

Static Analysis of a Linux Distribution

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How to find programming mistakes efficiently?

0 users (preferably volunteers)



1 Automatic Bug Reporting Tool (ABRT)



2 code review, automated tests, fuzzing

+++git

3 static analysis



Why do we use static analysis at Red Hat?

- ... to find programming mistakes soon enough – example:

```
Error: SHELLCHECK_WARNING:
/etc/rc.d/init.d/squid:136:10: warning: Use "${var:?}" to ensure this never expands to /* .
# 134|         RETVAL=$?
# 135|         if [ $RETVAL -eq 0 ] ; then
# 136|->             rm -rf $SQUID_PIDFILE_DIR/*
# 137|             start
# 138|         else
```

<https://bugzilla.redhat.com/1202858> – *[UNRELEASED] restarting testing build of squid results in deleting all files in hard-drive*

- Static analysis is required for Common Criteria certification.

Agenda

- 1 Code Review, Fuzzing
- 2 Linux Distribution, Reproducible Builds
- 3 Static Analysis of a Linux Distribution
- 4 Dynamic Analysis and Formal Verification

Code Review

- design (anti-)patterns
- error handling (OOM, permission denied, ...)
- validation of input data (headers, length, encoding, ...)
- sensitive data treatment (avoid exposing private keys, ...)
- use of crypto algorithms
- resource management

Fuzzing

- Feeding programs with unusual input.
- Can be combined with valgrind, GCC sanitizers, etc.
- [radamsa](#) – general purpose data fuzzer

```
$ cat file | radamsa | program
```

- [OSS-Fuzz](#) – continuous fuzzing of open source software
 - service provided by Google
 - many security issues detected e.g. in curl

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Linux Distribution

- operating system (OS)
- based on the Linux kernel




- a lot of other programs running in user space



- usually open source

Upstream vs. Downstream

- **Upstream** SW projects – usually independent
- **Downstream** distribution of upstream SW projects
 - Red Hat uses the RPM package manager 
 - Files on the file system owned by **RPM packages**:
 - Dependencies form an oriented graph over packages.
 - We can query package database.
 - We can verify installed packages.

Fedora vs. RHEL

- Fedora 
 - new features available early
 - driven by the community (developers, users, ...)

- RHEL (Red Hat Enterprise Linux) 
 - stability and security of existing deployments
 - driven by Red Hat (and its customers)

Where do RPM packages come from?

- Developers maintain source RPM packages (SRPMs).
- Binary RPMs can be built from SRPMs using `rpmbuild`:

```
rpmbuild --rebuild git-2.30.2-1.fc34.src.rpm
```

- Binary RPMs can be then installed on the system:

```
sudo dnf install git
```

Reproducible Builds

- Local builds are not reproducible.
- `mock` – chroot-based tool for building RPMs:

```
mock -r fedora-rawhide-x86_64 git-2.30.2-1.fc34.src.rpm
```

- `koji` – service for scheduling build tasks

```
koji build rawhide git-2.30.2-1.fc34.src.rpm
```

- Easy to hook static analyzers on the build process!

Agenda

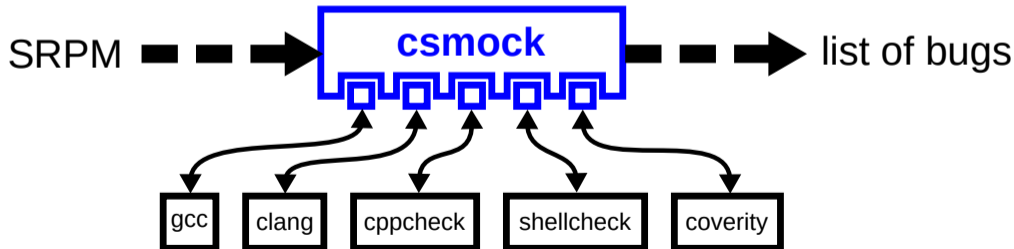
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Static Analysis at Red Hat in Numbers

- Preliminary scan of all RHEL-9 packages in February 2021.
- Analyzed 480 million LoC (Lines of Code) in 3700 packages.
- 98.6 % packages scanned successfully.
- Approx. 680 000 potential bugs detected in total.
- Approx. one potential bug per each 750 LoC.

