



PHYLOGENY RECONSTRUCTION: THE BASICS











A Simple Concept of Speciation



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The Coalescent Model





substitutions.

The evolution of biological sequences





The evolution of biological sequences





Biological sequences can change by *substitutions*



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Divergence and Coalescent





What we have...





What do we get:

acgggatcccattac acgggatcccaatac ccgggatagcttccattac acgacatatccactggattcc acccctatccactggattac

A collection of homologous sequence that vary slightly in their nucleotide or amino acid composition and in their length

What we want...





A collection of homologous sequence that vary slightly in their nucleotide or amino acid composition



Public Data Sources





The Problem: Finding the homologous positions

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An mathematical function able to measure the biological quality of an alignment...



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Related questions:

>What should a biologically correct alignment look like?

➤To what extent can we define and formalize its properties?



An mathematical function able to measure the biological quality of an alignment...

Related questions:

>What should a biologically correct alignment look like?

>To what extent can we define and formalize its properties?





A mathematical function meant to measure the biological quality of an alignment...

$$\sigma(\alpha) = \sum_{i=1}^{n} S(a_i, b_i)$$

 $\sigma(\alpha)$: the score of the pairwise alignment α

- n : length of α
- a_i : letter of sequence A at position i in α
- b_i : letter of sequence B at position i in α



A mathematical function meant to measure the biological quality of an alignment...

$$\sigma(\alpha) = \sum_{i=1}^{n} S(a_i, b_i)$$

Objective: find α that maximizes $\sigma(\alpha)$!



Given two sequences A ={ $a_1, a_2, ..., a_n$ } and B={ $b_1, b_2, ..., b_m$ } and a scoring function S such that

$$S(a_i, b_j) = \begin{cases} +5, if \ a_i = b_j \\ -2, if \ a_i \neq b_j \\ -6, \text{ for introduction of a gap} \end{cases}$$

then we look for that alignment, that gives us the highest score by summing up the column scores $S(a_i,b_j)$ for all columns of the alignment.



Given two sequences A ={ $a_1, a_2, ..., a_n$ } and B={ $b_1, b_2, ..., b_m$ } and a scoring function S such that

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A1: T G C T C G T A
T
$$-$$
 T C A T A
 $+5$ -6 -6 $+5$ $+5$ -2 $+5$ $+5$ = 11

continue.....

A2: T G C T C G T A
T - T - C A T A
$$+5$$
 -6 -2 -6 $+5$ -2 $+5$ +5 = 4

etc...



• There are far too many

> number of possible pairwise alignments: $\binom{2n}{n}$

> for two sequences of length N=300 there are 10^{179} possibilities



• There are far too many

number of possible pairwise alignments: $\binom{2n}{n}$ for two sequences of length N=300 there are 10¹⁷⁹

possibilities

Hence, we need a smart way to cut the computation short, like the **dynamic programming** approach for pairwise alignments by *Needleman and Wunsch* (1970).

Re-use of previous results





etc...



A **dynamic programming** approach usually includes:

- A mathematical description of the (biological) quality of an solution, i.e. an recursive objective function
- The computation of all intermediate values needed to obtain the globally optimal solution, thereby avoiding double-computations
- The reconstruction of the globally optimal solution from the values obtained in the previous step (backtracking)

The Needleman-Wunsch pair-wise alignment



Scoring function

$$S(a_i, b_j) = \begin{cases} +5, if \ a_i = b_j \\ -2, if \ a_i \neq b_j \\ -6, \text{ for introduction of a gap} \end{cases}$$

Objective function

$$\sigma(i, j) = \max \begin{cases} \sigma(i-1, j-1) + S(a_i, b_j) \\ \sigma(i, j-1) + S(gap, b_j) \\ \sigma(i-1, j) + S(a_i, gap) \end{cases}$$

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The Needleman-Wunsch algorithm





> $\sigma(i,j)$ is the optimal alignment score up to and including a_i and b_i

 $S(a_i, b_j) = \begin{cases} +5, if \ a_i = b_j \\ -2, if \ a_i \neq b_j \\ -6, \text{ for introduction of a gap} \end{cases}$

Needleman-Wunsch algorithm: Initialization



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 $S(a_i, b_j) = \begin{cases} +5, if \ a_i = b_j \\ -2, if \ a_i \neq b_j \\ -6, \text{ for introduction of a gap} \end{cases}$

