

File and disk encryption

Milan Brož xbroz@fi.muni.cz Petr Švenda svenda@fi.muni.cz Faculty of Informatics, Masaryk University



Data storage encryption

Lecture

- File and disk encryption
- Distributed storage encryption
- Abstraction layers, hardware acceleration
- Cryptography basic principles
 - Confidentiality and integrity protection
 - Encryption modes
 - Key management
- Attacks and common issues
- Laboratory





File and disk encryption

MOTIVATION & STORAGE LAYERS OVERVIEW

Motivation

Offline, "Data at Rest" protection notebook, external drives, data in cloud, backups

Key removal = easy data disposal

Confidentiality protection

often **policy** to encrypt mobile devices / sensitive data prevents data leaks (stolen device)

Data integrity protection (not often yet)

Overview

(Distributed) Storage Stack

layers accessing storage through blocks (sectors) distributed => adding network layer

Full Disk Encryption (FDE)

self-encrypted drives, (software) sector-level encryption

Filesystem-level encryption

general-purpose filesystem with encryption cryptographic file systems

Storage stack & encryption layers

Userspace	Application	(Application specific)
OS kernel "Hardware"	Virtual file-system (directories, files,)	File-system encryption
	Specific file-system (NTFS, ext4, XFS, APFS)	
	Volume Management (partitions, on-demand allocation, snapshots, deduplication,)	Disk (sector) encryption
	Block layer (sectors I/O)	
	Storage transport (USB, SCSI, SAS, SATA, FC, NVME)	HW-based encryption self-encrypted drives, inline (slot) encryption, chipset-based encryption, hardware security module
	Device drivers	
	Hardware (I/O controllers, disks, NAND chips,)	

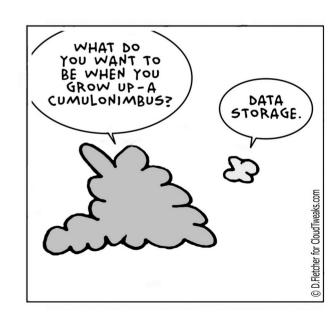
Software Defined Storage (SDS)

- commodity hardware with abstracted storage/network logic
- encryption is "just" one logic function
- usually combination with classic storage (and encryption)
- Distributed storage storage + network layer
 - must use also network layer encryption
 - note differences in network and storage encryption (replay attack resistance, integrity protection, ...)

Distributed Storage, Cloud & Encryption

Distributed storage – add network layer

- Shared volumes (disk encryption below)
- Clustered file-system (fs encryption)
- Distributed object store (object encryption)
- Cloud data storage

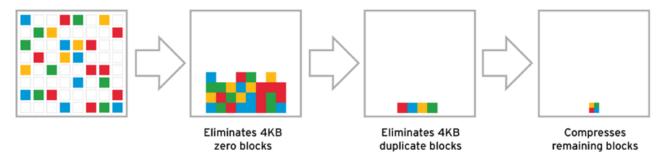




Cloud storage – common features

Deduplication – avoid to store repeated data

VDO data reduction processing



Compression

special case: zeroed blocks

Data snapshots (in time)

COW (copy on write)

Cloud storage & encryption

Encryption with storage backend, network access and compression & deduplication & snapshots ...

Encryption on client side (end-to-end)

- no efficiency for deduplication/compression
- ~ in future homomorphic encryption?

Encryption on server side

- confidentiality for clients is lost
- server has access to plaintext



Full Disk Encryption (FDE)

Block device – transparent disk sector level

- disk, partition, VM disk image
- ciphertext device / virtual plaintext device
- atomic unit is sector (512 bytes, 4k, 64k)
- consecutive sector numbers
- sectors encrypted independently

One key decrypts the whole device

- media (volume) key one per device
- unlocking passphrases/keys

Filesystem-level Encryption

File/Directory

- atomic unit is filesystem block (~ sector in FDE)
- blocks are encrypted independently
- Generic filesystems with encryption
 - some metadata can be kept in plaintext (name, size, ...)
- Cryptographic filesystems
 - metadata encrypted
 - ~ stacked layer over generic filesystem

Multiple keys / multiple users

File vs. disk encryption

Full disk encryption

- + for notebook, external drives (offline protection)
- + no user decision later what to encrypt, transparency
- + hibernation partition and swap encryption
- more users whole disk accessible
- key disclosure complete data leak
- usually no integrity protection
- +/- self-encrypted drives you have to trust hw

