

# IAo85: Satisfiability and Automated Reasoning

## Seminar 5

**Exercise 1** If you didn't write the bounded model checking exercise using incremental SMT calls, rewrite it as such. Finish the k-induction exercise from the last seminar.

**Exercise 2** You know the problem of Satisfiability Modulo Theories (SMT), where the task is to compute any model of the given formula. There is a related problem of Optimization Modulo Theories (OMT), where the task is to compute for a given formula  $\varphi$  and a term  $t$  a model that minimizes the value of  $t$ .

For example, given a formula

$$\varphi = (x = 0 \wedge y \geq 10) \vee (x = 2 \wedge y \geq 5)$$

over LIA and a term  $t = x + y$ , the solution of optimization modulo theories problem would be a model  $\mu(x = 2), \mu(y = 5)$ .

Design an algorithm that solves the OMT problem using potentially multiple queries to an SMT solver and implement it in pySMT.

**Exercise 3** Use the congruence closure algorithm to determine which of the following conjunctions of UF-literals are satisfiable. Also identify some of the implied equalities and disequalities.

1.  $x = y \wedge f(x) \neq f(y)$ ,
2.  $f(x) \neq x \wedge f^2(x) = x \wedge f^4(x) = x$ ,
3.  $f(x) \neq x \wedge f^3(x) = x \wedge f^5(x) = x$ .

**Exercise 4** Use the algorithm from the lecture to determine which of the following conjunctions of literals in difference logic over integers are satisfiable. Also identify some of the implied inequalities.

1.  $x - y = 5 \wedge z - y \leq 3 \wedge x - z < 4$ ,
2.  $a - b \leq 3 \wedge c - b \leq 10 \wedge d - a \leq 1 \wedge a - d \leq 5 \wedge c - a \leq 1 \wedge d - c \geq 3 \wedge d - b \geq 2$ .