The Needleman-Wunsch algorithm: Recursion



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$$S(a_i, b_j) = \begin{cases} +5, if \ a_i = b_j \\ -2, if \ a_i \neq b_j \\ -6, \text{ for introduction of a gap} \end{cases}$$



Needleman-Wunsch algorithm: Backtrack





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Needleman-Wunsch algorithm: Backtrack




Needleman-Wunsch algorithm: Backtrack





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Needleman-Wunsch algorithm: Backtrack





Alignment Score: 11

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Smith-Waterman pairwise local alignment

$$S(a_i, b_j) = \begin{cases} +5, if \ a_i = b_j \\ -2, if \ a_i \neq b_j \\ -6, \text{ for introduction of a gap} \end{cases}$$

$$\sigma(i, j) = \max \begin{cases} \sigma(i-1, j-1) + S(a_i, b_j) \\ \sigma(i, j-1) + S(gap) \\ \sigma(i-1, j) + S(gap) \\ 0 \end{cases}$$



Smith-Waterman pairwise local alignment





Alignment Score: 18

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Affine Gap costs



$$g(l) = g_o + l * g_e$$

$$\sigma(i, j) = \max \begin{cases} \sigma(i-1, j-1) + S(a_i, b_j) \\ \sigma(i, j-1) + S(gap, b_j) \\ \sigma(i-1, j) + S(a_i, gap) \end{cases}$$







Alternative Scoring Functions





PAM250:

C Cys	12																			
S Ser	0	2																		
T Thr	-2	1	3																	
P Pro	-3	1	0	6																
A Ala	-2	1	1	1	2	Î														
G Gly	-3	1	0	-1	1	5														
N Asn	-4	1	0	-1	0	0	2		_											
D Asp	-5	0	0	-1	0	1	2	4												
E Glu	-5	0	0	-1	0	0	1	3	4	L										
Q Gln	-5	-1	-1	0	0	-1	1	2	2	4										
H His	-3	-1	-1	0	-1	-2	2	1	1	3	6									
R Arg	-4	0	-1	0	-2	-3	0	-1	-1	1	2	6								
K Lys	-5	0	0	-1	-1	-2	1	0	0	1	0	3	5		2					
M Met				_						_										
INI INTEL	-5	-2	-1	-2	-1	-3	-2	-3	-2	-1	-2	0	0	6						
I Ile	-5 -2	-2 -1	-1 0	-2 -2	-1 -1	-3 -3	-2 -2	-3 -2	-2 -2	-1 -2	-2 -2	0 -2	0 -2	6 2	5					
I Ile L Leu	-5 -2 -6	-2 -1 -3	-1 0 -2	-2 -2 -3	-1 -1 -2	-3 -3 -4	-2 -2 -3	-3 -2 -4	-2 -2 -3	-1 -2 -2	-2 -2 -2	0 -2 -3	0 -2 -3	6 2 4	5 2	6	1			
I Ile L Leu V Val	-5 -2 -6 -2	-2 -1 -3 -1	-1 0 -2 0	-2 -2 -3 -1	-1 -1 -2 0	-3 -3 -4 -1	-2 -2 -3 -2	-3 -2 -4 -2	-2 -2 -3 -2	-1 -2 -2 -2	-2 -2 -2 -2	0 -2 -3 -2	0 -2 -3 -2	6 2 4 2	5 2 4	6 2	4		_	
I Ile L Leu V Val F Phe	-5 -2 -6 -2 -4	-2 -1 -3 -1	-1 0 -2 0 -3	-2 -2 -3 -1	-1 -1 -2 0 -5	-3 -3 -4 -1	-2 -2 -3 -2 -4	-3 -2 -4 -2 -6	-2 -2 -3 -2 -5	-1 -2 -2 -2	-2 -2 -2 -2 -2	0 -2 -3 -2 -4	0 -2 -3 -2 -5	6 2 4 2 0	5 2 4 1	6 2 2	4	9	i	
I Ile L Leu V Val F Phe Y Tyr	-5 -2 -6 -2 -4 0	-2 -1 -3 -1 -3 -3	-1 0 -2 0 -3 -3	-2 -2 -3 -1 -5 -5	-1 -1 -2 0 -5 -3	-3 -3 -4 -1 -5 -5	-2 -2 -3 -2 -4 -2	-3 -2 -4 -2 -6 -4	-2 -2 -3 -2 -5 -4	-1 -2 -2 -2 -5 -4	-2 -2 -2 -2 -2 0	0 -2 -3 -2 -4 -4	0 -2 -3 -2 -5 -4	6 2 4 2 0 -2	5 2 4 1 -1	6 2 2 -1	4 -1 -2	9 7	10	
I Ile L Leu V Val F Phe Y Tyr W Trp	-5 -2 -6 -2 -4 0 -8	-2 -1 -3 -1 -3 -3 -2	-1 0 -2 0 -3 -3	-2 -2 -3 -1 -5 -5 -6	-1 -1 -2 0 -5 -3 -6	-3 -3 -4 -1 -5 -5 -7	-2 -2 -3 -2 -4 -2 -4 -4	-3 -2 -4 -2 -6 -4 -7	-2 -2 -3 -2 -5 -4 -7	-1 -2 -2 -5 -4 -5	-2 -2 -2 -2 -2 0 -3	0 -2 -3 -2 -4 -4 2	0 -2 -3 -2 -5 -4 -3	6 2 4 2 0 -2 -4	5 2 4 1 -1 -5	6 2 -1 -2	4 -1 -2 -6	9 7 0	10 0	17

