

## Disk failures in the real world: What does an MTTF of 1,000,000 hours mean to you?

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### Abstract

Component failure in large-scale IT installations is becoming an ever larger problem as the number of components in a single cluster approaches a million.

In this paper, we present and analyze field-gathered disk replacement data from a number of large production systems, including high-performance computing sites and internet services sites. About 100,000 disks are covered by this data, some for an entire lifetime of five years. The data include drives with SCSI and FC, as well as SATA interfaces. The mean time to failure (MTTF) of those drives, as specified in their datasheets, ranges from 1,000,000 to 1,500,000 hours, suggesting a nominal annual failure rate of at most 0.88%.

We find that in the field, annual disk replacement rates typically exceed 1%, with 2-4% common and up to 13% observed on some systems. This suggests that field replacement is a fairly different process than one might predict based on datasheet MTTF.

We also find evidence, based on records of disk replacements in the field, that failure rate is not constant with age, and that, rather than a significant infant mortality effect, we see a significant early onset of wear-out degradation. That is, replacement rates in our data grew constantly with age, an effect often assumed not to set in until after a nominal lifetime of 5 years.

Interestingly, we observe little difference in replacement rates between SCSI, FC and SATA drives, potentially an indication that disk-independent factors, such as operating conditions, affect replacement rates more than component specific factors. On the other hand, we see only one instance of a customer rejecting an entire population of disks as a bad batch, in this case because of media error rates, and this instance involved SATA disks.

Time between replacement, a proxy for time between failure, is not well modeled by an exponential distribution and exhibits significant levels of correlation, including autocorrelation and long-range dependence.

### 1 Motivation

Despite major efforts, both in industry and in academia, high reliability remains a major challenge in running large-scale IT systems, and disaster prevention and cost of actual disasters make up a large fraction of the total cost of ownership. With ever larger server clusters, maintaining high levels of reliability and availability is a growing problem for many sites, including high-performance computing systems and internet service providers. A particularly big concern is the reliability of storage systems, for several reasons. First, failure of storage can not only cause temporary data unavailability, but in the worst case it can lead to permanent data loss. Second, technology trends and market forces may combine to make storage system failures occur more frequently in the future [24]. Finally, the size of storage systems in modern, large-scale IT installations has grown to an unprecedented scale with thousands of storage devices, making component failures the norm rather than the exception [7].

Large-scale IT systems, therefore, need better system design and management to cope with more frequent failures. One might expect increasing levels of redundancy designed for specific failure modes [3, 7], for example. Such designs and management systems are based on very simple models of component failure and repair processes [22]. Better knowledge about the statistical properties of storage failure processes, such as the distribution of time between failures, may empower researchers and designers to develop new, more reliable and available storage systems.

Unfortunately, many aspects of disk failures in real systems are not well understood, probably because the owners of such systems are reluctant to release failure data or do not gather such data. As a result, practitioners usually rely on vendor specified parameters, such as mean-time-to-failure (MTTF), to model failure processes, although many are skeptical of the accuracy of

those models [4, 5, 33]. Too much academic and corporate research is based on anecdotes and back of the envelope calculations, rather than empirical data [28].

The work in this paper is part of a broader research agenda with the long-term goal of providing a better understanding of failures in IT systems by collecting, analyzing and making publicly available a diverse set of real failure histories from large-scale production systems. In our pursuit, we have spoken to a number of large production sites and were able to convince several of them to provide failure data from some of their systems.

In this paper, we provide an analysis of seven data sets we have collected, with a focus on storage-related failures. The data sets come from a number of large-scale production systems, including high-performance computing sites and large internet services sites, and consist primarily of hardware replacement logs. The data sets vary in duration from one month to five years and cover in total a population of more than 100,000 drives from at least four different vendors. Disks covered by this data include drives with SCSI and FC interfaces, commonly represented as the most reliable types of disk drives, as well as drives with SATA interfaces, common in desktop and nearline systems. Although 100,000 drives is a very large sample relative to previously published studies, it is small compared to the estimated 35 million enterprise drives, and 300 million total drives built in 2006 [1]. Phenomena such as bad batches caused by fabrication line changes may require much larger data sets to fully characterize.

We analyze three different aspects of the data. We begin in Section 3 by asking how disk replacement frequencies compare to replacement frequencies of other hardware components. In Section 4, we provide a quantitative analysis of disk replacement rates observed in the field and compare our observations with common predictors and models used by vendors. In Section 5, we analyze the statistical properties of disk replacement rates. We study correlations between disk replacements and identify the key properties of the empirical distribution of time between replacements, and compare our results to common models and assumptions. Section 6 provides an overview of related work and Section 7 concludes.

## 2 Methodology

### 2.1 What is a disk failure?

While it is often assumed that disk failures follow a simple fail-stop model (where disks either work perfectly or fail absolutely and in an easily detectable manner [22, 24]), disk failures are much more complex in reality. For example, disk drives can experience latent sector faults or transient performance problems. Often it

is hard to correctly attribute the root cause of a problem to a particular hardware component.

Our work is based on hardware replacement records and logs, i.e. we focus on disk conditions that lead a drive customer to treat a disk as permanently failed and to replace it. We analyze records from a number of large production systems, which contain a record for every disk that was replaced in the system during the time of the data collection. To interpret the results of our work correctly it is crucial to understand the process of how this data was created. After a disk drive is identified as the likely culprit in a problem, the operations staff (or the computer system itself) perform a series of tests on the drive to assess its behavior. If the behavior qualifies as faulty according to the customer’s definition, the disk is replaced and a corresponding entry is made in the hardware replacement log.

The important thing to note is that there is not one unique definition for when a drive is faulty. In particular, customers and vendors might use different definitions. For example, a common way for a customer to test a drive is to read all of its sectors to see if any reads experience problems, and decide that it is faulty if any one operation takes longer than a certain threshold. The outcome of such a test will depend on how the thresholds are chosen. Many sites follow a “better safe than sorry” mentality, and use even more rigorous testing. As a result, it cannot be ruled out that a customer may declare a disk faulty, while its manufacturer sees it as healthy. This also means that the definition of “faulty” that a drive customer uses does not necessarily fit the definition that a drive manufacturer uses to make drive reliability projections. In fact, a disk vendor has reported that for 43% of all disks returned by customers they find no problem with the disk [1].

It is also important to note that the failure behavior of a drive depends on the operating conditions, and not only on component level factors. For example, failure rates are affected by environmental factors, such as temperature and humidity, data center handling procedures, workloads and “duty cycles” or powered-on hours patterns.