Examples: SED (hw), LUKS, VeraCrypt, BitLocker, FileVault

File vs. disk encryption

Filesystem based encryption

- + multiple users
- +/- user can decide what to encrypt
- + copied files keeps encryption in-place
- + more effective (only really used blocks)
- + should provide integrity protection (not always!)
- more complicated sw, usually more bugs
- unusable for swap partitions

Examples: Linux fscrypt API, ZFS, APFS (Apple fs)

Examples of HW-based encryption

- Self-encrypting drives (SED), OPAL2 standard
 - Encryption on the same chip providing media access
- Inline encryption
 - Slots for keys (through OS context)
- Chipset-based encryption
 - Encryption on controller chip (e.g. USB bridge)
- Hardware acceleration
 - AES-NI, accelerators, ASICs, GPUs, ...
- Secure hardware I tokens
 - HSM, TPM, SmartCards, ...



File and disk encryption

CRYPTOGRAPHY

Cryptography algorithms primitives

Symmetric encryption

block ciphers

cipher block mode + initial vector / tweaks

hash algorithms

Key management and key storage

Random Number Generators (RNG)

Key Derivation Functions (KDF)

asymmetric cryptography

Deniable encryption / Steganography

Data confidentiality & integrity

Confidentiality

Data are available only to authorized users.

Integrity

Data are consistent.

Data has not been modified by unauthorized user.

=> all modifications must be detected.

Note: replay attack (revert to old valid data) detection cannot be provided without separate trusted store

Data integrity / authenticated encryption

Poor man's authentication (= no authentication)

- User is able to detect unexpected change
- Very limited, cannot prevent old content replacement

Integrity – additional overhead

- Where to store integrity data?
- Encryption + separate integrity data
- Authenticated modes (combines both)
- Tamper Evident Counter (TEC)
- Merkle tree



File and disk encryption

DATA ENCRYPTION MODES

Symmetric encryption (examples)

AES, Cammelia, Adiantum, Serpent, Twofish, (Specks, Kuznyechik, ...)

Encryption-only modes

- Storage encryption mostly CBC, XTS
- Length-preserving encryption, block tweak

Authenticated modes (encryption + integrity)

Integrity protection often on higher layer.

Standards

IEEE 1619 – encryption modes for storage

NIST Special Publications (SP) –

ciphers,modes, KDF, password handling, ...

OPAL2 – self-encrypted drives

IEEE 1667 – authentication

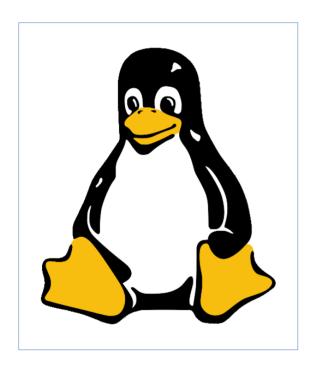
FIPS 140-2, 140-3, Common Criteria (CC)

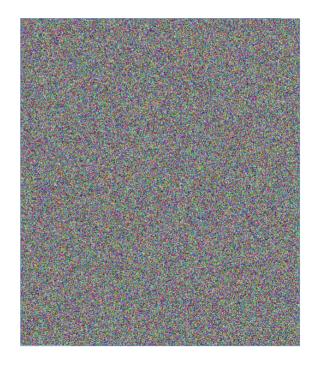
Propagation of plaintext changes

A change in plaintext should transform to randomly-looking change in the whole ciphertext sector. Solutions?

- Ignore it, and decrease granularity of change
 => change location inside ciphertext sector
- Use wide mode (encryption block size = sector size)
 - requires at least 2x encryption loop
 - modes are patent encumbered
- Use additional operations
 - Elephant diffuser in Windows Bitlocker
 - Google Adiantum (cipher composition)