Analysis of RPM Packages

- Command-line tool to run static analyzers on RPM packages.
- One interface, one output format, plug-in API for (static) analyzers.
- Fully open-source, available in Fedora and CentOS.



csmock – Supported Static Analyzers

	C	C++	C#	Java	Go	JavaScript	PHP	Python	Ruby	Shell
<code>gcc</code>	✓	✓								
<code>gcc -fanalyzer</code>	✓									
<code>clang --analyze</code>	✓	✓								
<code>cppcheck</code>	✓	✓								
<code>coverity</code>	✓	✓	✓	✓	✓	✓	✓	✓	✓	
<code>shellcheck</code>										✓
<code>pylint</code>								✓		
<code>bandit</code>								✓		
<code>smatch</code>	✓									

Need more?

<https://github.com/mre/awesome-static-analysis#user-content-programming-languages-1>

What is important for developers?

The static analyzers need to:

- be fully automatic
- provide reasonable signal to noise ratio
- provide reproducible and consistent results
- be approximately as fast as compilation of the package
- support differential scans:
 - added/fixed bugs in an update?
 - <https://github.com/kdudka/csdiff>



csmock – Output Format

Error: RESOURCE_LEAK (CWE-772):

```
src/fptr.c:450: alloc_fn: Storage is returned from allocation function "calloc".
src/fptr.c:450: var_assign: Assigning: "e" = storage returned from "calloc(24UL, 1UL)".
src/fptr.c:450: overwrite_var: Overwriting "e" in "e = calloc(24UL, 1UL)" leaks the storage that "e" points to.
# 448|         if ((f = (struct opd_fptr *) l->u.refp[i]->ent)->ent == NULL)
# 449|             {
# 450|->         e = calloc (sizeof (struct opd_ent), 1);
# 451|             if (e == NULL)
# 452|                 {
```

Error: CPPCHECK_WARNING (CWE-401):

```
src/fptr.c:464: error[memleak]: Memory leak: e
# 462|     }
# 463|
# 464|-> return ret;
# 465| }
```

Error: RESOURCE_LEAK (CWE-772):

```
src/fptr.c:450: alloc_fn: Storage is returned from allocation function "calloc".
src/fptr.c:450: var_assign: Assigning: "e" = storage returned from "calloc(24UL, 1UL)".
src/fptr.c:464: leaked_storage: Variable "e" going out of scope leaks the storage it points to.
# 462|     }
# 463|
# 464|-> return ret;
# 465| }
```

csmock – Output Format

RESOURCE_LEAK (CWE-772): → checker
 src/fptr.c:450: alloc_fn: Storage is returned from allocation function "calloc".
 src/fptr.c:450: var_assign: Assigning: "e" = storage returned from "calloc(24UL, 1UL)".
 src/fptr.c:450: overwrite_var: Overwriting "e" in "e = calloc(24UL, 1UL)" leaks the storage that "e" points to.
 # 448| if ((f = (struct opd_ptr *) 1->u.refp[i]->ent)->ent == NULL)
 # 449| {
 # 450|-> e = calloc (sizeof (struct opd_ent), 1); → key event
 # 451| if (e == NULL)
 # 452| {
 Error: CPPCHECK_WARNING (CWE-401) → CWE ID
 src/fptr.c:464: error[memLeak]: Memory leak: e
 # 462| }
 # 463|
 # 464|-> return ret; → location info
 # 465| } → other events
 Error: RESOURCE_LEAK (CWE-772):
 src/fptr.c:450: alloc_fn: Storage is returned from allocation function "calloc".
 src/fptr.c:450: var_assign: Assigning: "e" = storage returned from "calloc(24UL, 1UL)".
 src/fptr.c:464: leaked_storage: Variable "e" going out of scope leaks the storage it points to.
 # 462| }
 # 463|
 # 464|-> return ret;
 # 465| } → message associated with the key event



csmock – Output Format (Trace Events)