We would also like to point out that the failure behavior of disk drives, even if they are of the same model, can differ, since disks are manufactured using processes and parts that may change. These changes, such as a change in a drive’s firmware or a hardware component or even the assembly line on which a drive was manufactured, can change the failure behavior of a drive. This effect is often called the effect of batches or vintage. A bad batch can lead to unusually high drive failure rates or unusually high rates of media errors. For example, in the HPC3 data set (Table 1) the customer had 11,000 SATA drives replaced in Oct. 2006 after observing a high fre-

Data set	Type of cluster	Duration	#Disk events	# Servers	Disk Count	Disk Parameters	MTTF (Mhours)	Date of first Deploym.	ARR (%)
HPC1	HPC	08/01 - 05/06	474	765	2,318	18GB 10K SCSI	1.2	08/01	4.0
			124	64	1,088	36GB 10K SCSI	1.2		2.2
HPC2	HPC	01/04 - 07/06	14	256	520	36GB 10K SCSI	1.2	12/01	1.1
HPC3	HPC	12/05 - 11/06	103	1,532	3,064	146GB 15K SCSI	1.5	08/05	3.7
			4	N/A	144	73GB 15K SCSI	1.5		3.0
			253	N/A	11,000	250GB 7.2K SATA	1.0		3.3
HPC4	Various HPC clusters	09/03 - 08/06	269	N/A	8,430	250GB SATA	1.0	09/03	2.2
			7	N/A	2,030	500GB SATA	1.0	11/05	0.5
			9	N/A	3,158	400GB SATA	1.0	09/05	0.8
COM1	Int. serv.	May 2006	84	N/A	26,734	10K SCSI	1.0	2001	2.8
COM2	Int. serv.	09/04 - 04/06	506	9,232	39,039	15K SCSI	1.2	2004	3.1
COM3	Int. serv.	01/05 - 12/05	2	N/A	56	10K FC	1.2	N/A	3.6
			132	N/A	2,450	10K FC	1.2	N/A	5.4
			108	N/A	796	10K FC	1.2	N/A	13.6
			104	N/A	432	10K FC	1.2	1998	24.1

Table 1: Overview of the seven failure data sets. Note that the disk count given in the table is the number of drives in the system at the end of the data collection period. For some systems the number of drives changed during the data collection period, and we account for that in our analysis. The disk parameters 10K and 15K refer to the rotation speed in revolutions per minute; drives not labeled 10K or 15K probably have a rotation speed of 7200 rpm.

quency of media errors during writes. Although it took a year to resolve, the customer and vendor agreed that these drives did not meet warranty conditions. The cause was attributed to the breakdown of a lubricant leading to unacceptably high head flying heights. In the data, the replacements of these drives are not recorded as failures.

In our analysis we do not further study the effect of batches. We report on the field experience, in terms of disk replacement rates, of a set of drive customers. Customers usually do not have the information necessary to determine which of the drives they are using come from the same or different batches. Since our data spans a large number of drives (more than 100,000) and comes from a diverse set of customers and systems, we assume it also covers a diverse set of vendors, models and batches. We therefore deem it unlikely that our results are significantly skewed by “bad batches”. However, we caution the reader not to assume all drives behave identically.

## 2.2 Specifying disk reliability and failure frequency

Drive manufacturers specify the reliability of their products in terms of two related metrics: the *annualized failure rate (AFR)*, which is the percentage of disk drives in a population that fail in a test scaled to a per year estimation; and the *mean time to failure (MTTF)*. The AFR of a new product is typically estimated based on accelerated life and stress tests or based on field data from earlier products [2]. The MTTF is estimated as the number of power on hours per year divided by the AFR. A

common assumption for drives in servers is that they are powered on 100% of the time. Our data set providers all believe that their disks are powered on and in use at all times. The MTTFs specified for today’s highest quality disks range from 1,000,000 hours to 1,500,000 hours, corresponding to AFRs of 0.58% to 0.88%. The AFR and MTTF estimates of the manufacturer are included in a drive’s datasheet and we refer to them in the remainder as the *datasheet AFR* and the *datasheet MTTF*.

In contrast, in our data analysis we will report the *annual replacement rate (ARR)* to reflect the fact that, strictly speaking, disk replacements that are reported in the customer logs do not necessarily equal disk failures (as explained in Section 2.1).

## 2.3 Data sources

Table 1 provides an overview of the seven data sets used in this study. Data sets HPC1, HPC2 and HPC3 were collected in three large cluster systems at three different organizations using supercomputers. Data set HPC4 was collected on dozens of independently managed HPC sites, including supercomputing sites as well as commercial HPC sites. Data sets COM1, COM2, and COM3 were collected in at least three different cluster systems at a large internet service provider with many distributed and separately managed sites. In all cases, our data reports on only a portion of the computing systems run by each organization, as decided and selected by our sources.

It is important to note that for some systems the number of drives in the system changed significantly during

the data collection period. While the table provides only the disk count at the end of the data collection period, our analysis in the remainder of the paper accounts for the actual date of these changes in the number of drives. Second, some logs also record events other than replacements, hence the number of disk events given in the table is not necessarily equal to the number of replacements or failures. The ARR values for the data sets can therefore not be directly computed from Table 1.

Below we describe each data set and the environment it comes from in more detail.

**HPC1** is a five year log of hardware replacements collected from a 765 node high-performance computing cluster. Each of the 765 nodes is a 4-way SMP with 4 GB of memory and **three to four 18GB 10K rpm SCSI drives**. Of these nodes, 64 are used as filesystem nodes containing, in addition to the three to four 18GB drives, 17 36GB 10K rpm SCSI drives. The applications running on this system are typically large-scale scientific simulations or visualization applications. The data contains, for each hardware replacement that was recorded during the five year lifetime of this system, when the problem started, which node and which hardware component was affected, and a brief description of the corrective action.

HPC2 is a record of disk replacements observed on the compute nodes of a 256 node HPC cluster. Each node is a 4-way SMP with 16 GB of memory and contains two 36GB 10K rpm SCSI drives, except for eight of the nodes, which contain eight 36GB 10K rpm SCSI drives each. The applications running on this system are typically large-scale scientific simulations or visualization applications. For each disk replacement, the data set records the number of the affected node, the start time of the problem, and the slot number of the replaced drive.

**HPC3** is a record of disk replacements observed on a 1,532 node HPC cluster. Each node is equipped with eight CPUs and 32GB of memory. **Each node**, except for four login nodes, **has two 146GB 15K rpm SCSI disks**. In addition, 11,000 7200 rpm 250GB SATA drives are used in an external shared filesystem and 144 73GB 15K rpm SCSI drives are used for the filesystem metadata. The applications running on this system are typically large-scale scientific simulations or visualization applications. For each disk replacement, the data set records the day of the replacement.

The HPC4 data set is a warranty service log of disk replacements. It covers three types of SATA drives used in dozens of separately managed HPC clusters. For the first type of drive, the data spans three years, for the other two types it spans a little less than a year. The data records, for each of the 13,618 drives, when it was first shipped and when (if ever) it was replaced in the field.

COM1 is a log of hardware failures recorded by an internet service provider and drawing from multiple dis-

tributed sites. Each record in the data contains a timestamp of when the failure was repaired, information on the failure symptoms, and a list of steps that were taken to diagnose and repair the problem. The data does not contain information on when each failure actually happened, only when repair took place. The data covers a population of 26,734 10K rpm SCSI disk drives. The total number of servers in the monitored sites is not known.

COM2 is a warranty service log of hardware failures recorded on behalf of an internet service provider aggregating events in multiple distributed sites. Each failure record contains a repair code (e.g. “Replace hard drive”) and the time when the repair was finished. Again there is no information on the start time of each failure. The log does not contain entries for failures of disks that were replaced in the customer site by hot-swapping in a spare disk, since the data was created by the warranty processing, which does not participate in on-site hot-swap replacements. To account for the missing disk replacements we obtained numbers for the periodic replenishments of on-site spare disks from the internet service provider. The size of the underlying system changed significantly during the measurement period, starting with 420 servers in 2004 and ending with 9,232 servers in 2006. We obtained quarterly hardware purchase records covering this time period to estimate the size of the disk population in our ARR analysis.

The COM3 data set comes from a large external storage system used by an internet service provider and comprises four populations of different types of FC disks (see Table 1). While this data was gathered in 2005, the system has some legacy components that were as old as from 1998 and were known to have been physically moved after initial installation. We did not include these “obsolete” disk replacements in our analysis. COM3 differs from the other data sets in that it provides only aggregate statistics of disk failures, rather than individual records for each failure. The data contains the counts of disks that failed and were replaced in 2005 for each of the four disk populations.

