

## Algebraic techniques (fingerprinting)

- Friedvald's technique for matrix multiplication ✓
- Polynomial comparison: Schwartz-Zippel thm
- SZ thm.  $\Rightarrow$  Friedvald's technique. ✓

## Matrix multiplication

Given  $n \times n$  matrices  $A, B$  and  $C$  over a finite field  $\mathbb{F}_p$ .  
Finite fields are finite sets of numbers with a well defined addition and multiplication. They exist for all prime-power sizes.  $\mathbb{F}_p$  for prime  $p, \mathbb{F}_p = \{0, \dots, p-1\}, +, + \pmod p$ .

Verify whether  $A \cdot B = C$

Naive solution

Multiply  $A \cdot B$  and compare to  $C$   
 $O(n^3)$   $O(n^2)$   
 $[O(n^{2.523})]$

Suppose you want to check whether a new multiplication algorithm is correct, with randomized technique  $A \cdot B \stackrel{?}{=} C$  can be done in  $O(n^2)$ .

1.) Choose  $\vec{v} \in \{0,1\}^n$  at random and calculate

$$\underbrace{(A \cdot (B \cdot \vec{v}))}_{O(n^2)} \quad \text{and} \quad \underbrace{(C \cdot \vec{v})}_{O(n^2)} \quad \text{and compare the results} \\ O(n)$$

$$(A \cdot B - C) \cdot \vec{v} \stackrel{?}{=} \vec{0}$$

2.) If the results are equal alg. outputs 'YES'  $\neq$   
if not alg. outputs 'NO'

3.) output NO  $\Rightarrow A \cdot B \neq C$  w.p. 1  
output YES  $\Rightarrow A \cdot B \neq C$  w.p.  $= \frac{1}{2}$

### ANALYSIS:

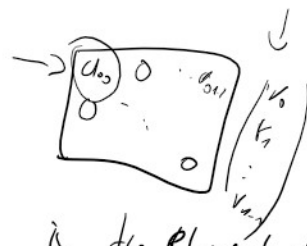
$\rightarrow$  We can reduce the problem to finding whether

$$D = A \cdot B - C \text{ is identically } 0 \quad D = \begin{pmatrix} 0 & \dots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \dots & 0 \end{pmatrix}$$

$\rightarrow D \cdot \vec{v} = \vec{0}$  for all strings  $\vec{v}$ .

$\rightarrow D \neq 0 \Rightarrow D$  has a non-zero element

$\Pr(\text{Algorithm outputs 'YES'} \mid D \neq 0)$



WLOG assume that the non-zero element of  $D$  is the element  $d_{11}$ .

Let's calculate the first element of  $\vec{z} = D \cdot \vec{v}$   $\vec{z} = (z_1, z_2, \dots, z_n)$

$$e_0 = \begin{bmatrix} d \\ d_{00} \end{bmatrix} \cdot r_0 + d_{01} \cdot r_1 + d_{02} \cdot r_2 + \dots + d_{0, n-1} \cdot r_{n-1} \stackrel{?}{=} 0$$

$$r_0 = \frac{d_{01} \cdot r_1 + d_{02} \cdot r_2 + \dots + d_{0, n-1} \cdot r_{n-1}}{-d_{00}}$$

for all  $(r_0, \dots, r_{n-1})$  R.H.S is a fixed value in  $\mathbb{F}_p$   $\{0, \dots, p-1\}$

$r_0$  is chosen from  $\{0, 1\}$   $\leftarrow$

$$\Pr(e_0 = 0 \mid D \neq 0) \leq \frac{1}{2}$$

Is the choice of  $\vec{v} \in \{0, 1\}^n$  special?