Encryption example output

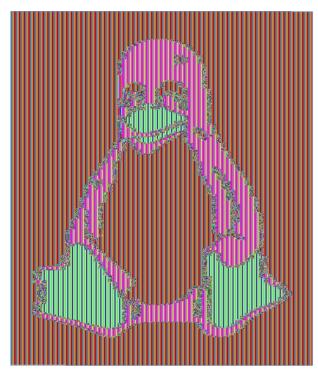




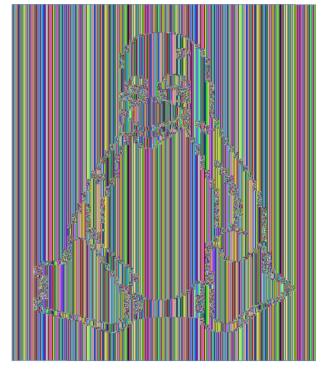
plaintext

ciphertext

Wrongly used encryption – patterns, leaks



ECB mode

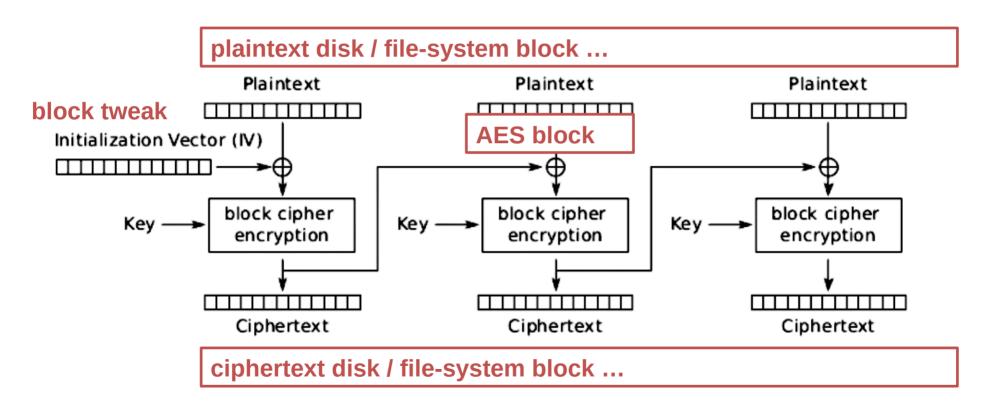


AES-XTS & constant IV

Cipher-Block-Chaining (CBC) mode

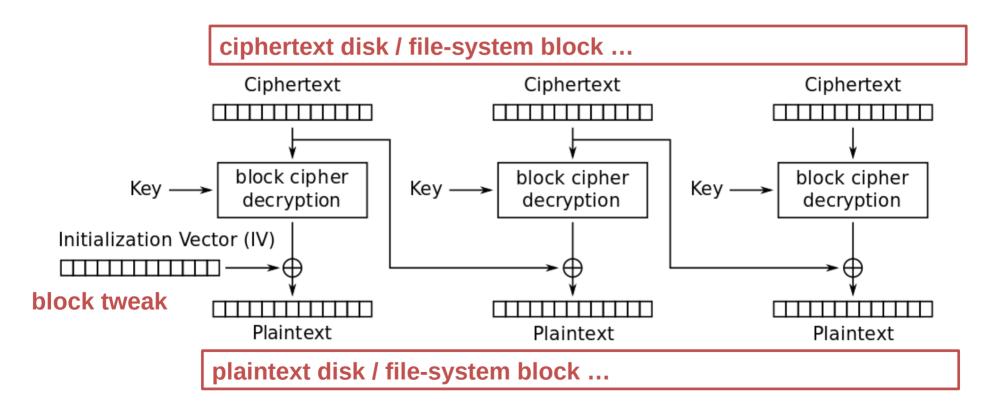
- Blocks cannot be encrypted in parallel
- Blocks can be decrypted in parallel
- Tweak must be non-predictable (watermarking!)