## 2.4 Statistical methods

We characterize an empirical distribution using two import metrics: the mean and the squared coefficient of variation ( $C^2$ ). The squared coefficient of variation is a measure of the variability of a distribution and is defined as the squared standard deviation divided by the squared mean. The advantage of using the squared coefficient of variation as a measure of variability, rather than the variance or the standard deviation, is that it is normalized by the mean, and so allows comparison of variability across distributions with different means.

We also consider the empirical cumulative distribu-

tion function (CDF) and how well it is fit by four probability distributions commonly used in reliability theory: the exponential distribution; the Weibull distribution; the gamma distribution; and the lognormal distribution. We parameterize the distributions through maximum likelihood estimation and evaluate the goodness of fit by visual inspection, the negative log-likelihood and the chi-square tests.

We will also discuss the hazard rate of the distribution of time between replacements. In general, the hazard rate of a random variable  $t$  with probability distribution  $f(t)$  and cumulative distribution function  $F(t)$  is defined as [25]

$$h(t) = \frac{f(t)}{1 - F(t)}$$

Intuitively, if the random variable  $t$  denotes the time between failures, the hazard rate  $h(t)$  describes the instantaneous failure rate as a function of the time since the most recently observed failure. An important property of  $t$ 's distribution is whether its hazard rate is constant (which is the case for an exponential distribution) or increasing or decreasing. A constant hazard rate implies that the probability of failure at a given point in time does not depend on how long it has been since the most recent failure. An increasing hazard rate means that the probability of a failure increases, if the time since the last failure has been long. A decreasing hazard rate means that the probability of a failure decreases, if the time since the last failure has been long.

The hazard rate is often studied for the distribution of lifetimes. It is important to note that we will focus on the hazard rate of the *time between disk replacements*, and not the hazard rate of disk lifetime distributions.

Since we are interested in correlations between disk failures we need a measure for the degree of correlation. The autocorrelation function (ACF) measures the correlation of a random variable with itself at different time lags  $l$ . The ACF, for example, can be used to determine whether the number of failures in one day is correlated with the number of failures observed  $l$  days later. The autocorrelation coefficient can range between 1 (high positive correlation) and -1 (high negative correlation). A value of zero would indicate no correlation, supporting independence of failures per day.

Another aspect of the failure process that we will study is long-range dependence. Long-range dependence measures the memory of a process, in particular how quickly the autocorrelation coefficient decays with growing lags. The strength of the long-range dependence is quantified by the Hurst exponent. A series exhibits long-range dependence if the Hurst exponent,  $H$ , is  $0.5 < H < 1$ . We use the Selfis tool [14] to obtain estimates of the Hurst parameter using five different methods: the absolute value method, the variance method, the R/S method,

HPC1	
Component	%
CPU	44
Memory	29
<b>Hard drive</b>	<b>16</b>
PCI motherboard	9
Power supply	2

Table 2: *Node outages that were attributed to hardware problems broken down by the responsible hardware component. This includes all outages, not only those that required replacement of a hardware component.*

the periodogram method, and the Whittle estimator. A brief introduction to long-range dependence and a description of the Hurst parameter estimators is provided in [15].

### 3 Comparing disk replacement frequency with that of other hardware components

The reliability of a system depends on all its components, and not just the hard drive(s). A natural question is therefore what the relative frequency of drive failures is, compared to that of other types of hardware failures. To answer this question we consult data sets HPC1, COM1, and COM2, since these data sets contain records for all types of hardware replacements, not only disk replacements. Table 3 shows, for each data set, a list of the ten most frequently replaced hardware components and the fraction of replacements made up by each component. We observe that while the actual fraction of disk replacements varies across the data sets (ranging from 20% to 50%), it makes up a significant fraction in all three cases. In the HPC1 and COM2 data sets, disk drives are the most commonly replaced hardware component accounting for 30% and 50% of all hardware replacements, respectively. In the COM1 data set, disks are a close runner-up accounting for nearly 20% of all hardware replacements.

While Table 3 suggests that disks are among the most commonly replaced hardware components, it does not necessarily imply that disks are less reliable or have a shorter lifespan than other hardware components. The number of disks in the systems might simply be much larger than that of other hardware components. In order to compare the reliability of different hardware components, we need to normalize the number of component replacements by the component's population size.

Unfortunately, we do not have, for any of the systems, exact population counts of all hardware components. However, we do have enough information in HPC1 to estimate counts of the four most frequently replaced hard-

HPC1		COM1		COM2	
Component	%	Component	%	Component	%
<b>Hard drive</b>	<b>30.6</b>	Power supply	34.8	<b>Hard drive</b>	<b>49.1</b>
Memory	28.5	Memory	20.1	Motherboard	23.4
Misc/Unk	14.4	<b>Hard drive</b>	<b>18.1</b>	Power supply	10.1
CPU	12.4	Case	11.4	RAID card	4.1
PCI motherboard	4.9	Fan	8.0	Memory	3.4
Controller	2.9	CPU	2.0	SCSI cable	2.2
QSW	1.7	SCSI Board	0.6	Fan	2.2
Power supply	1.6	NIC Card	1.2	CPU	2.2
MLB	1.0	LV Power Board	0.6	CD-ROM	0.6
SCSI BP	0.3	CPU heatsink	0.6	Raid Controller	0.6

Table 3: *Relative frequency of hardware component replacements for the ten most frequently replaced components in systems HPC1, COM1 and COM2, respectively. Abbreviations are taken directly from service data and are not known to have identical definitions across data sets.*

ware components (CPU, memory, disks, motherboards). We estimate that there is a total of 3,060 CPUs, 3,060 memory dimms, and 765 motherboards, compared to a disk population of 3,406. Combining these numbers with the data in Table 3, we conclude that for the HPC1 system, the rate at which in five years of use a memory dimm was replaced is roughly comparable to that of a hard drive replacement; a CPU was about 2.5 times less often replaced than a hard drive; and a motherboard was 50% less often replaced than a hard drive.

The above discussion covers only failures that required a hardware component to be replaced. When running a large system one is often interested in any hardware failure that causes a node outage, not only those that necessitate a hardware replacement. We therefore obtained the HPC1 troubleshooting records for any node outage that was attributed to a hardware problem, including problems that required hardware replacements as well as problems that were fixed in some other way. Table 2 gives a breakdown of all records in the troubleshooting data, broken down by the hardware component that was identified as the root cause. We observe that 16% of all outage records pertain to disk drives (compared to 30% in Table 3), making it the third most common root cause reported in the data. The two most commonly reported outage root causes are CPU and memory, with 44% and 29%, respectively.

For a complete picture, we also need to take the severity of an anomalous event into account. A closer look at the HPC1 troubleshooting data reveals that a large number of the problems attributed to CPU and memory failures were triggered by parity errors, i.e. the number of errors is too large for the embedded error correcting code to correct them. In those cases, a simple reboot will bring the affected node back up. On the other hand, the majority of the problems that were attributed to hard

disks (around 90%) lead to a drive replacement, which is a more expensive and time-consuming repair action.

Ideally, we would like to compare the frequency of hardware problems that we report above with the frequency of other types of problems, such software failures, network problems, etc. Unfortunately, we do not have this type of information for the systems in Table 1. However, in recent work [27] we have analyzed failure data covering any type of node outage, including those caused by hardware, software, network problems, environmental problems, or operator mistakes. The data was collected over a period of 9 years on more than 20 HPC clusters and contains detailed root cause information. We found that, for most HPC systems in this data, more than 50% of all outages are attributed to hardware problems and around 20% of all outages are attributed to software problems. Consistently with the data in Table 2, the two most common hardware components to cause a node outage are memory and CPU. The data of this recent study [27] is not used in this paper because it does not contain information about storage replacements.