Many others...





Both, Needleman-Wunsch and Smith-Waterman alignment methods are **exact** methods since they guarantee a globally optimal solution for the optimization problem!

Drawback: Computational expensive, i.e. O(nm) in time and memory



Solutions:

omit regions from the grid, that cannot contribute to the optimal alignment (reduction of the search space, by remaining exact)





Solutions:

 use of heuristics (more rigorous reduction of the search space, sacrificing the guaranteed optimal solution for search speed)

Hashing



• Lookup method for finding an alignment

Pos:	1	2	3	4	5	6	7	8	9	10	11
Seq 1:	k	С	S	р	t	а	•	•	•	•	•
Seq 2:	•	•	•	•	•	a	С	S	р	r	k

Amino acid	Pos in Seq 1	Pos in Seq 2	Offset
k	1	11	10
С	2	7	-5
S	3	8	-5
р	4	9	-5
t	5	-	-
а	6	6	0
r	-	10	-

Hashing



• Lookup method for finding an alignment

Pos:	1	2	3	4	5	6	7	8	9	10	11
Seq 1:	k	С	S	р	t	a	•	•	•	•	•
Seq 2:	•	•	•	•	•	a	С	S	р	r	k

Amino acid	Pos in Seq 1	Pos in Seq 2	Offset
k	1	11	10
С	2	7	-5
S	3	8	-5
р	4	9	-5
t	5	-	-
а	6	6	0
r	-	10	-

Resulting alignment: Seq 1: k c s p t a

Seq 2: a c s p r k

What we are really looking for:



Ther_tengcongensis Clos_acetobutylicum Clos_tetani Desu_desulfuricans Vibr_vulnificus Caul_crescentus Micr_degradans Vibr_cholerae Shew_oneidensis Rat_beta1_sGC Rat_beta2_sGC Nost_punctiforme Nost_sp. consensus>50

Ther_tengcongensis Clos_acetobutylicum Clos_tetani Desu_desulfuricans Vibr_vulnificus Caul_crescentus Micr_degradans Vibr_cholerae Shew_oneidensis Rat_beta1_sGC Rat_beta2_sGC Nost_punctiforme Nost_sp. consensus>50

Ther_tengcongensis Clos_acetobutylicum Clos_tetani Desu_desulfuricans Vibr_vulnificus Caul_crescentus Micr_degradans Vibr_cholerae Shew_oneidensis Rat_beta1_sGC Rat_beta2_sGC Nost_punctiforme Nost_sp. consensus>50





Optimal Solution: Extend Needleman-Wunsch or Smith-Waterman to multiple sequences

How to construct Multiple Sequence Alignments?

Optimal Solution: Extend Needleman-Wunsch or Smith-Waterman to multiple sequences

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How to construct Multiple Sequence Alignments?

