

Essential Information Theory

PA154 Jazykové modelování (1.3)

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March 2, 2021

Source: Introduction to Natural Language Processing (600.465)
Jan Hajič, CS Dept., Johns Hopkins Univ.
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The Formula

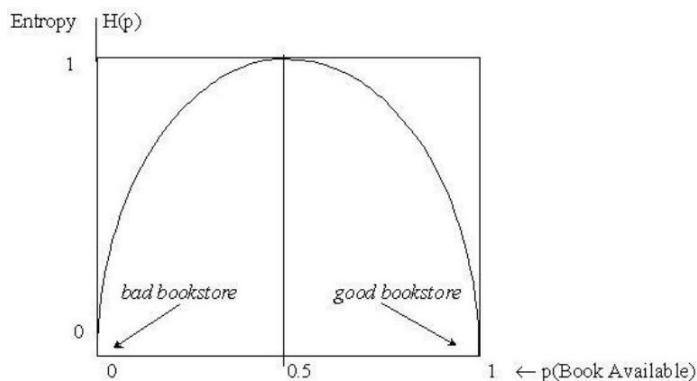
- Let $p_X(x)$ be a distribution of random variable X
- Basic outcomes (alphabet) Ω

Entropy

$$H(X) = -\sum_{x \in \Omega} p(x) \log_2 p(x)$$

- Unit: bits (\log_{10} : nats)
- Notation: $H(X) = H_p(X) = H(p) = H_X(p) = H(p_X)$

Example: Book Availability



The Notion of Entropy

- Entropy – “chaos”, fuzziness, opposite of order, ...
 - you know it
 - it is much easier to create “mess” than to tidy things up...
- Comes from physics:
 - Entropy does not go down unless energy is used
- Measure of **uncertainty**:
 - if low ... low uncertainty

Entropy

The higher the entropy, the higher uncertainty, but the higher “surprise” (information) we can get out of experiment.

Using the Formula: Example

- Toss a fair coin: $\Omega = \{head, tail\}$
 - $p(head) = .5, p(tail) = .5$
 - $H(p) = -0.5 \log_2(0.5) + (-0.5 \log_2(0.5)) = 2 \times ((-0.5) \times (-1)) = 2 \times 0.5 = 1$
- Take fair, 32-sided die: $p(x) = \frac{1}{32}$ for every side x
 - $H(p) = -\sum_{i=1, \dots, 32} p(x_i) \log_2 p(x_i) = -32(p(x_1) \log_2 p(x_1))$
(since for all i $p(x_i) = p(x_1) = \frac{1}{32}$)
 $= -32 \times (\frac{1}{32} \times (-5)) = 5$ (now you see why it's called **bits**?)
- Unfair coin:
 - $p(head) = .2 \dots H(p) = .722$
 - $p(head) = .1 \dots H(p) = .081$

The Limits

- When $H(p) = 0$?
 - if a result of an experiment is **known** ahead of time:
 - necessarily:

$$\exists x \in \Omega; p(x) = 1 \ \forall y \in \Omega; y \neq x \Rightarrow p(y) = 0$$
- Upper bound?
 - none in general
 - for $|\Omega| = n: H(p) \leq \log_2 n$
 - nothing can be more uncertain than the uniform distribution

Entropy and Expectation

- Recall:
 - $E(X) = \sum_{x \in X(\Omega)} p_x(x) \times x$
- Then:

$$E\left(\log_2\left(\frac{1}{p(x)}\right)\right) = \sum_{x \in X(\Omega)} p_x(x) \log_2\left(\frac{1}{p_x(x)}\right) = -\sum_{x \in X(\Omega)} p_x(x) \log_2 p_x(x) = H(p_x) = \textit{notation } H(p)$$

Perplexity: motivation

- Recall:
 - 2 equiprobable outcomes: $H(p) = 1$ bit
 - 32 equiprobable outcomes: $H(p) = 5$ bits
 - 4.3 billion equiprobable outcomes: $H(p) \cong 32$ bits
- What if the outcomes are not equiprobable?
 - 32 outcomes, 2 equiprobable at 0.5, rest impossible:
 - $H(p) = 1$ bit
 - any measure for comparing the entropy (i.e. uncertainty/difficulty of prediction) (also) for random variables with different number of outcomes?

Perplexity

- Perplexity:
 - $G(p) = 2^{H(p)}$
- ... so we are back at 32 (for 32 eqp. outcomes), 2 for fair coins, etc.
- it is easier to imagine:
 - NLP example: vocabulary size of a vocabulary with uniform distribution, which is equally hard to predict
- the "wilder" (biased) distribution, the better:
 - lower entropy, lower perplexity

Joint Entropy and Conditional Entropy

- Two random variables: X (space Ω), Y (Ψ)
- Joint entropy:
 - no big deal: $((X, Y)$ considered a single event):

$$H(X, Y) = -\sum_{x \in \Omega} \sum_{y \in \Psi} p(x, y) \log_2 p(x, y)$$

- Conditional entropy:

$$H(Y|X) = -\sum_{x \in \Omega} \sum_{y \in \Psi} p(x, y) \log_2 p(y|x)$$

recall that $H(X) = E\left(\log_2 \frac{1}{p_x(x)}\right)$
(weighted "average", and weights are not conditional)

Conditional Entropy (Using the Calculus)

- other definition:

$$\begin{aligned} H(Y|X) &= \sum_{x \in \Omega} p(x) H(Y|X=x) = \\ &\quad \text{for } H(Y|X=x), \text{ we can use} \\ &\quad \text{the single-variable definition (x ~ constant)} \\ &= \sum_{x \in \Omega} p(x) \left(-\sum_{y \in \Psi} p(y|x) \log_2 p(y|x)\right) = \\ &= -\sum_{x \in \Omega} \sum_{y \in \Psi} p(y|x) p(x) \log_2 p(y|x) = \\ &= -\sum_{x \in \Omega} \sum_{y \in \Psi} p(x, y) \log_2 p(y|x) \end{aligned}$$

Properties of Entropy I

- Entropy is non-negative:
 - $H(X) \geq 0$
 - proof: (recall: $H(X) = -\sum_{x \in \Omega} p(x) \log_2 p(x)$)
 - $\log_2(p(x))$ is negative or zero for $x \leq 1$,
 - $p(x)$ is non-negative; their product $p(x) \log(p(x))$ is thus negative,
 - sum of negative numbers is negative,
 - and $-f$ is positive for negative f
- Chain rule:
 - $H(X, Y) = H(Y|X) + H(X)$, as well as
 - $H(X, Y) = H(X|Y) + H(Y)$ (since $H(Y, X) = H(X, Y)$)

Properties of Entropy II

- Conditional Entropy is better (than unconditional):
 - ▶ $H(Y|X) \leq H(Y)$
- $H(X, Y) \leq H(X) + H(Y)$ (follows from the previous (in)equalities)
 - ▶ equality iff X, Y independent
 - ▶ (recall: X, Y independent iff $p(X, Y) = p(X)p(Y)$)

- $H(p)$ is concave (remember the book availability graph?)

- ▶ concave function f over an interval (a, b) :

$$\forall x, y \in (a, b), \forall \lambda \in [0, 1]:$$

$$f(\lambda x + (1 - \lambda)y) \geq \lambda f(x) + (1 - \lambda)f(y)$$

- ▶ function f is convex if $-f$ is concave

- for proofs and generalizations, see Cover/Thomas