## 4 Disk replacement rates

### 4.1 Disk replacements and MTTF

In the following, we study how field experience with disk replacements compares to datasheet specifications of disk reliability. Figure 1 shows the datasheet AFRs (horizontal solid and dashed line), the observed ARR for each of the seven data sets and the weighted average ARR for all disks less than five years old (dotted line). For HPC1, HPC3, HPC4 and COM3, which cover different types of disks, the graph contains several bars, one for each type of disk, in the left-to-right order of the corresponding top-to-bottom entries in Table 1. Since at this

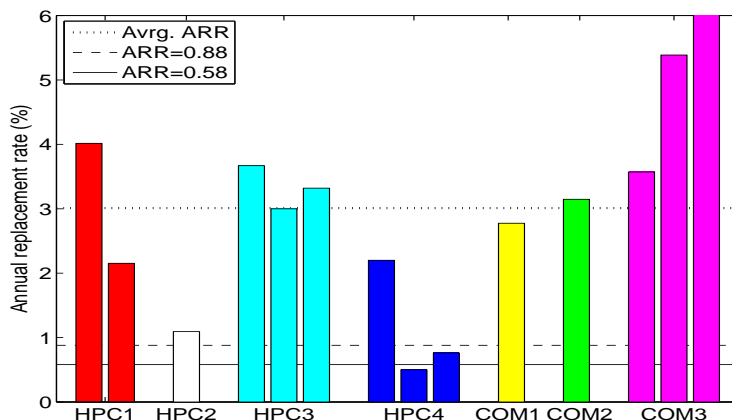


Figure 1: Comparison of datasheet AFRs (solid and dashed line in the graph) and ARRs observed in the field. Each bar in the graph corresponds to one row in Table 1. The dotted line represents the weighted average over all data sets. Only disks within the nominal lifetime of five years are included, i.e. there is no bar for the COM3 drives that were deployed in 1998. The third bar for COM3 in the graph is cut off – its ARR is 13.5%.

point we are not interested in wearout effects after the end of a disk’s nominal lifetime, we have included in Figure 1 only data for drives within their nominal lifetime of five years. In particular, we do not include a bar for the fourth type of drives in COM3 (see Table 1), which were deployed in 1998 and were more than seven years old at the end of the data collection. These possibly “obsolete” disks experienced an ARR, during the measurement period, of 24%. Since these drives are well outside the vendor’s nominal lifetime for disks, it is not surprising that the disks might be wearing out. All other drives were within their nominal lifetime and are included in the figure.

Figure 1 shows a significant discrepancy between the observed ARR and the datasheet AFR for all data sets. While the datasheet AFRs are between 0.58% and 0.88%, the observed ARRs range from 0.5% to as high as 13.5%. That is, the observed ARRs by data set and type, are by up to a factor of 15 higher than datasheet AFRs.

Most commonly, the observed ARR values are in the 3% range. For example, the data for HPC1, which covers almost exactly the entire nominal lifetime of five years exhibits an ARR of 3.4% (significantly higher than the datasheet AFR of 0.88%). The average ARR over all data sets (weighted by the number of drives in each data set) is 3.01%. Even after removing all COM3 data, which exhibits the highest ARRs, the average ARR was still 2.86%, 3.3 times higher than 0.88%.

It is interesting to observe that for these data sets there is no significant discrepancy between replacement rates for SCSI and FC drives, commonly represented as the most reliable types of disk drives, and SATA drives, frequently described as lower quality. For example, the

ARRs of drives in the HPC4 data set, which are exclusively SATA drives, are among the lowest of all data sets. Moreover, the HPC3 data set includes both SCSI and SATA drives (as part of the same system in the same operating environment) and they have nearly identical replacement rates. Of course, these HPC3 SATA drives were decommissioned because of media error rates attributed to lubricant breakdown (recall Section 2.1), our only evidence of a bad batch, so perhaps more data is needed to better understand the impact of batches in overall quality.

It is also interesting to observe that the only drives that have an observed ARR below the datasheet AFR are the second and third type of drives in data set HPC4. One possible reason might be that these are relatively new drives, all less than one year old (recall Table 1). Also, these ARRs are based on only 16 replacements, perhaps too little data to draw a definitive conclusion.

A natural question arises: why are the observed disk replacement rates so much higher in the field data than the datasheet MTTF would suggest, even for drives in the first years of operation. As discussed in Sections 2.1 and 2.2, there are multiple possible reasons.

First, customers and vendors might not always agree on the definition of when a drive is “faulty”. The fact that a disk was replaced implies that it failed some (possibly customer specific) health test. When a health test is conservative, it might lead to replacing a drive that the vendor tests would find to be healthy. Note, however, that even if we scale down the ARRs in Figure 1 to 57% of their actual values, to estimate the fraction of drives returned to the manufacturer that fail the latter’s health test [1], the resulting AFR estimates are still more than a factor of two higher than datasheet AFRs in most cases.



Second, datasheet MTTFs are typically determined based on accelerated (stress) tests, which make certain assumptions about the operating conditions under which the disks will be used (e.g. that the temperature will always stay below some threshold), the workloads and “duty cycles” or powered-on hours patterns, and that certain data center handling procedures are followed. In practice, operating conditions might not always be as ideal as assumed in the tests used to determine datasheet MTTFs. A more detailed discussion of factors that can contribute to a gap between expected and measured drive reliability is given by Elerath and Shah [6].

Below we summarize the key observations of this section.

**Observation 1:** Variance between datasheet MTTF and disk replacement rates in the field was larger than we expected. The weighted average ARR was 3.4 times larger than 0.88%, corresponding to a datasheet MTTF of 1,000,000 hours.

**Observation 2:** For older systems (5-8 years of age), data sheet MTTFs underestimated replacement rates by as much as a factor of 30.

**Observation 3:** Even during the first few years of a system’s lifetime (< 3 years), when wear-out is not expected to be a significant factor, the difference between datasheet MTTF and observed time to disk replacement was as large as a factor of 6.

**Observation 4:** In our data sets, the replacement rates of SATA disks are not worse than the replacement rates of SCSI or FC disks. This may indicate that disk-independent factors, such as operating conditions, usage and environmental factors, affect replacement rates more than component specific factors. However, the only evidence we have of a bad batch of disks was found in a collection of SATA disks experiencing high media error rates. We have too little data on bad batches to estimate the relative frequency of bad batches by type of disk, although there is plenty of anecdotal evidence that bad batches are not unique to SATA disks.

## 4.2 Age-dependent replacement rates

One aspect of disk failures that single-value metrics such as MTTF and AFR cannot capture is that in real life failure rates are not constant [5]. Failure rates of hardware products typically follow a “bathtub curve” with high failure rates at the beginning (infant mortality) and the end (wear-out) of the lifecycle. Figure 2 shows the failure rate pattern that is expected for the life cycle of hard drives [4, 5, 33]. According to this model, the first year

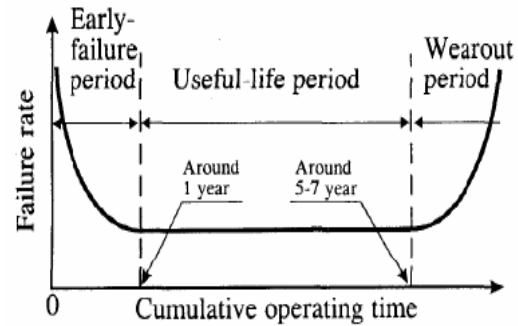


Figure 2: Lifecycle failure pattern for hard drives [33].

of operation is characterized by early failures (or infant mortality). In years 2-5, the failure rates are approximately in steady state, and then, after years 5-7, wear-out starts to kick in.

