

Algebraic techniques

- Freivald's technique for matrix multiplication
- Polynomial comparison: Schwartz-Zippel theorem
- SQ Lhm. \Rightarrow Freivald's technique

Matrix comparison

Given $n \times n$ matrices A, B and C over a finite field \mathbb{F}_p .

Finite fields are finite set of numbers with well defined multiplication and addition. For prime p $\mathbb{F}_p = \{0, \dots, p-1\}$ and $+, \times \pmod p$.

\mathbb{F}_p exist only for prime power p .

Verify whether

$$A \cdot B = C$$

Multiply $A \cdot B$ and compare to C .

$$O(n^3) \quad [O(n^{2.373})] \quad \leftarrow$$

↑

Suppose you want to check whether your matrix multiplication algorithm works correctly. With randomized technique we can solve the problem in $O(n^2)$.

Alg.

1) Choose $\vec{r} \in \{0,1\}^n$ at random and calculate

$$\frac{A \cdot (\overline{B \cdot \vec{r}})}{O(n^2)} \text{ and } \frac{(\overline{C \cdot \vec{r}})}{O(n^2)} \text{ compare the results } O(n)$$

2) • If the vectors are equal, alg. says matrices are equal ($A \cdot B = C$)
• if not then $A \cdot B \neq C$

3) output NO $\Rightarrow A \cdot B \neq C$ w.p 1

output YES \Rightarrow w.p. smaller or equal to $\frac{1}{2}$.

ANALYSIS:

\rightarrow We can reduce the problem to finding whether

$$D = A \cdot B - C \text{ is identically } 0 \quad D = \begin{pmatrix} 0 & \dots & 0 \\ \vdots & & \vdots \\ 0 & \dots & 0 \end{pmatrix}$$

$\rightarrow D \cdot \vec{r} = 0$ for all strings

$\rightarrow D \neq 0 \Rightarrow D$ has a non-zero element $O_{i_1, p_1, i_2, p_2, \dots, i_k, p_k}$

$$P_r (\text{Algorithm says 'YES'} \mid D \neq 0)$$

WLOG assume that non-zero element is in the top left corner of D .

$$D = \begin{pmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ 0 & & & 0 \\ & & & \\ & & & 0 \end{pmatrix}$$

The argument can be formulated for any position of non-zero element.

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lets calculate the first element of

$$e = D \cdot \vec{v} \quad (\text{if } e \text{ is all zeros ALS says 'YES'})$$

$$e_1 = d_1 \cdot v_1 + d_2 \cdot v_2 + \dots + d_n \cdot v_n \quad (\text{if equal to zero we get wrong answer})$$

$$\underbrace{v_1}_{\neq 0} = \frac{d_2 v_2 + \dots + d_n v_n}{-d_1} \quad (\text{non-zero})$$

R.H.S is a fixed value (principle of delay (deferred decision))

v_1 is chosen from $\{0, 1\}$.

$$\Pr(e_1 > 0 \mid D \neq 0) \leq \frac{1}{2}.$$

Is the choice of $v \in \{0, 1\}^n$ special?

How about $v \in S \subseteq \mathbb{F} \quad |S|=2$

How about $v \in S \subseteq \mathbb{F} \quad |S|=k$

$$\Downarrow$$

$$\Pr \text{ of error} \leq \frac{1}{k}$$

Note that this technique can be used for any matrix

identity $X \stackrel{?}{=} Y$ if X and Y are given explicitly

Example:

$$\begin{pmatrix} 1232 & 4758 \\ 891 & 932 \end{pmatrix} \stackrel{?}{=} \begin{pmatrix} 8 & 7 \\ 14 & 58 \end{pmatrix} \pmod{3} \quad \begin{pmatrix} 0 \\ 1 \end{pmatrix} \pmod{3} \quad (0, 1, 2)$$

Example:

$$\begin{pmatrix} 1233 & 4798 \\ 891 & 932 \end{pmatrix} \stackrel{?}{=} \begin{pmatrix} 8 & 7 \\ 14 & 98 \end{pmatrix} \pmod{3} \quad \begin{matrix} \text{mod } 3 \\ (0, 1) \end{matrix}$$

$$4798 + 932 = 98 + 7 \pmod{3}$$

Polynomials

$P(x) \in \mathbb{F}_p[x]$ (set of all polynomials over \mathbb{F}_p)

$$f(x) = \sum_{i=0}^{\infty} a_i x^i \pmod{p}$$

$\forall a_i \in \mathbb{F}_p$

Is polynomial $P(x)$ identically 0?

Are $P_1(x)$ and $P_2(x)$ equal?

$$P_1(x) - P_2(x) \stackrel{?}{=} 0$$

Verify whether $P_1(x) \cdot P_2(x) = P_3(x)$

$$P_1(x) \cdot P_2(x) - P_3(x) \stackrel{?}{=} 0$$

\rightarrow if $P(x)$ is identically 0, then $\forall a, P(a) = 0$

\rightarrow if $P(x)$ is not identically 0? How many a give $P(a) = 0$?

roots of polynomial(s.)