Error: **RESOURCE LEAK** (CWE-772):

```
src/fptr.c:447: cond_true: Condition "i < 1->nrefs", taking true branch.
src/fptr.c:448: cond_true: Condition "(f = (struct opd_fptr *)1->u.refp[i]->ent)->ent == NULL", taking true branch.
src/fptr.c:450: alloc_fn: Storage is returned from allocation function "calloc".
src/fptr.c:450: var_assign: Assigning: "e" = storage returned from "calloc(24UL, 1UL)".
src/fptr.c:451: cond_false: Condition "e == NULL", taking false branch.
src/fptr.c:456: if_end: End of if statement.
src/fptr.c:462: loop: Jumping back to the beginning of the loop.
src/fptr.c:447: loop_begin: Jumped back to beginning of loop.
src/fptr.c:447: cond_true: Condition "i < 1->nrefs", taking true branch.
src/fptr.c:448: cond_true: Condition "(f = (struct opd_fptr *)1->u.refp[i]->ent)->ent == NULL", taking true branch.
src/fptr.c:450: overwrite_var: Overwriting "e" in "e = calloc(24UL, 1UL)" leaks the storage that "e" points to.
# 448|     if ((f = (struct opd_fptr *) 1->u.refp[i]->ent)->ent == NULL)
# 449|     {
# 450|->         e = calloc (sizeof (struct opd_ent), 1);
# 451|         if (e == NULL)
# 452|         {
```



Example of a Fix

```
--- a/src/fptr.c
+++ b/src/fptr.c
@@ -438,28 +438,29 @@
 GElf_Addr
 opd_size (struct prelink_info *info, GElf_Word entsize)
 {
     struct opd_lib *l = info->ent->opd;
     int i;
     GElf_Addr ret = 0;
     struct opd_ent *e;
     struct opd_fptr *f;

     for (i = 0; i < l->nrefs; ++i)
         if ((f = (struct opd_fptr *) l->u.refp[i]->ent)->ent == NULL)
             {
                 e = calloc (sizeof (struct opd_ent), 1);
                 if (e == NULL)
                     {
                         error (0, ENOMEM, "%s: Could not create OPD table",
                                info->ent->filename);
                         return -1;
                     }

                 e->val = f->val;
                 e->gp = f->gp;
                 e->opd = ret | OPD_ENT_NEW;
+                 f->ent = e;
                 ret += entsize;
             }

     return ret;
 }
```

Upstream vs. Enterprise

Different approaches to static analysis:

- **Upstream**
 - Fix as many bugs as possible.
 - False positive ratio increases over time!
- **Enterprise**
 - Run differential scans to verify code changes.
 - Up to 10% of bugs usually detected as new in an update.
 - Up to 10% of them usually confirmed as real by developers.

Agenda

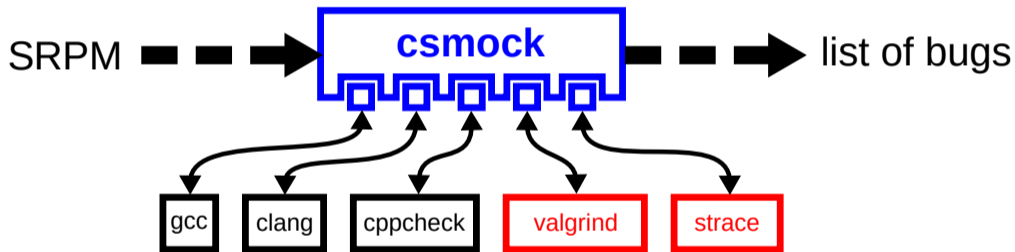
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Dynamic Analysis

- Executes code in a modified run-time environment.
- Embedded in compilers: address sanitizer, thread sanitizer, UB sanitizer, ...
- Standalone tools: valgrind, strace, ...
- Not so easy to automate as static analysis.
- Good to have some test-suite to begin with.

Dynamic Analysis of RPM Packages

- Experimental csmock plug-ins for valgrind and strace:



```
$ sudo yum install csmock-plugin-valgrind
```

```
$ csmock -t valgrind -r fedora-rawhide-x86_64 *.src.rpm
```

Tests Embedded in RPM Packages

```
$ fedpkg clone -a logrotate
$ cd logrotate
$ grep -A8 '%build' logrotate.spec
%build
mkdir build && cd build
%global _configure ../configure
%configure --with-state-file-path=%{_localstatedir}/lib/logrotate/logrotate.status
%make_build

%check
%make_build -C build -s check

$ fedpkg srpm
$ rpmbuild --rebuild *.src.rpm
```

Dynamic Analysis of RPM Packages – Simple Approach

- Dynamic analyzers usually support tracing of child processes.
- Let's combine it together:
 - `valgrind --trace-children=yes rpmbuild --rebuild *.src.rpm`
 - `strace --follow-forks rpmbuild --rebuild *.src.rpm`
- But did we want to dynamically analyze `rpmbuild`, `bash`, `make`, etc.?
 - This makes the analysis extremely slow.
 - We get reports unrelated to `*.src.rpm`.