The common concern, that MTTFs do not capture infant mortality, has led the International Disk drive Equipment and Materials Association (IDEMA) to propose a new standard for specifying disk drive reliability, based on the failure model depicted in Figure 2 [5, 33]. The new standard requests that vendors provide four different MTTF estimates, one for the first 1-3 months of operation, one for months 4-6, one for months 7-12, and one for months 13-60.

The goal of this section is to study, based on our field replacement data, how disk replacement rates in large-scale installations vary over a system’s life cycle. Note that we only see customer visible replacement. Any infant mortality failure caught in the manufacturing, system integration or installation testing are probably not recorded in production replacement logs.

The best data sets to study replacement rates across the system life cycle are HPC1 and the first type of drives of HPC4. The reason is that these data sets span a long enough time period (5 and 3 years, respectively) and each cover a reasonably homogeneous hard drive population, allowing us to focus on the effect of age.

We study the change in replacement rates as a function of age at two different time granularities, on a per-month and a per-year basis, to make it easier to detect both short term and long term trends. Figure 3 shows the annual replacement rates for the disks in the compute nodes of system HPC1 (left), the file system nodes of system HPC1 (middle) and the first type of HPC4 drives (right), at a yearly granularity.

We make two interesting observations. First, replacement rates in all years, except for year 1, are larger than the datasheet MTTF would suggest. For example, in HPC1’s second year, replacement rates are 20% larger than expected for the file system nodes, and a factor of



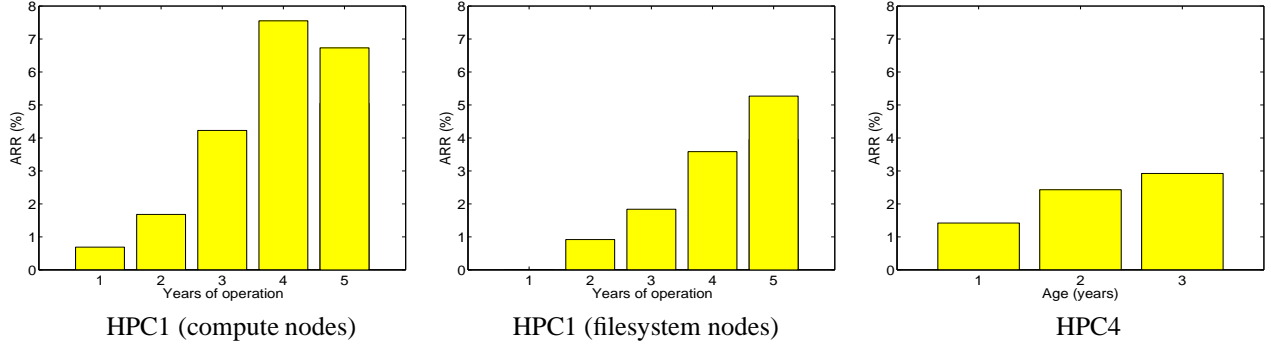


Figure 3: ARR for the first five years of system HPC1’s lifetime, for the compute nodes (left) and the file system nodes (middle). ARR for the first type of drives in HPC4 as a function of drive age in years (right).

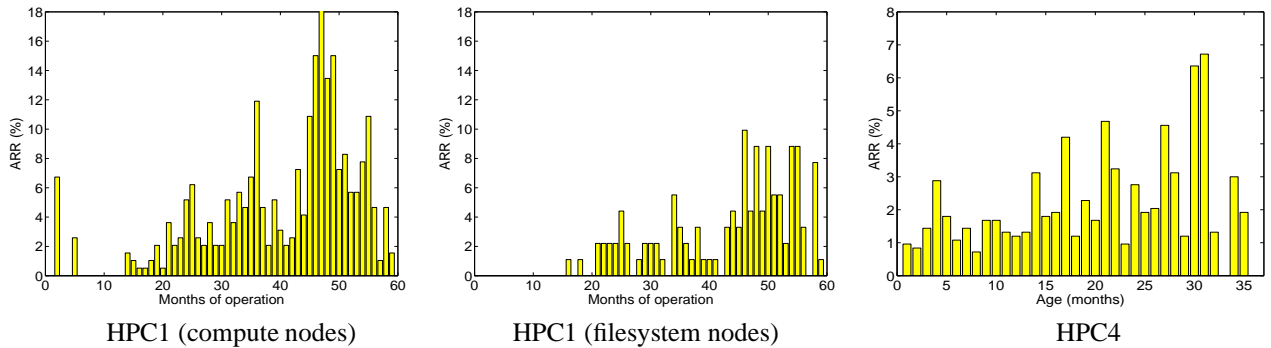


Figure 4: ARR per month over the first five years of system HPC1’s lifetime, for the compute nodes (left) and the file system nodes (middle). ARR for the first type of drives in HPC4 as a function of drive age in months (right).

two larger than expected for the compute nodes. In year 4 and year 5 (which are still within the nominal lifetime of these disks), the actual replacement rates are 7–10 times higher than the failure rates we expected based on datasheet MTTF.

The second observation is that replacement rates are rising significantly over the years, even during early years in the lifecycle. Replacement rates in HPC1 nearly double from year 1 to 2, or from year 2 to 3. This observation suggests that wear-out may start much earlier than expected, leading to steadily increasing replacement rates during most of a system’s useful life. This is an interesting observation because it does not agree with the common assumption that after the first year of operation, failure rates reach a steady state for a few years, forming the “bottom of the bathtub”.

Next, we move to the per-month view of replacement rates, shown in Figure 4. We observe that for the HPC1 file system nodes there are no replacements during the first 12 months of operation, i.e. there’s no detectable infant mortality. For HPC4, the ARR of drives is not higher in the first few months of the first year than the last few months of the first year. In the case of the HPC1 compute nodes, infant mortality is limited to the

first month of operation and is not above the steady state estimate of the datasheet MTTF. Looking at the lifecycle after month 12, we again see continuously rising replacement rates, instead of the expected “bottom of the bathtub”.

Below we summarize the key observations of this section.

**Observation 5:** Contrary to common and proposed models, hard drive replacement rates do not enter steady state after the first year of operation. Instead replacement rates seem to steadily increase over time.

**Observation 6:** Early onset of wear-out seems to have a much stronger impact on lifecycle replacement rates than infant mortality, as experienced by end customers, even when considering only the first three or five years of a system’s lifetime. We therefore recommend that wear-out be incorporated into new standards for disk drive reliability. The new standard suggested by IDEMA does not take wear-out into account [5, 33].

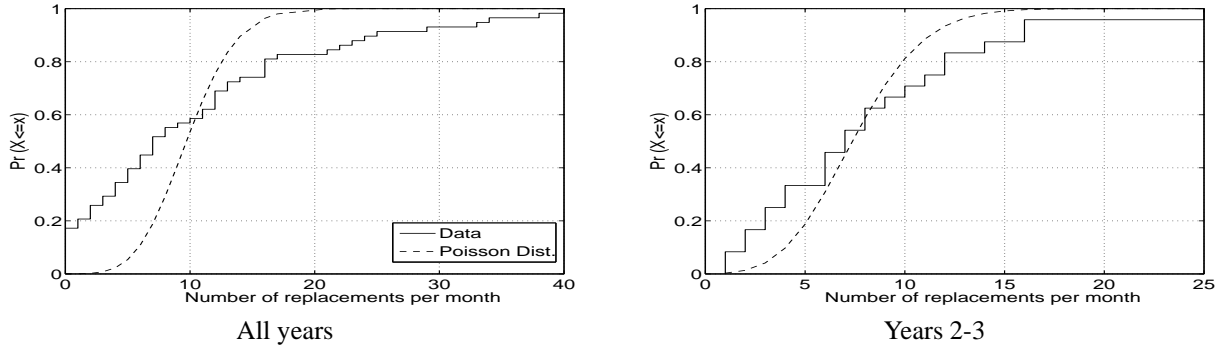


Figure 5: CDF of number of disk replacements per month in HPC1

## 5 Statistical properties of disk failures

In the previous sections, we have focused on aggregate statistics, e.g. the average number of disk replacements in a time period. Often one wants more information on the statistical properties of the time between failures than just the mean. For example, determining the expected time to failure for a RAID system requires an estimate on the probability of experiencing a second disk failure in a short period, that is while reconstructing lost data from redundant data. This probability depends on the underlying probability distribution and maybe poorly estimated by scaling an annual failure rate down to a few hours.