CBC encryption





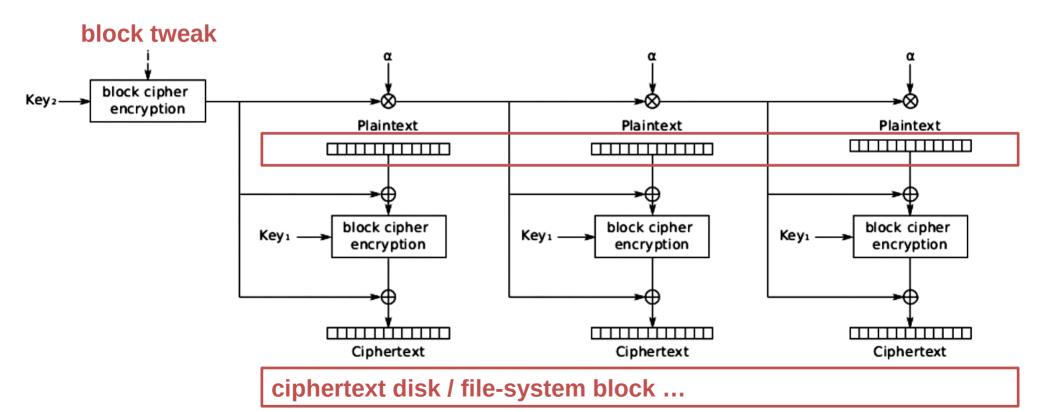
CBC decryption



XOR-Encrypt-XOR (XEX/XTS) mode

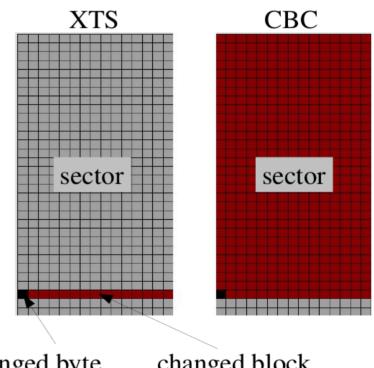
- Encryption/decryption can be run in parallel
- Two keys 512-bit key means AES-256
- Tweak can be predictable nonce sector number (offset)
- Used in most of FDE systems today (2022)
- Not a wide mode plaintext change located in AES block
- Trade-off for performance

XTS mode encryption/decryption



CBC vs XTS change propagation

- XTS is trade-off for performance
- For storage, data always aligned to encryption blocks
 XTS: no ciphertext stealing
- Initial vector/tweak is important
- CBC is phased out today

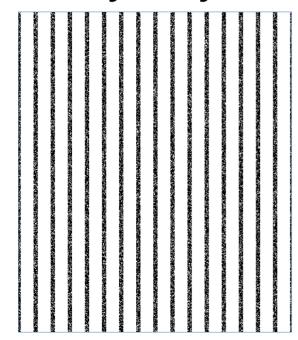


changed byte
(in plaintext)

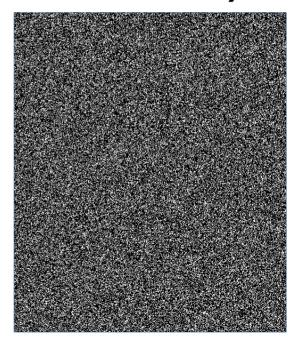
changed block (in ciphertext)

AES-XTS IV mode – sector# vs random

Every 64 byte changed (ciphertext differences)



IV is sector number

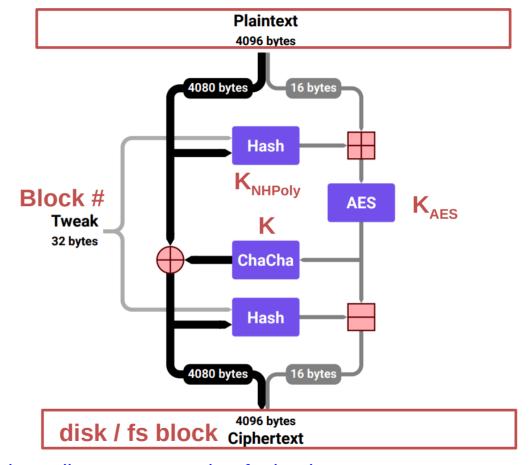


randomized IV

Adiantum

- Low-end mobile device disk / file encryption
- Wide "mode"
- HBSB composition:
 - Hash NHPoly1305)
 - Block Cipher AES
 - Stream Cipher XChaCha12,20
 - Hash NHPoly1305
- Key derivation

 $K_{AES} | K_{NHPoly} = XChaCha(K,1|0..0)$



https://eprint.iacr.org/2018/720

https://security.googleblog.com/2019/02/introducing-adiantum-encryption-for.html

Steganography / deniable encryption

Plausible deniability:

existence of encrypted data is deniable if adversary cannot prove that it exists

Steganography

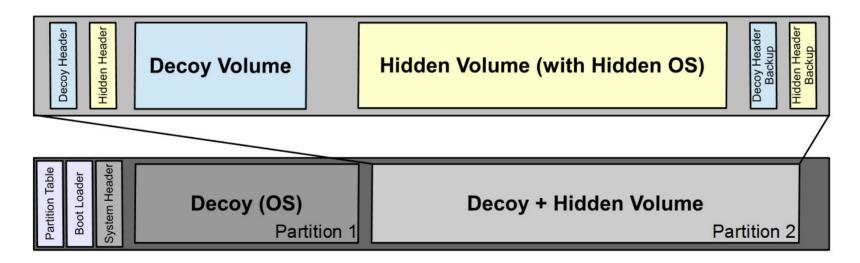
hiding data in another data object

Steganographic file-systems

Deniable disk encryption

Trivial example: VeraCrypt hidden disk

- FAT linear allocation (other fs are very problematic)
- Hide another disk in unallocated space



