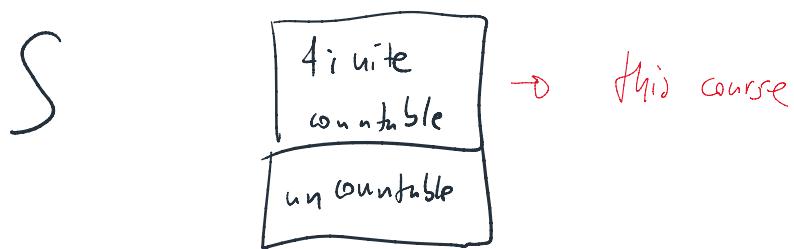


PROBABILITY THEORY - A CRASH COURSE

1.) PROBABILITY SPACE - a set of all possible outcomes of a random experiment



Example: a set of all n-bit strings

2.) Probability events: $E \subseteq S$

Example: a string with exactly 3 symbols '1'.

3.) Probability function $p: S \rightarrow [0,1]$

$$\text{if } i \in S \quad p(i) \geq 0$$

$$\sum_{i \in S} p(i) = 1$$

Example: $\boxed{p(i) = \frac{1}{2^n}}$ (uniform distribution over all strings i of size n)

$$p(E) = \sum_{i \in E} p(i)$$

Example: What's the probability to obtain a 5-bit string with exactly 3 symbols '1'?

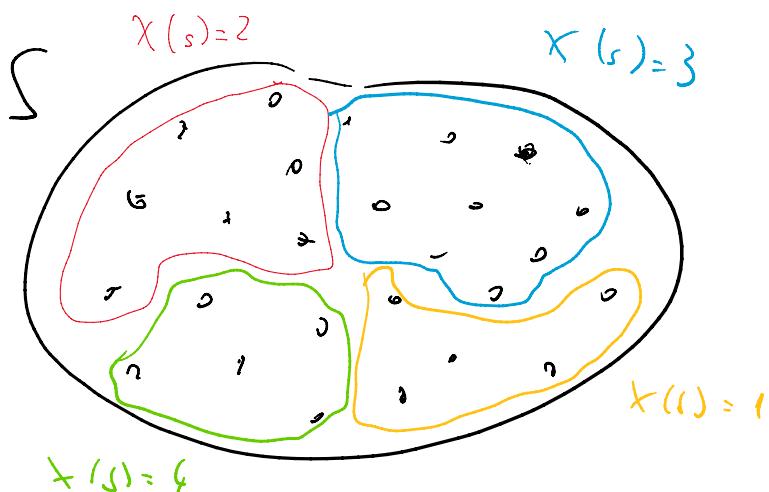
$$P(E) = \sum_{i \in E} P(i) = \sum_{i \in E} \frac{1}{32} = \binom{5}{3} \cdot \frac{1}{32} = \frac{10}{32}$$

RANDOM VARIABLE

(X, Y, Z)

$$X: S \rightarrow \mathbb{R}$$

Essentially X is a partition of S into mutually exclusive and collectively exhaustive set of events.



Example: X is a number of symbols '1' in an n-bit random string

Example: For $n=4$ what is the distribution of X ?

$$\Pr(X=0) = \frac{1}{16}$$

$$\Pr(X=1) = \frac{4}{16}$$

$$\Pr(X>1) = \frac{11}{16}$$

$$\Pr(t=2) = \frac{6}{16}$$

$$\Pr(x \geq 2 \wedge x < 4) = \frac{10}{16}$$

$$\Pr(t=3) = \frac{1}{4}$$

$$\Pr(x=4) = \frac{1}{16}$$

Expectation of R.V.s

$$E(x) = E(X) = \sum_{\substack{i \in R \\ i \in \text{Im}(X)}} i \cdot \Pr(X=i)$$

Example: $E(X) = \sum_{i=0}^4 i \cdot \Pr(t=i) = 0 \cdot \frac{1}{16} + 1 \cdot \frac{4}{16} + 2 \cdot \frac{6}{16} + 3 \cdot \frac{4}{16} + 4 \cdot \frac{1}{16}$