The most common assumption about the statistical characteristics of disk failures is that they form a Poisson process, which implies two key properties:

1. Failures are independent.
2. The time between failures follows an exponential distribution.

The goal of this section is to evaluate how realistic the above assumptions are. We begin by providing statistical evidence that disk failures in the real world are unlikely to follow a Poisson process. We then examine each of the two key properties (independent failures and exponential time between failures) independently and characterize in detail how and where the Poisson assumption breaks. In our study, we focus on the HPC1 data set, since this is the only data set that contains precise timestamps for when a problem was detected (rather than just timestamps for when repair took place).

### 5.1 The Poisson assumption

The Poisson assumption implies that the number of failures during a given time interval (e.g. a week or a month) is distributed according to the Poisson distribution. Figure 5 (left) shows the empirical CDF of the number of

disk replacements observed per month in the HPC1 data set, together with the Poisson distribution fit to the data’s observed mean.

We find that the Poisson distribution does not provide a good visual fit for the number of disk replacements per month in the data, in particular for very small and very large numbers of replacements in a month. For example, under the Poisson distribution the probability of seeing  $\geq 20$  failures in a given month is less than 0.0024, yet we see 20 or more disk replacements in nearly 20% of all months in HPC1’s lifetime. Similarly, the probability of seeing zero or one failure in a given month is only 0.0003 under the Poisson distribution, yet in 20% of all months in HPC1’s lifetime we observe zero or one disk replacement.

A chi-square test reveals that we can reject the hypothesis that the number of disk replacements per month follows a Poisson distribution at the 0.05 significance level. All above results are similar when looking at the distribution of number of disk replacements per day or per week, rather than per month.

One reason for the poor fit of the Poisson distribution might be that failure rates are not steady over the lifetime of HPC1. We therefore repeat the same process for only part of HPC1’s lifetime. Figure 5 (right) shows the distribution of disk replacements per month, using only data from years 2 and 3 of HPC1. The Poisson distribution achieves a better fit for this time period and the chi-square test cannot reject the Poisson hypothesis at a significance level of 0.05. Note, however, that this does not necessarily mean that the failure process during years 2 and 3 does follow a Poisson process, since this would also require the two key properties of a Poisson process (independent failures and exponential time between failures) to hold. We study these two properties in detail in the next two sections.

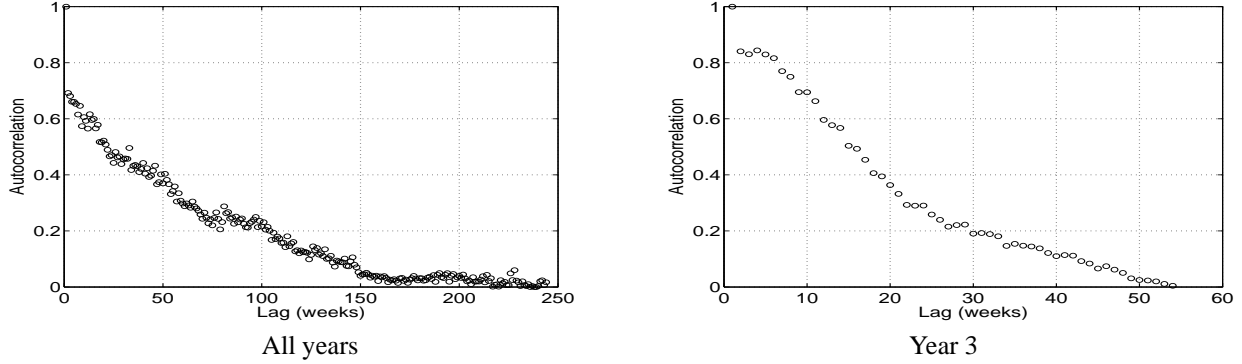


Figure 6: Autocorrelation function for the number of disk replacements per week computed across the entire lifetime of the HPC1 system (left) and computed across only one year of HPC1’s operation (right).

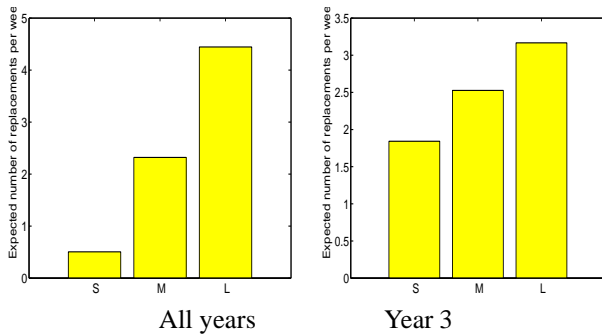


Figure 7: Expected number of disk replacements in a week depending on the number of disk replacements in the previous week.

## 5.2 Correlations

In this section, we focus on the first key property of a Poisson process, the independence of failures. Intuitively, it is clear that in practice failures of disks in the same system are never completely independent. The failure probability of disks depends for example on many factors, such as environmental factors, like temperature, that are shared by all disks in the system. When the temperature in a machine room is far outside nominal values, all disks in the room experience a higher than normal probability of failure. The goal of this section is to statistically quantify and characterize the correlation between disk replacements.

We start with a simple test in which we determine the correlation of the number of disk replacements observed in successive weeks or months by computing the correlation coefficient between the number of replacements in a given week or month and the previous week or month. For data coming from a Poisson processes we would expect correlation coefficients to be close to 0. Instead we find significant levels of correlations, both at the monthly and the weekly level.

The correlation coefficient between consecutive weeks is 0.72, and the correlation coefficient between consecutive months is 0.79. Repeating the same test using only the data of one year at a time, we still find significant levels of correlation with correlation coefficients of 0.4-0.8.

Statistically, the above correlation coefficients indicate a strong correlation, but it would be nice to have a more intuitive interpretation of this result. One way of thinking of the correlation of failures is that the failure rate in one time interval is predictive of the failure rate in the following time interval. To test the strength of this prediction, we assign each week in HPC1’s life to one of three buckets, depending on the number of disk replacements observed during that week, creating a bucket for weeks with small, medium, and large number of replacements, respectively<sup>1</sup>. The expectation is that a week that follows a week with a “small” number of disk replacements is more likely to see a small number of replacements, than a week that follows a week with a “large” number of replacements. However, if failures are independent, the number of replacements in a week will not depend on the number in a prior week.

Figure 7 (left) shows the expected number of disk replacements in a week of HPC1’s lifetime as a function of which bucket the preceding week falls in. We observe that the expected number of disk replacements in a week varies by a factor of 9, depending on whether the preceding week falls into the first or third bucket, while we would expect no variation if failures were independent. When repeating the same process on the data of only year 3 of HPC1’s lifetime, we see a difference of a close to factor of 2 between the first and third bucket.