Deniable encryption problems

Side-channels

tracking activity that cannot be explained for decoy system

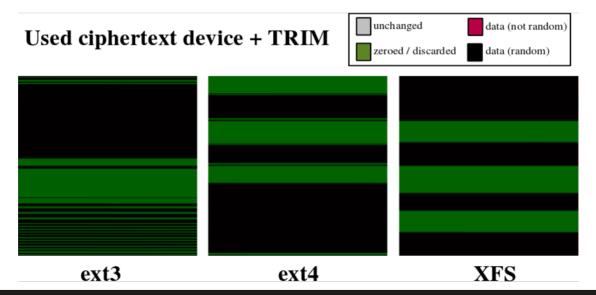
- Software: link to recently open documents, ... Suspicious parameters (FAT), disabled TRIM, ...
- Hardware: internal SSD block allocations (access to "unused" areas)

Incompatibility with new drives (TRIM)

Note: flash storage HW is much more complicated (NAND chips management, wear-leveling, ...). With low-level access you can detect suspicious patterns.

TRIM / discard and encryption

- TRIM informs SSD drive about unused space
- Unused space is detectable
- Pattern recognition (fs type) example





File and disk encryption

KEY MANAGEMENT

Long-term key generation and key store

Encryption key (~ Media Encryption Key – MEK)

- Used to encrypt device
 - change means complete reencryption
- Usually generated by a secure RNG

Unlocking key (~ Key Encryption Key – KEK)

- Key wrap (MEK remains the same)
- Can be derived from passphrase
 - PBKDF2 (Password Based Key Derivation)
 - scrypt, Argon2 (memory-hard KDFs) dictionary and brute-force resistance

Key storage

Outside of encrypted device / filesystem

- Another device, file, token, SmartCard, TPM, HSM
- On a key server (network)
- Protected by another key key wrap, key encapsulation

On the same disk (with encrypted data)

Metadata on-disk – key slots

Integration with key management tools

LDAP, Active Directory, ...

Combination of above

Key removal and recovery

Key removal (wipe of key) = data disposal

- intended (secure disk disposal)
- unintended (error) => complete lost of data

Key recovery

- trade-off between security and user-friendly approach
- metadata backups
- multiple metadata copies
- Key Escrow (key backup to different system)
- recovery key to regenerate encryption key



File and disk encryption

ATTACK EXAMPLES

Attacks always get better, they never get worse.

- Against algorithm design
 - wrongly used encryption mode, IV
- To implementation
 - insufficient entropy (broken RNG)
 - weak derivation from weak passwords
 - side channels
- Obtaining key or passphrase in open form
 - Cold Boot
 - "Black bag analysis" Malware, key-logger
 - social engineering, "Rubber-hose cryptoanalysis"

Integrity attacks

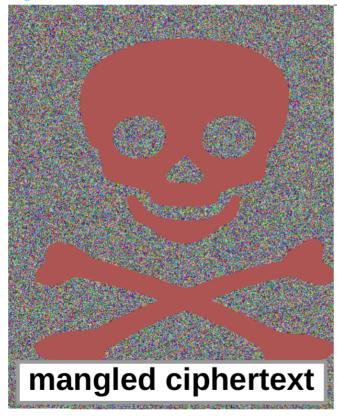
No integrity protection

- Inserted random block
 - => undetected data corruption
- Inserted block from other part of disk
- Random error (RAM bit flip)
 - => "silent data corruption"

Weak integrity protection

Inserted previous content of (ciphertext) block
 replay attack

Integrity attacks





FDE attacks – real-world examples

- Some chipsets use ECB mode
- Weak key derivation (brute-force possible)
- Trivial unlocking mode (1-bit password is ok/bad)
- Weak key-escrow (backup key in EEPROM, ...)
- SED switch power attacks
- SED ransomware and unconfigured passphrase
- Cold boot key in memory
- Key loggers
- Weak RNG (key is not random)
- LUKS2 reencryption (forced decryption)





Laboratory – FDE attack examples

Basic understanding of FDE

VeraCrypt, LUKS

Scanning memory image for encryption key

ColdBoot attack principle

HW key-logger attack

Why you have to trust your HW

Optional: flawed algorithm and watermarking

Revealing legacy TrueCrypt hidden disk existence (CBC)