DFA on WolfSSL Ed25519

Seminar 11, PB173

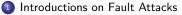
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2 Recall on Differential Fault Analysis (RSA)

Real-world example: fault attacks on WolfSSL



Attack categories

- Fault Injection was already partially covered on seminar 1.
- Even Differential Fault Analysis (DFA) was briefly mentioned, but I will recall it now.



Plan

1 Introductions on Fault Attacks

Recall on Differential Fault Analysis (RSA)

Real-world example: fault attacks on WolfSSL



Attack categories

- Side-channel attacks
 - use some physical (analog) characteristics
 - the target is running in normal conditions
- Faults: use abnormal conditions causing malfunctions in the system
- Micro-probing: accessing the chip surface directly in order to observe, learn and manipulate the device
- Reverse engineering



Types of implementation attacks

Active vs passive:

- Passive i.e. eavesdropping: the device operates within its specification
- Active i.e. tampering: the key is recovered by exploiting some abnormal behavior e.g. power glitches or laser pulses

Invasiveness:

- Non-invasive aka low-cost:
 - power/EM measurements
 - Coldboot attacks: data remanence in memories cooling down is increasing the retention time
 - Rowhammer is essentially a fault attack
- Semi-invasive: the device is de-packaged but no direct contact exists with the chip e.g. optical attacks
- Invasive aka expensive: the strongest type is bus probing

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Methods

- Variation in supply voltage i.e. glitching
 - Can cause a processor skip instruction
 - Actively investigated by smartcard industry
 - So-called unloopers were used to activate the infinity loop in PayTV smartcards
- Variation in the external clock: may cause data misread or an instruction miss
- Change in temperature
 - The temperature threshold is defined for which the chip will work properly
 - Can cause changes in RAM content
- White light: photons induce faults
- X-rays and ion beams



Goals

- Insert computational fault
 - Null key
 - Wrong crypto result (Differential Fault Analysis DFA)
- Change software decisions
 - Force approval of false PIN
 - Reverse life cycle state PayTV and old phone cards
 - Enforce access rights
 - Break secure boot



Practical Fault Injection Aspects and what we concentrate on in this lecture

- Most common FI: voltage and EM (due to its price)
 - https://github.com/newaetech/chipshouter-picoemp
- Differential Fault Analysis (DFA)
 - We mention a few advanced recent methods that strongly relate to SCA
- Glitching decisions:
 - secure boot
 - obtaining memory dumps
 - enabling debug interfaces



DFA

- Bellcore attack in 1995
 - Differential faults on RSA-CRT signatures
 - Requires 1 correct and 1 wrong signature
- Attack on <u>DES</u> in 1997 by Biham and Shamir
- Special attacks on <u>AES</u>, ECC etc.
- Fault attacks on key transfer

DFA on cryptosystems

- Basic DFA scenario:
 - adversary obtains a pair of ciphertexts that are derived by encrypting the same plaintext (one is correct value and the other is faulty)
 - two encryptions are identical up to the point where the fault occurred
 - \rightarrow two ciphertexts can be regarded as the outputs of a reduced-round iterated block cipher where the inputs are unknown but show a small (and possibly known) differential
- DFA on DES
 - the original attack of Biham and Shamir exploits computational errors occurring in the final rounds of the cipher
 - assumes that one bit of the right half of the DES internal state is flipped at a random position



Recall from seminar 1: RSA with CRT

Optimization of computing a signature giving about 4-fold speedup: $n = p \cdot q$ Signature: $s = m^d \mod n$

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\begin{array}{ll} \text{Pre-computed values } d_{\rho}:=d \mod (p-1) \qquad d_q:=d \mod (q-1) \\ i_q:=q^{-1} \mod p \end{array}
```

 $s_p := m^{d_p} \mod p$ $s_q := m^{d_q} \mod q$

Garner's method (1965) to recombine s_p and s_q : $s = s_q + q \cdot (i_q(s_p - s_q) \mod p)$

Where to glitch?

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Almost anywhere :-) computations of s_p and s_q.
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If error is in s_p then the adversary can recover q as follows: $q = \gcd(n, s - \hat{s})$.

Plan

Introductions on Fault Attacks

Recall on Differential Fault Analysis (RSA)

Real-world example: fault attacks on WolfSSL



Ed25519

- Instance of EdDSA, which was proposed to "fix the unnecessary requirements on randomness" in ECDSA
- Does not depend on a "good" source of randomness, but instead derives a secret deterministically (hashing the msg and a long-term auxiliary key)
- Widely adopted by TLS1.3, Zcash, SSH, Tor, Signal, WolfSSL etc. (check "Things that use Ed25519")
- Turns out to be easy to attacks in some real-world deployments i.e. WolfSSL

Niels Samwel, Lejla Batina, Guido Bertoni, Joan Daemen and Ruggero Susella: *Breaking Ed25519 in WolfSSL*, CTRSA2018. Niels Samwel, Lejla Batina: *Practical Fault Injection on Deterministic Signatures: the Case of EdDSA*, Africacrypt 2018.

CROC

Ed25519

Algorithm 1 Ed25519 key setup and signature generation

Key setup.

- 1: Hash k such that $H(k)=(h_0,h_1,\ldots,h_{2b-1})=(a,b)$
- 2: $a = (h_0, \ldots, h_{b-1})$, Private scalar
- 3: $b = (h_b, \ldots, h_{2b-1})$, Auxiliary key
- 4: Compute public key: A = aB.

Signature generation.

- 5: Compute ephemeral private key: r = H(b, M).
- 6: Compute ephemeral public key: R = rB.
- 7: Compute h = H(R, A, M) and convert to integer.
- 8: Compute: $S = (r + ha) \mod I$.
- 9: Signature pair: (R, S).



The Attack

Two signatures, original (R, S) and faulty (R', S'):

$$S = r + ha$$

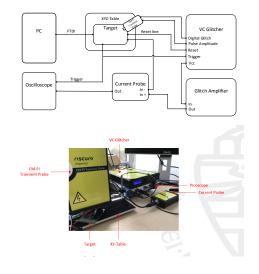
 $S' = r + h'a$

$$S - ha = S' - h'a$$

$$a = \frac{S - S'}{h - h'}$$

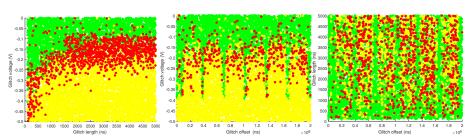


Setup





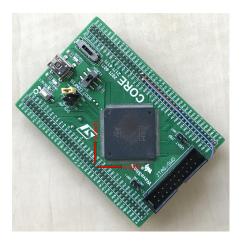
Results



Voltage fault injection results, Normal (green), Inconclusive (yellow), Successful (red).



Results





x-axis



Conclusion

Two real physical side-channel attacks were actually performed against Ed25519

- Side-channel analysis of Ed25519 with 4000 traces
- Fault injection on Ed25519 with 100% success rate for EM FI and 70% for voltage glitching out of 10 000 measurements
- For both attacks there exist inexpensive countermeasures



Questions