So far, we have only considered correlations between successive time intervals, e.g. between two successive weeks. A more general way to characterize correlations is to study correlations at different time lags by using the autocorrelation function. Figure 6 (left) shows the autocorrelation function for the number of disk replacements

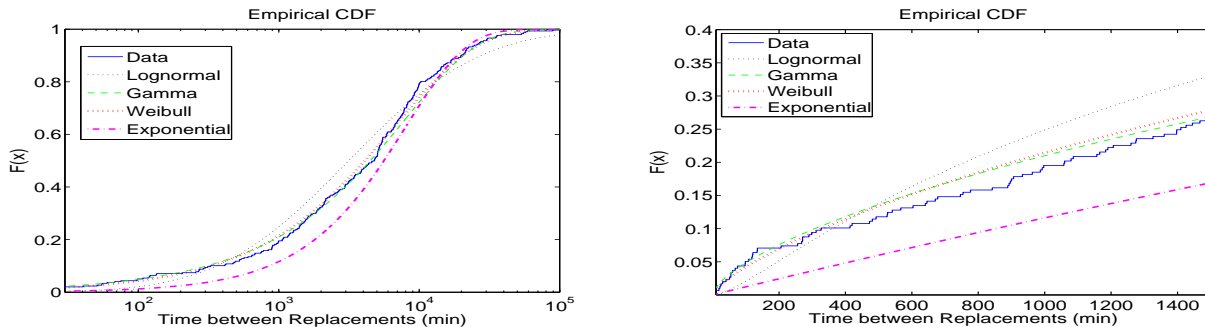


Figure 8: *Distribution of time between disk replacements across all nodes in HPC1.*

per week computed across the HPC1 data set. For a stationary failure process (e.g. data coming from a Poisson process) the autocorrelation would be close to zero at all lags. Instead, we observe strong autocorrelation even for large lags in the range of 100 weeks (nearly 2 years).

We repeated the same autocorrelation test for only parts of HPC1’s lifetime and find similar levels of autocorrelation. Figure 6 (right), for example, shows the autocorrelation function computed only on the data of the third year of HPC1’s life. Correlation is significant for lags in the range of up to 30 weeks.

Another measure for dependency is long range dependence, as quantified by the Hurst exponent  $H$ . The Hurst exponent measures how fast the autocorrelation functions drops with increasing lags. A Hurst parameter between 0.5–1 signifies a statistical process with a long memory and a slow drop of the autocorrelation function. Applying several different estimators (see Section 2) to the HPC1 data, we determine a Hurst exponent between 0.6-0.8 at the weekly granularity. These values are comparable to Hurst exponents reported for Ethernet traffic, which is known to exhibit strong long range dependence [16].

**Observation 7:** Disk replacement counts exhibit significant levels of autocorrelation.

**Observation 8:** Disk replacement counts exhibit long-range dependence.

### 5.3 Distribution of time between failure

In this section, we focus on the second key property of a Poisson failure process, the exponentially distributed time between failures. Figure 8 shows the empirical cumulative distribution function of time between disk replacements as observed in the HPC1 system and four distributions matched to it.

We find that visually the gamma and Weibull distributions are the best fit to the data, while exponential and

lognormal distributions provide a poorer fit. This agrees with results we obtain from the negative log-likelihood, that indicate that the Weibull distribution is the best fit, closely followed by the gamma distribution. Performing a Chi-Square-Test, we can reject the hypothesis that the underlying distribution is exponential or lognormal at a significance level of 0.05. On the other hand the hypothesis that the underlying distribution is a Weibull or a gamma cannot be rejected at a significance level of 0.05.

Figure 8 (right) shows a close up of the empirical CDF and the distributions matched to it, for small time-between-replacement values (less than 24 hours). The reason that this area is particularly interesting is that a key application of the exponential assumption is in estimating the time until data loss in a RAID system. This time depends on the probability of a second disk failure during reconstruction, a process which typically lasts on the order of a few hours. The graph shows that the exponential distribution greatly underestimates the probability of a second failure during this time period. For example, the probability of seeing two drives in the cluster fail within one hour is four times larger under the real data, compared to the exponential distribution. The probability of seeing two drives in the cluster fail within the same 10 hours is two times larger under the real data, compared to the exponential distribution.

The poor fit of the exponential distribution might be due to the fact that failure rates change over the lifetime of the system, creating variability in the observed times between disk replacements that the exponential distribution cannot capture. We therefore repeated the above analysis considering only segments of HPC1’s lifetime. Figure 9 shows as one example the results from analyzing the time between disk replacements in year 3 of HPC1’s operation. While visually the exponential distribution now seems a slightly better fit, we can still reject the hypothesis of an underlying exponential distribution at a significance level of 0.05. The same holds for other 1-year and even 6-month segments of HPC1’s lifetime. This leads us to believe that even during shorter segments

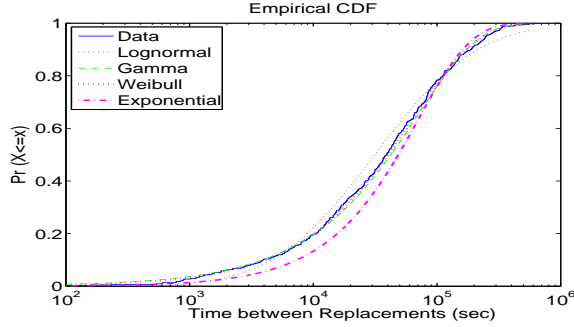


Figure 9: Distribution of time between disk replacements across all nodes in HPC1 for only year 3 of operation.

of HPC1’s lifetime the time between replacements is not realistically modeled by an exponential distribution.

While it might not come as a surprise that the simple exponential distribution does not provide as good a fit as the more flexible two-parameter distributions, an interesting question is what properties of the empirical time between failure make it different from a theoretical exponential distribution. We identify as a first differentiating feature that the data exhibits higher variability than a theoretical exponential distribution. The data has a  $C^2$  of 2.4, which is more than two times higher than the  $C^2$  of an exponential distribution, which is 1.

A second differentiating feature is that the time between disk replacements in the data exhibits decreasing hazard rates. Recall from Section 2.4 that the hazard rate function measures how the time since the last failure influences the expected time until the next failure. An increasing hazard rate function predicts that if the time since a failure is long then the next failure is coming soon. And a decreasing hazard rate function predicts the reverse. The table below summarizes the parameters for the Weibull and gamma distribution that provided the best fit to the data.

HPC1 Data	Distribution / Parameters			
	Weibull		Gamma	
	Shape	Scale	Shape	Scale
Compute nodes	0.73	0.037	0.65	176.4
Filesystem nodes	0.76	0.013	0.64	482.6
All nodes	0.71	0.049	0.59	160.9

Disk replacements in the filesystem nodes, as well as the compute nodes, and across all nodes, are fit best with gamma and Weibull distributions with a shape parameter less than 1, a clear indicator of decreasing hazard rates.

Figure 10 illustrates the decreasing hazard rates of the time between replacements by plotting the expected remaining time until the next disk replacement (Y-axis) as a function of the time since the last disk replacement (X-axis). We observe that right after a disk was replaced the

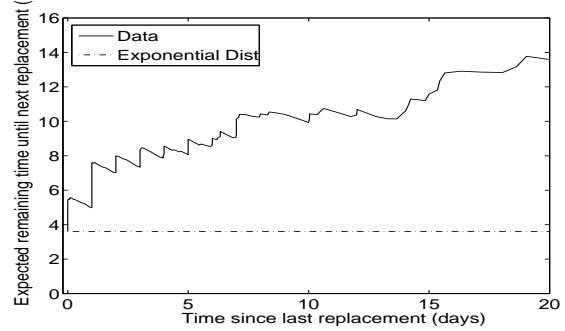


Figure 10: Illustration of decreasing hazard rates

expected time until the next disk replacement becomes necessary was around 4 days, both for the empirical data and the exponential distribution. In the case of the empirical data, after surviving for ten days without a disk replacement the expected remaining time until the next replacement had grown from initially 4 to 10 days; and after surviving for a total of 20 days without disk replacements the expected time until the next failure had grown to 15 days. In comparison, under an exponential distribution the expected remaining time stays constant (also known as the memoryless property).

