Pr[X \geq 2 \wedge X < 4] = \frac{10}{16}$$

$$\Pr[X=2] = \frac{6}{16}$$

$$\Pr[X=3] = \frac{4}{16}$$

$$\Pr[X=4] = \frac{1}{16}$$

### EXPECTATION OF RANDOM VARIABLES

$$E(X) = \sum_{\substack{i \in \mathbb{R} \\ i \in \text{Im}(X)}} i \cdot \Pr(X=i)$$

**EXAMPLE** :  $E(X) = \sum_{i=0}^4 i \cdot \Pr[X=i] = 0 \cdot \Pr[X=0] + 1 \cdot \Pr[X=1] + 2 \cdot \Pr[X=2] + 3 \cdot \Pr[X=3] + 4 \cdot \Pr[X=4]$

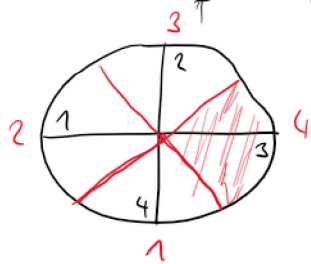
$$= 0 + \frac{4}{16} + \frac{12}{16} + \frac{12}{16} + \frac{4}{16}$$

$$= 2$$

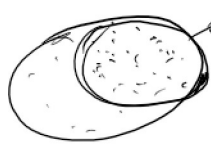
### CONDITIONAL PROBABILITIES

GIVEN 2 random variables  $X$  and  $Y$  over the same random experiment define conditional probability of  $X$  given  $Y = y$

$$Pr(X=i|Y=j) = \frac{Pr(X=i, Y=j)}{Pr(Y=j)} \quad (Pr(Y=j) \neq 0)$$



Intuitively we are creating a new probability space equal to an event  $(Y=j)$



$\frac{1}{Pr(Y=j)}$  renormalizes the original probabilities

### Examples

$$S = \{0,1\}^4$$

$X$  - number of '1'

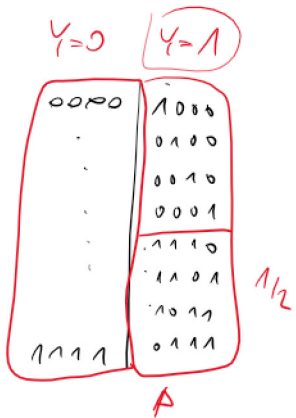
$Y$  - parity of the string  $\begin{cases} \text{even} = 0 \\ \text{odd} = 1 \end{cases}$

$$Pr\{Y=0\} = 1/2$$

$$Pr\{Y=1\} = 1/2$$

$$Pr\{X=3|Y=0\} = 0 = \frac{Pr\{X=3, Y=0\}}{Pr\{Y=0\}}$$

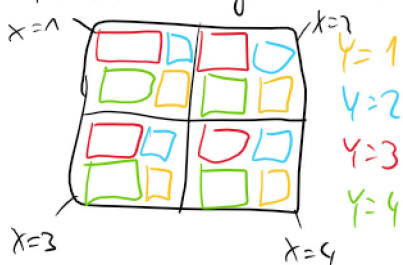
$$Pr\{X=3|Y=1\} = \frac{Pr\{X=3, Y=1\}}{Pr\{Y=1\}} = \frac{1/16}{1/2} = \frac{1}{2}$$



### INDEPENDENCE

$X$  and  $Y$  are independent if for all  $i, j$

$$Pr(X=i|Y=j) = Pr(X=i)$$



### EXAMPLE

ARE  $Y$  and  $X$  from the previous example independent?

$n \rightarrow$

$$X=4$$

### EXAMPLE

ARE  $Y$  and  $X$  from the previous example independent?