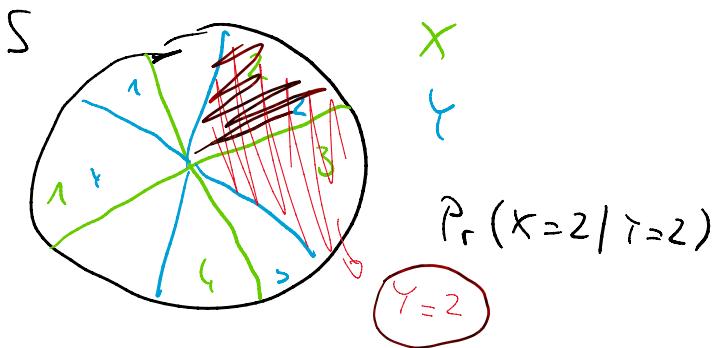
$$= 2$$

Conditional probabilities

Given 2 random variables X and Y over the same random experiment

define conditional probability of X given $Y=j$

$$\Pr(X=i | Y=j) = \Pr(X=i \wedge Y=j) / \Pr(Y=j) \quad [\Pr(Y=j) \neq 0]$$



Intuitively we are creating a new probability space equal to event

$\gamma = j$.

Example:

$$\mathcal{S} = \{0, 1\}^4$$

$X = \text{number of symbols } '1'$

$\gamma = \text{parity of the string}$

$$\Pr(\gamma=0) = 1/2$$

$$\Pr(\gamma=1) = 1/2$$

$$\rightarrow \Pr[X=3 | \gamma=0] = 0 = \Pr[X=3 \wedge \gamma=0]$$

$$\Pr[X=3 | \gamma=1] = \frac{\Pr[X=3, \gamma=1]}{\Pr[\gamma=1]} = \frac{8/16}{1/2} = \frac{8}{16} = \frac{1}{2}$$

Independence of random variables

X and γ are independent if for all $i, j \in \mathbb{R}$

$$\Pr(X=i | \gamma=j) = \Pr(X=i)$$

$x=1$	$x=2$								
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Red	Blue								
Green	Orange								
Red	Blue								
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$x=3$									

$\gamma=1$
 $\gamma=2$
 $\gamma=3$
 $\gamma=4$

Example: Are X and γ from the previous example independent?

$$\Pr(X=3) = 1/4 \neq \Pr(X=3 | \gamma=0) = 0$$

Z - two values of the first bit of $\{0,1\}^3$

Are X and Y independent?

$$\Pr[Z=1] = \frac{1}{2}$$

$$\Pr[Z=0] = \frac{1}{2}$$

$$\Pr[Z=1 | Y=1] = \frac{1}{2}$$

$$\Pr[Z=0 | Y=1] = \frac{1}{2}$$

$$\Pr[Z=1 | Y=0] = \frac{1}{2}$$

$$\Pr[Z=0 | Y=0] = \frac{1}{2}$$

1 0 0 0		0 1 1 1
1 1 1 0		0 1 0 0
1 1 0 1		0 0 1 0
1 0 1 1		0 0 0 1

1 1 1 1	
1 0 0 1	
1 0 1 0	
1 1 0 0	

LINEARITY OF EXPECTATION

$$W = X + Y + Z$$

$$E(W) = E(X+Y+Z) = E(X) + E(Y) + E(Z)$$

Example

$$W = X+Y+Z$$

$$\begin{aligned} E(W) &= \sum_{i,j,\xi} (i+j+\xi) \cdot \Pr(X=i, Y=j, Z=\xi) \\ &= \underset{\substack{\text{=} \\ \text{Z}}}{E(X)} + \underset{\substack{\text{=} \\ \text{Y}}}{E(Y)} + \underset{\substack{\text{=} \\ \text{Z}}}{E(Z)} = 3 \end{aligned}$$

$$E(X_1 \cdot X_2) \neq E(X_1) E(X_2)$$

The law of total probability

r.v. X and Y

$$\begin{aligned} \Pr(X=i) &= \sum_{j \in \text{Im}(Y)} \Pr(X=i | Y=j) \cdot \Pr(Y=j) \\ &= \sum_{j \in \text{Im}(Y)} \Pr(X=i \wedge Y=j) \end{aligned}$$

RQUICKSORT

IN: A collection of numbers S

OUT: Ordered list of numbers in S

1.) if S contains a single element output S

2.) choose a pivot $z \in S$ uniformly at random

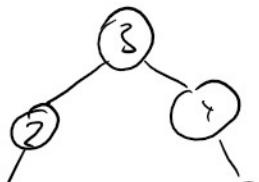
3.) Create S_1 which contains all $s \in S$ $s \leq z$