$P(x)$ has exactly $\deg(P(x)) \rightarrow$ the highest exponent

$$3x^3 + 7x + x + 78x^2 + 3 + 9 + 8 + \dots \pmod{3}$$

||| ?

$$\left. \begin{array}{l} 3x^3 + 7x^2 + \dots \\ 4y^2 + 6y^2 + \dots \end{array} \right\}$$

0

Algorithm. Choose $r \in S \subseteq F$ at random and evaluate $P(r)$. if $P(r) = 0$ say $P(x)$ is identically 0, otherwise

$P(x)$ is not identically 0.

$$\Pr(\text{wrong answer}) \leq \frac{\# \text{ roots}}{|S|} = \frac{\deg(P(x))}{|S|} \leq \frac{n}{|S|} \quad \text{if } \deg(P(x)) = n$$

Similar argument for multivariate polynomials

$$P[x_1, \dots, x_n] \in \mathbb{F}_p[x_1, \dots, x_n]$$

$$P[x_1, \dots, x_n] = \left. \begin{aligned} & c_{0 \dots 0} + c_{10 \dots 0} x_1 + c_{010 \dots 0} x_2 \\ & + c_{200 \dots 0} (x_1^2 x_2) + \dots + c_{a_1 \dots a_n} x_1^{a_1} \dots x_n^{a_n} \end{aligned} \right\}$$

$x_1^2 x_2^3 x_3^7 \rightarrow$ this is a polynomial term

$$\deg(x_1^2 x_2^3 x_3^7) = 7 \quad (\text{sum of all exponents})$$

Total degree of $P(x_1, \dots, x_n)$ = the largest degree over all the terms

Schwartz-Zippel theorem

Let $Q[x_1, \dots, x_n] \in \mathbb{F}[x_1, \dots, x_n]$ of total degree d .

Fix any $S \subseteq \mathbb{F}$ and let r_1, \dots, r_n to be chosen at random

from S_0 . Then:

$$\Pr(Q(v_1, \dots, v_n) = 0 \mid Q(x_1, \dots, x_n) \neq 0) \leq \frac{d}{|S|}$$

Proof by induction in the number of variables

l. B_0 - done above

l. H. - this holds for $n-1$ variables

l. S_0 - show it holds for n variables

Let highest degree of x_n in Q be $k \leq d$

$$Q(x_1, \dots, x_n) = \sum_{i=0}^k x_n^i \cdot Q_i(x_1, \dots, x_{n-1})$$

$$q(x_n) = Q(v_1, \dots, v_{n-1}, x_n)$$

$$\deg(q) = k$$

$$\Pr(q(v_n) = 0 \mid Q_2(v_1, \dots, v_{n-1}) \neq 0) \leq \frac{k}{|S|}$$

$\neq \text{roots}$

From IH

$$\Pr[Q_2(v_1, \dots, v_{n-1}) = 0] \leq \frac{d-2}{|S|}$$

Q_i contains all terms with x_n with power i

$$Q(x_1, x_2) = x_1 x_2 + 3x_1 x_2^2 + 4x_1 x_2^3 + x_1^2 x_2 + 7x_1^2 x_2^2 + 3x_1^2 x_2^3 + x_2 + x_2^3$$

$$= \lambda_1 [x_2 + 3x_2^2 + 4x_2^3] = Q_1(x_2)$$

$$+ \lambda_1^2 [x_2 + 7x_2^2 + 3x_2^3] = Q_2(x_2)$$

$$+ [x_2^2 + x_2^3] = Q_3(x_2)$$

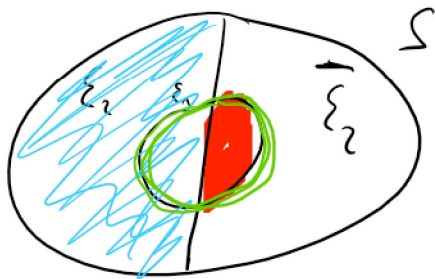
This implies the result.

For two events $\xi_1 = q(v_n) = 0$ $\xi_2 = Q_2(v_1, \dots, v_{n-1}) = 0$

for an $\xi - \xi_1$

for any ξ_1, ξ_2

$$\Pr\{\xi_1\} \leq \Pr\{\xi_1 | \xi_2\} + \Pr\{\xi_2\}$$



Homework

if in $Q[x_1, \dots, x_n]$ $\deg(x_i) = d_i$

and $v_i \in S_i \subseteq F$

Probability that $Q(v_1, \dots, v_n) = 0$ given $Q \neq 0$

is upper bounded by $\frac{d_1}{|S_1|} + \frac{d_2}{|S_2|} + \dots + \frac{d_n}{|S_n|}$

all S_i are identical

$$= \frac{\sum d_i}{|S|} > \frac{d}{|S|}$$

value from Schwartz theorem

$$x_1^2 x_2^1 + x_1^1 x_2^2$$

total degree is 3

$$\text{and } d_1 + d_2 = 4$$

$\boxed{S-2} \Rightarrow$ Freivald's technique

F.f. = decide whether an $n \times n$ matrix $Q = \begin{pmatrix} q_{11} & \dots & q_{1n} \\ \vdots & & \vdots \\ q_{n1} & \dots & q_{nn} \end{pmatrix}$
is identically 0.

$$\begin{aligned} \text{Define } Q[x_1, \dots, x_n] &= Q \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix} \leftarrow \\ &= q_{11}x_1 + q_{12}x_2 + \dots + q_{1n}x_n \\ &\quad + q_{21}x_1 + q_{22}x_2 + \dots + q_{2n}x_n \\ &\quad \vdots \\ &\quad + q_{n1}x_1 + q_{n2}x_2 + \dots + q_{nn}x_n \end{aligned}$$

for Q identically 0 matrix $(\Leftrightarrow) Q[x_1, \dots, x_n]$ is a zero polynomial

Choose $v \in \mathcal{F}^n$ and from S-2 theorem

$$\Pr\{Q[v_1, \dots, v_n] = 0 \mid Q[x_1, \dots, x_n] \neq 0\} \leq \frac{\deg Q}{|S|} = \frac{1}{2}$$