$X$  - number of '1'

$Y$  - parity

$$\# \Pr(X=3) = 1/4$$

$$\Pr(X=3 | Y=0) = 0$$

$Z$  is the value of the first bit of  $\{0, 1\}^4$   
is  $Z$  and  $Y$  independent?

$$\Pr(Z=1) = 1/2$$

$$\Pr(Z=0) = 1/2$$

$$\Pr(Z=1 | Y=1) = 1/2$$

$$\Pr(Z=0 | Y=1) = 1/2$$

$$\Pr(Z=1 | Y=0) = 1/2$$

$$\Pr(Z=0 | Y=0) = 1/2$$

### LINEARITY OF EXPECTATION

$$W = X + Y + Z$$

$W$  is a well defined r.v.

$$E(W) = E(X + Y + Z) = E(X) + E(Y) + E(Z)$$

$$E\left(\sum_i a_i X_i\right) = \sum_i a_i E(X_i)$$

$$E\left(\sum_i X_i\right) = \sum_i E(X_i) \quad \text{this is often much easier to calculate}$$

$$E(W) = \sum_{i,j,k} (i+j+k) \cdot \Pr(X=i, Y=j, Z=k)$$

$$= E(X) + E(Y) + E(Z)$$

$$= 2 + 1/2 + 1/2 = 3$$

$$E(X_1 \cdot X_2) \neq E(X_1) \cdot E(X_2)$$

## THE LAW OF TOTAL PROBABILITY

r.v.  $X$  and  $Y$

$$Pr(X=i) = \sum_{j \in \text{Im}(Y)} Pr(X=i | Y=j) \cdot Pr(Y=j)$$

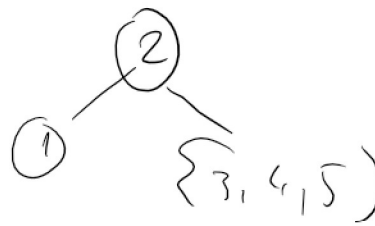
" "  
 $Pr(X=i, Y=j)$

## QUICK SORT

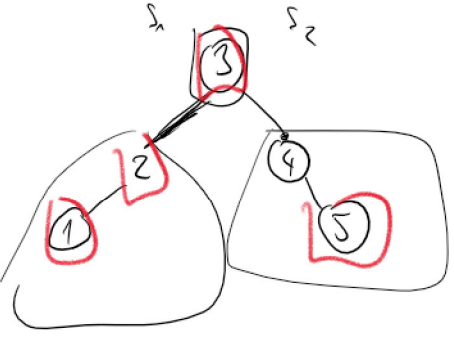
IN: Collection of numbers  $S$   
 OUT: ordered list of elements in  $S$

- 0.) if  $S$  contains a single element output it
- 1.) Choose a **pivot**  $p \in S$  **uniformly at random**
- 2.) Create  $S_1$  which contains all  $s \in S$   $s < p$   
 Create  $S_2$  which contains all  $s \in S$   $s > p$
- 3.) Output  $(\overbrace{\text{quicksort}(S_1)}, p, \overbrace{\text{quicksort}(S_2)})$

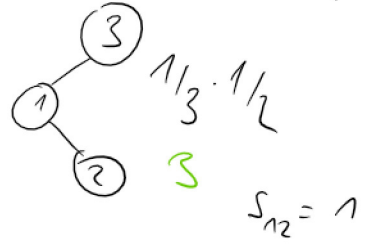
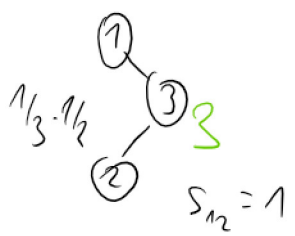
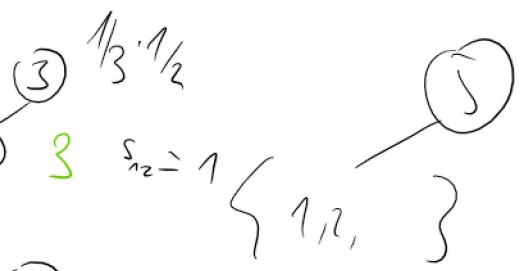
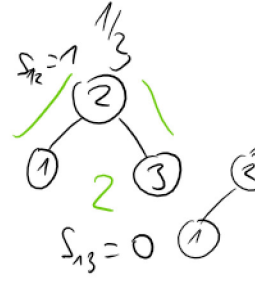
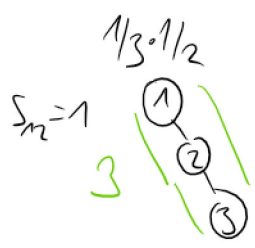
Probability space is a set of ordered trees with  $|S|=n$  nodes