How about  $\vec{v} \in S \subseteq \mathbb{F}_p^n$   $|S| = 2$

How about  $\vec{v} \in S \subseteq \mathbb{F}_p^n$   $|S| = \epsilon \Rightarrow \Pr(\text{error}) \leq \frac{1}{\epsilon}$

Note that this technique can be used for any matrix identity

$X = Y$  if  $X$  and  $Y$  are given explicitly

## Polynomials

$P(x) \in \mathbb{F}_p[x]$  (set of all polynomials over  $\mathbb{F}_p$ )

$$f(x) = \sum_{i=0}^{\infty} a_i x^i \pmod{p} \quad \forall_i a_i \in \mathbb{F}_p$$

$$3x^2 + 7x + 78x^2 + \dots$$

• Is polynomial  $p(x)$  identically 0?

$$\begin{aligned} &3x^2 + 7x + 78x^2 + 319 \\ &+ 8 \pmod{3} \end{aligned}$$

• Are  $P_1(x)$  and  $P_2(x)$  equal?

$$P_1(x) - P_2(x) \equiv 0? \quad \checkmark$$

• Verifying whether  $P_1(x) \cdot P_2(x) \stackrel{?}{=} P_3(x)$

$$P_1(x) \cdot P_2(x) - P_3(x) \stackrel{?}{=} 0 \quad \checkmark$$

→ if  $P(x) \equiv 0$ , then  $\forall a \quad P(a) = 0$

→ if  $P(x) \not\equiv 0$  how many  $a$  s.t.  $P(a) = 0$ ?

↓  
roots of polynomial

$P(x)$  has at most  $\deg(P(x))$  distinct roots

↓  
the highest exponent

### Algorithm

Choose  $r \in S \subseteq F$  at random and evaluate  $P(r)$

if  $P_r(0) = 0$  s.t.  $P(x) \equiv 0$ , otherwise  $P(x) \not\equiv 0$ .

$$\Pr(\text{error}) \leq \frac{\#\text{roots}}{|S|} = \frac{\deg(P(x))}{|S|} \leq \frac{n}{|S|} \quad \deg(P(x)) = n$$

Similar argument for multivariate polynomials exists: Schwartz-Zippel theorem

$$P\{x_1, \dots, x_n\} \in \mathbb{F}_p[x_1, \dots, x_n]$$

$$P\{x_1, \dots, x_n\} = \underbrace{c_{0000}}_n + \underbrace{c_{1000}}_n x_1 + \underbrace{c_{0100}}_n x_2 + \dots + \underbrace{c_{0001}}_n x_n \\ + \underbrace{c_{1100}}_n (x_1 x_2) - \dots - \dots - \underbrace{c_{a_1 a_2 \dots a_n}}_n x_1^{a_1} x_2^{a_2} \dots x_n^{a_n}$$

$$c_{a_1 \dots a_n} \in \mathbb{F}_p$$

$x_1^2 x_2^3 x_3 x_7 \rightsquigarrow$  Polynomial terms

$$\deg(x_1^2 x_2^3 x_3 x_7) = 7$$

Total degree  $P(x_1, \dots, x_n) =$  the largest degree over all it's terms

### Schwartz-Zippel thm.

Let  $Q\{x_1, \dots, x_n\} \in \mathbb{F}_p[x_1, \dots, x_n]$  of total degree  $d$ .  
Fix any  $S \subseteq \mathbb{F}_p$  and let  $v_1, \dots, v_n$  be chosen at random from  $S$ .

then:

$$\Pr(Q(v_1, \dots, v_n) = 0 \mid Q(x_1, \dots, x_n) \neq 0) \leq \frac{d}{|S|}$$

Proof by induction in the number of variables

I.B. done above

I.H. this holds for  $n-1$  variables

I.S. show that this holds for  $n$  variables

$$Q(x_1, \dots, x_n) = \sum_{i=0}^k x_n^i Q_i(x_1, \dots, x_{n-1})$$

Principle of deferred decision allows us to choose  $v_1, \dots, v_{n-1}$  before choosing  $v_n$

$$Q(x_1, x_2) = x_1 x_2 + 3x_1 x_2^2 + 4x_1 x_2^3 + x_1^2 x_2 + 7x_1^2 x_2^2 + 3x_1^2 x_2^3 + x_2 + x_2^3$$

$$= x_1 \cdot (x_2 + 3x_2^2 + 4x_2^3) + x_1^2 \cdot (x_2 + 7x_2^2 + 3x_2^3) + (x_2 + x_2^3)$$