Dynamic Analysis of RPM Packages – Better Approach

- Produce binaries that will launch a dynamic analyzer for themselves.
- We can use a compiler wrapper to instrument the build of an RPM package:

```
$ export PATH=$(cswrap --print-path-to-wrap):$PATH
$ export CSWRAP_ADD_CFLAGS=-Wl,--dynamic-linker,/usr/bin/csexec-loader
$ export CSEXEC_WRAP_CMD=valgrind
$ rpmbuild --rebuild *.src.rpm
```

- Only binaries produced in `%build` will run through valgrind in `%check`.

Program Interpreter

- Program interpreter specified by [shebang](#):

```
$ head -1 /usr/bin/yum
```

```
#!/usr/bin/python3
```

```
$ /usr/bin/yum [...] → /usr/bin/python3 /usr/bin/yum [...]
```

- Program interpreter specified by ELF header:

```
$ file /sbin/logrotate
```

```
/sbin/logrotate: ELF 64-bit LSB shared object, x86-64, version 1 (SYSV),  
dynamically linked, interpreter /lib64/ld-linux-x86-64.so.2, BuildID[sha1]=...
```

- ELF interpreter can be set to a custom value when linking the binary:

```
$ file ./logrotate
```

```
./logrotate: ELF 64-bit LSB shared object, x86-64, version 1 (SYSV),  
dynamically linked, interpreter /usr/bin/csexec-loader, BuildID[sha1]=...
```

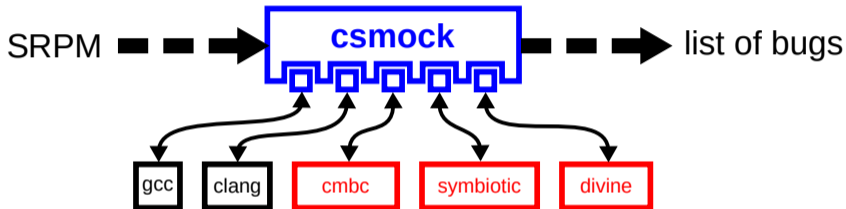
Wrapper of Dynamic Linker – Implementation

- `csexec` works as a wrapper of the system dynamic linker:
<https://github.com/kdudka/cswrap/wiki/csexec>
- `$CSEXEC_WRAP_CMD` can specify a dynamic analyzer to use.
- `csexec` runs the system dynamic linker explicitly (to eliminate self-loop):
`./logrotate [...] → valgrind /lib64/ld-linux-x86-64.so.2 ./logrotate [...]`
- `csexec` uses the `--argv0` option of the system dynamic linker if available:
<https://sourceware.org/git/?p=glibc.git;a=commitdiff;h=c6702789>
- `csexec` emulates the original target of the `/proc/self/exe` symlink.

Wrapper of Dynamic Linker – Evaluation

- No completely unrelated bug reports.
- Minimal performance overhead.
- Minimal interference with commonly used testing frameworks.
- Able to successfully run upstream test-suite of [GNU coreutils](#) (without valgrind).
- Some tests fail if we wrap them by valgrind though:
 - a test that verifies the count [open file descriptors](#)
 - a test that intentionally sets non-existing [\\$TMPDIR](#)
 - ...
- [\[TODO: demo\]](#)

Formal Verification of RPM Packages



- **AUFOVER** (Automation of Formal Verification) project, supported by Technology Agency of the Czech Republic:
<https://starfos.tacr.cz/en/project/TH04010192>
- **SV-COMP** (Competition on Software Verification):
<https://sv-comp.sosy-lab.org/2021/results/results-verified/>



Slides Available Online

<https://kdudka.fedorapeople.org/muni21.pdf>