Note, that the above result is not in contradiction with the increasing replacement rates we observed in Section 4.2 as a function of drive age, since here we look at the distribution of the time between disk replacements in a cluster, not disk lifetime distributions (i.e. how long did a drive live until it was replaced).

**Observation 9:** The hypothesis that time between disk replacements follows an exponential distribution can be rejected with high confidence.

**Observation 10:** The time between disk replacements has a higher variability than that of an exponential distribution.

**Observation 11:** The distribution of time between disk replacements exhibits decreasing hazard rates, that is, the expected remaining time until the next disk was replaced grows with the time it has been since the last disk replacement.

## 6 Related work

There is very little work published on analyzing failures in real, large-scale storage systems, probably as a result of the reluctance of the owners of such systems to release failure data.

Among the few existing studies is the work by Talagala et al. [29], which provides a study of error logs in a research prototype storage system used for a web server and includes a comparison of failure rates of different hardware components. They identify SCSI disk enclosures as the least reliable components and SCSI disks as one of the most reliable component, which differs from our results.

In a recently initiated effort, Schwarz et al. [28] have started to gather failure data at the Internet Archive, which they plan to use to study disk failure rates and bit rot rates and how they are affected by different environmental parameters. In their preliminary results, they report ARR values of 2–6% and note that the Internet Archive does not seem to see significant infant mortality. Both observations are in agreement with our findings.

Gray [31] reports the frequency of uncorrectable read errors in disks and finds that their numbers are smaller than vendor data sheets suggest. Gray also provides ARR estimates for SCSI and ATA disks, in the range of 3–6%, which is in the range of ARRs that we observe for SCSI drives in our data sets.

Pinheiro et al. analyze disk replacement data from a large population of serial and parallel ATA drives [23]. They report ARR values ranging from 1.7% to 8.6%, which agrees with our results. The focus of their study is on the correlation between various system parameters and drive failures. They find that while temperature and utilization exhibit much less correlation with failures than expected, the value of several SMART counters correlate highly with failures. For example, they report that after a scrub error drives are 39 times more likely to fail within 60 days than drives without scrub errors and that 44% of all failed drives had increased SMART counts in at least one of four specific counters.

Many have criticized the accuracy of MTTF based failure rate predictions and have pointed out the need for more realistic models. A particular concern is the fact that a single MTTF value cannot capture life cycle patterns [4, 5, 33]. Our analysis of life cycle patterns shows that this concern is justified, since we find failure rates to vary quite significantly over even the first two to three years of the life cycle. However, the most common life cycle concern in published research is underrepresenting infant mortality. Our analysis does not support this. Instead we observe significant underrepresentation of the early onset of wear-out.

Early work on RAID systems [8] provided some statistical analysis of time between disk failures for disks used in the 1980s, but didn't find sufficient evidence to reject the hypothesis of exponential times between failure with high confidence. However, time between failure has been analyzed for other, non-storage data in several studies [11, 17, 26, 27, 30, 32]. Four of the studies use

distribution fitting and find the Weibull distribution to be a good fit [11, 17, 27, 32], which agrees with our results. All studies looked at the hazard rate function, but come to different conclusions. Four of them [11, 17, 27, 32] find decreasing hazard rates (Weibull shape parameter  $< 0.5$ ). Others find that hazard rates are flat [30], or increasing [26]. We find decreasing hazard rates with Weibull shape parameter of 0.7-0.8.

Large-scale failure studies are scarce, even when considering IT systems in general and not just storage systems. Most existing studies are limited to only a few months of data, covering typically only a few hundred failures [13, 20, 21, 26, 30, 32]. Many of the most commonly cited studies on failure analysis stem from the late 80's and early 90's, when computer systems were significantly different from today [9, 10, 12, 17, 18, 19, 30].

## 7 Conclusion

Many have pointed out the need for a better understanding of what disk failures look like in the field. Yet hardly any published work exists that provides a large-scale study of disk failures in production systems. As a first step towards closing this gap, we have analyzed disk replacement data from a number of large production systems, spanning more than 100,000 drives from at least four different vendors, including drives with SCSI, FC and SATA interfaces. Below is a summary of a few of our results.

- Large-scale installation field usage appears to differ widely from nominal datasheet MTTF conditions. The field replacement rates of systems were significantly larger than we expected based on datasheet MTTFs.
- For drives less than five years old, field replacement rates were larger than what the datasheet MTTF suggested by a factor of 2–10. For five to eight year old drives, field replacement rates were a factor of 30 higher than what the datasheet MTTF suggested.
- Changes in disk replacement rates during the first five years of the lifecycle were more dramatic than often assumed. While replacement rates are often expected to be in steady state in year 2-5 of operation (bottom of the “bathtub curve”), we observed a continuous increase in replacement rates, starting as early as in the second year of operation.
- In our data sets, the replacement rates of SATA disks are not worse than the replacement rates of SCSI or FC disks. This may indicate that disk-independent factors, such as operating conditions, usage and environmental factors, affect replacement



rates more than component specific factors. However, the only evidence we have of a bad batch of disks was found in a collection of SATA disks experiencing high media error rates. We have too little data on bad batches to estimate the relative frequency of bad batches by type of disk, although there is plenty of anecdotal evidence that bad batches are not unique to SATA disks.

- The common concern that MTTFs underrepresent infant mortality has led to the proposal of new standards that incorporate infant mortality [33]. Our findings suggest that the underrepresentation of the early onset of wear-out is a much more serious factor than underrepresentation of infant mortality and recommend to include this in new standards.
- While many have suspected that the commonly made assumption of exponentially distributed time between failures/replacements is not realistic, previous studies have not found enough evidence to prove this assumption wrong with significant statistical confidence [8]. Based on our data analysis, we are able to reject the hypothesis of exponentially distributed time between disk replacements with high confidence. We suggest that researchers and designers use field replacement data, when possible, or two parameter distributions, such as the Weibull distribution.
- We identify as the key features that distinguish the empirical distribution of time between disk replacements from the exponential distribution, higher levels of variability and decreasing hazard rates. We find that the empirical distributions are fit well by a Weibull distribution with a shape parameter between 0.7 and 0.8.
- We also present strong evidence for the existence of correlations between disk replacement interarrivals. In particular, the empirical data exhibits significant levels of autocorrelation and long-range dependence.

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## Notes

<sup>1</sup>More precisely, we choose the cutoffs between the buckets such that each bucket contains the same number of samples (i.e. weeks) by using the 33th percentile and the 66th percentile of the empirical distribution as cutoffs between the buckets.

<sup>2</sup>This report was prepared as an account of work sponsored by an agency of the United States Government. Neither the United States Government nor any agency thereof, nor any of their employees, makes any warranty, express or implied, or assumes any legal liability or responsibility for the accuracy, completeness, or usefulness of any information, apparatus, product, or process disclosed, or represents that its use would not infringe privately owned rights. Reference herein to any specific commercial product, process, or service by trade name, trademark, manufacturer, or otherwise does not necessarily constitute or imply its endorsement, recommendation, or favoring by the United States Government or any agency thereof. The views and opinions of authors expressed herein do not necessarily state or reflect those of the United States Government or any agency thereof.

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