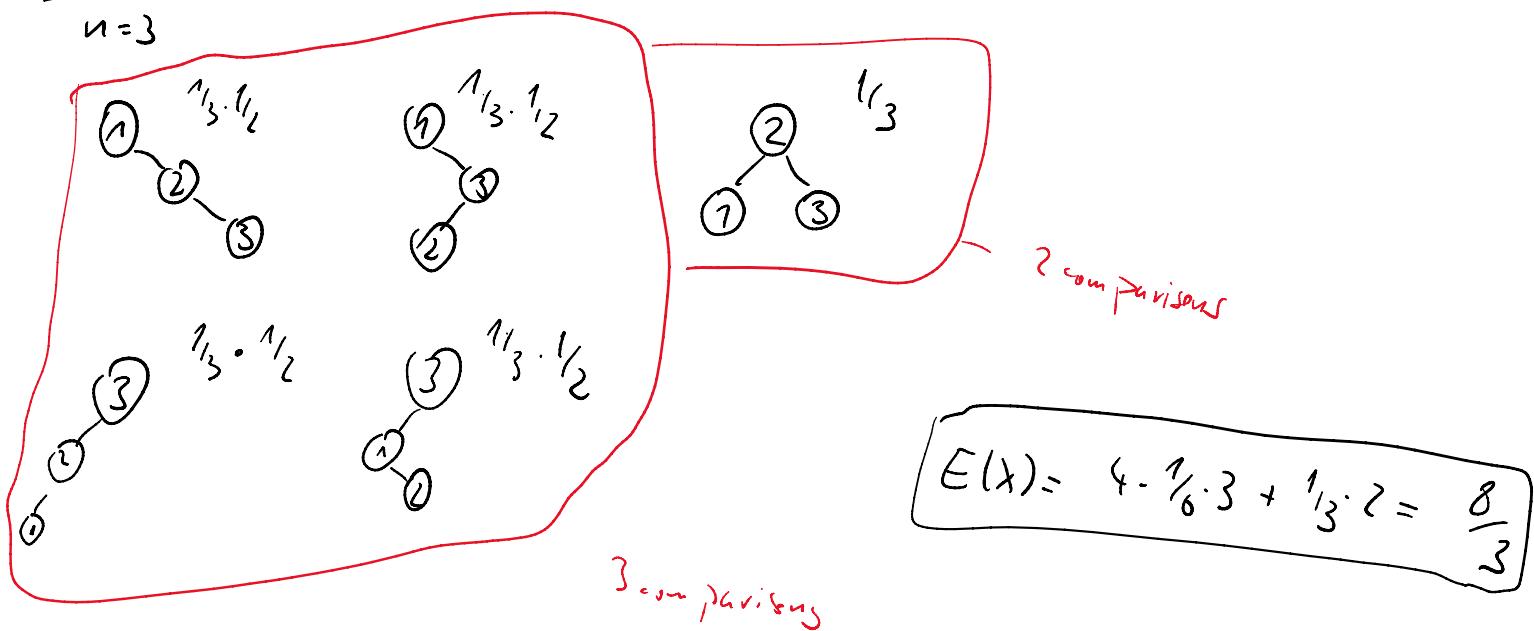
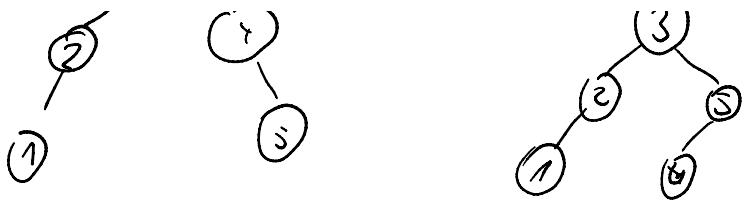
Create S_2 which contains all $s \in S$ $s > z$

4.) Output $(\text{quicksort}(S_1), z, \text{quicksort}(S_2))$

$$S = \{1, 2, 3, 4, 5\}$$

Each run can be identified with an ordered binary tree





X_i := number of comparisons with input of size i .

$$f(i) \geq E(X_i)$$

Can we find an upper bound $f(i)$ on $E(X_i)$ that shows $E(X_i)$ scales well?

$$f(i) \in O(i \log i)$$