$$S = \{ \underbrace{1, 2}_{S_1}, \underbrace{3, 4, 5}_{S_2} \}$$



$$1/5 \cdot 1/2 = 1/10 =$$



$$1/5 \cdot 1/2 = 1/10$$

$$\textcircled{2} \rightarrow S_{12} = 1$$

$$\textcircled{2} \rightarrow S_{12} = 1$$

$X$  = RV. - number of comparisons in a given tree

$$E(X_{(n=3)}) = \underbrace{1/3 \cdot 1/2 \cdot 3}_{\Pr\{x=3\} \cdot 3} + \underbrace{1/3 \cdot 2}_{\Pr\{x=2\}} = \frac{8}{3}$$

$$E(X_3) \quad E(X_4) \dots \quad E(X_{100})$$

$$f(n) = E(X_n)$$

Q How does scale with  $n$ ?  $O(n \log n)$

$\forall i, j \in S$  define a random variable  $S_{ij}$

$$S_{ij} = \begin{cases} 1 & \text{if } i \text{ and } j \text{ are compared in a given run} \\ 0 & \text{otherwise} \end{cases}$$

$$X = \sum_{i < j} S_{ij} \rightarrow \text{total number of comparisons}$$

$$E(X) = E\left(\sum_{i < j} S_{ij}\right) = \sum_{i < j} E(S_{ij}) \leq \sum_{i < j} P_{ij} \in O(n \log n)$$

$$E(S_{ij}) = 0 \cdot \Pr(S_{ij}=0) + 1 \cdot \Pr(S_{ij}=1) = \Pr(S_{ij}=1) < \frac{P_{ij}}{2}$$

Step 1 What is the probability that  $i$  and  $j$  get compared?

$$\Pr(S_{ij}=1) = \frac{2}{|S|}$$

Step 2  $|S_1|, |S_2|, |S_3|$  Probability of  $i, j$  to get compared in step 2

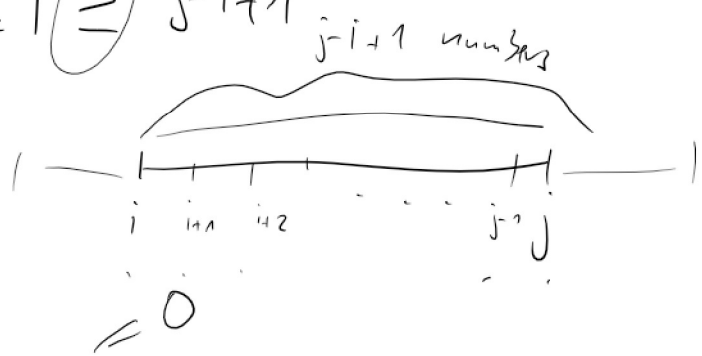
Step 2  $|S_1|, |S_2|$  of  $i, j$  to get compared in step 2

$\Pr(S_{ij}=1 | i < y < j) = 0$  depends on  $y$ .  
 $\Pr(S_{ij}=1 | y < i < j) = \frac{2}{|S_2|} < \frac{2}{j-i-1}$   $i, j \in S_1$   
 $\Pr(S_{ij}=1 | i < j < y) = \frac{2}{|S_1|} < \frac{2}{j-i-1}$

**Step 2** let's bound the probability of comparing

$i, j$  conditioned on the event of both  **$i, j$  being in a same subset**  $S_k$   $|S_k| \geq j-i+1$

$$\frac{2}{|S_k|} < \frac{2}{j-i+1}$$



$$\begin{aligned}
 \Pr(S_{ij}=1) &= \Pr(S_{ij}=1 | \text{they are not in the same set}) \cdot \Pr(\text{not in the same set}) \\
 &\quad + \Pr(S_{ij}=1 | \text{they are in the same set}) \cdot \Pr(\text{they are in the same set}) \\
 &= \Pr(S_{ij}=1 | \text{they are in the same set}) \cdot \Pr(\text{they are in the same set}) \\
 &\leq \Pr(S_{ij}=1 | \text{they are in the same set}) \\
 &\leq \frac{2}{j-i+1}
 \end{aligned}$$

$$E(X) \leq \sum_{i < j} \frac{2}{j-i+1} = \dots = n \cdot \sum_{i=1}^n \frac{1}{i} = \Theta(n \log n)$$

$\downarrow$   
 in the slides  $\Theta(\log n)$