$Q_1$   
 $Q_2$   
 $Q_0$

$$q(x_n) = Q(v_1, \dots, v_{n-1}, x_n) = \sum_{i=0}^k x_n^i Q_i(v_1, \dots, v_{n-1})$$

If  $Q \neq 0$  then there is at least one value  $i$ , such that  $Q_i \neq 0$ . Let  $\epsilon$  be largest such  $i$ .

$$\Pr(q(v_n) = 0 \mid Q_\epsilon(v_1, \dots, v_{n-1}) \neq 0, Q \neq 0) < \frac{\epsilon}{|S|}$$

from l.h.

$$\Pr\{Q_\epsilon(v_1, \dots, v_{n-1}) = 0 \mid Q \neq 0\} \leq \frac{d-\epsilon}{|S|}$$

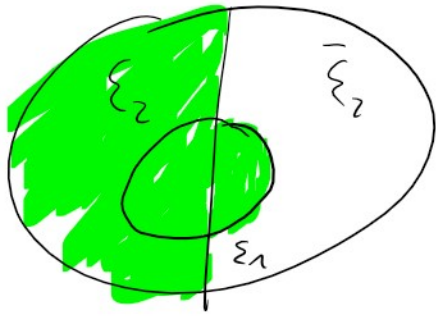
This implies the result.

For two events  $\mathcal{E}_1 = \{q(v_n) = 0 \mid Q \neq 0\}$

$$\mathcal{E}_2 = \{Q_\epsilon(v_1, \dots, v_{n-1}) = 0 \mid Q \neq 0\}$$

$\Downarrow$

$$\Pr\{\mathcal{E}_1\} \leq \Pr\{\mathcal{E}_1 \mid \overline{\mathcal{E}_2}\} + \Pr\{\mathcal{E}_2\}$$



QED

it in  $Q[x_1, \dots, x_n]$   $\deg(x_i) = d_i$

and  $v_i \in S_i \subseteq F$

$$\Pr \{ Q[x_1, \dots, x_n] = 0 \mid Q \neq 0 \} \leq \frac{d_1}{|S_1|} + \frac{d_2}{|S_2|} + \dots + \frac{d_n}{|S_n|}$$

if all  $|S_i|$  are identical  $= \frac{\sum_i d_i}{|S_i|} \geq \frac{d}{|S_i|}$   $\neq$

$\Rightarrow$  Freivald's matrix equality

F.t.  $Q \begin{pmatrix} a_{00} & \dots & a_{0n} \\ \vdots & & \vdots \\ a_{nn} & \dots & a_{nn} \end{pmatrix}$  is identically 0?

$$Q[x_0, \dots, x_{n-1}] \quad Q \begin{pmatrix} x_0 \\ \vdots \\ x_{n-1} \end{pmatrix}$$

$$= a_{0,0}x_0 + a_{0,1}x_1 + \dots + a_{0,n-1}x_{n-1}$$

$$\begin{aligned}
&= a_{0,0}x_0 + a_{0,1}x_1 + \dots + a_{0,n-1}x_{n-1} \\
&+ a_{1,0}x_0 + a_{1,1}x_1 + \dots + a_{1,n-1}x_{n-1} \\
&\quad \vdots \\
&+ a_{n-1,0}x_0 + \dots + a_{n-1,n-1}x_{n-1}
\end{aligned}$$

for  $Q \equiv 0 \Leftrightarrow Q[x_1, \dots, x_n] \equiv 0$

From S-2

$$\Pr [Q[x_1, \dots, x_n] = 0 \mid Q[x_1, \dots, x_n] \neq 0] \leq \frac{\deg Q}{|S|} = \frac{1}{2}$$