Optimal Solution: Extend Needleman-Wunsch or Smith-Waterman to multiple sequences

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But O(n^m) in time and memory:

Computationally not feasible... 4 sequences of length 1000 -> 1TB RAM

A new objective function: Sum of Pairs

Seq1: AGA--CTA Seq2: G-A--CTT Seq3: AGAAACTT CIBIV

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A new objective function: Sum of Pairs



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A new objective function: Sum of Pairs



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SUM OF PAIRS SCORE: 16



- The sequences are added stepwise. Thus, never more than two sequences (or multiple sequence alignments) are simultaneously aligned
- Sequences or MSAs are aligned using Dynamic
 Programming

Progressive Alignment Strategies (ClustalW)



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Scoring for the alignment of two alignments

$$\sigma(a^{i},b^{j}) = \frac{1}{n+m} \sum_{x=1}^{n} \sum_{y=1}^{m} S(a^{i}_{x},b^{j}_{y}) \times \omega_{x} \times \omega_{y}$$

 $\sigma(a^i, b^j)$:score for aligning column i from alignment (or sequence) \boldsymbol{a} to
column j from alignment or sequence \boldsymbol{b} n,mnumber of sequences in alignments \boldsymbol{a} and \boldsymbol{b} , respectively $S(a^i_x, b^j_y)$ score for aligning position \boldsymbol{i} in sequence \boldsymbol{x} from alignment \boldsymbol{a} to
position \boldsymbol{j} in sequence \boldsymbol{y} from alignment \boldsymbol{b} $\omega_{\boldsymbol{x}}, \omega_{\boldsymbol{y}}$ respective weights of the sequences \boldsymbol{x} and \boldsymbol{y}

Scoring for the alignment of two alignments

$$\sigma(a^{i},b^{j}) = \frac{1}{n+m} \sum_{x=1}^{n} \sum_{y=1}^{m} S(a^{i}_{x},b^{j}_{y}) \times \omega_{x} \times \omega_{y}$$



With sequence weights: Score = $(S(t,v)*\omega_1\omega_5$ + $S(t,i)*\omega_1\omega_6$ + $S(1,v)*\omega_2\omega_5$ + $S(1,i)*\omega_2\omega_6$ + $S(k,v)*\omega_3\omega_5$ + $S(k,i)*\omega_3\omega_6$ + $S(k,v)*\omega_4\omega_5$ + $S(k,i)*\omega_4\omega_6)/8$

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Features of ClustalW



- progressive strategy
- Distance based generation of a guide tree (approximative or exact)
- tree-guided (NJ) alignment
- change of the scoring matrix as the alignment proceeds
 (adaptation to increasing divergence of the sequences)
- dynamic variation of gap penalties in position- and residue-specific manner
 - gap opening penalties are locally reduced in stretches of 5 or more hydrophilic residues (indicative of loop or random coil regions).
 - gap penalties are locally increased within eight residues of existing gaps.
- sequence weighting

(Known) Problem of ClustalW: Local Optima

a.k.a: Once a gap always a gap

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Iterative Alignment Strategy



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Stochastic Iterative Alignment





Stochastic Iterative Alignment (SAGA)





Non-Stochastic Iterative Alignment

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Point: The initial alignment is modified by splitting it into two groups and re-aligning them with dynamic programming.



Example: Prrp, both, alignment (inner loop) and tree/weight (outer loop) are optimized.

Consistency based algorithm



Point: The optimal MSA is defined as the one that agrees the most with all optimal pair-wise alignments

Features:

- > does not depend on a specific substitution rate
- \succ can apply any method capable to align two sequences
- > position dependent, i.e. the score associated with the alignment of two residues depends on their position within the sequence rather that their individual nature
- rationale: given a set of independent observations, the constellation most often observed is often closer to the truth

Consistency based Objective Function For alignEment Evaluation (COFFEE)

The Principle of T-Coffee







Table 2. Some elements of validation on BAliBASE.									
Method	Ref1	Ref2	Ref3	Ref4	Ref5	Total			
DiAlign	71.0	25.2	35.1	74.7	80.4	57.3			
ClustalW	78.5	32.2	42.5	65.7	74.3	58.7			
Prrp	78.6	32.5	50.2	51.1	82.7	59.0			
T-Coffee	80.7	37.3	52.9	83.2	88.7	68.7			

Each method in the Method column was used to align the 141 test-sets contained in BAIiBASE. The alignments were then compared with the reference BAIiBASE alignment using aln_compare [34]. Ref1–5 indicates the five BAIiBASE categories. Results obtained in each category were averaged. All the observed differences are statistically significant, as assessed by the Wilcoxon rank-based test [34,47]. Ref1 contains a homogenous set of sequences, ref2 contains a homogenous group of sequences and an outlayer, ref3 contains two distantly related groups of sequences. Ref4 contains sequences that require long internal gaps to be properly aligned and ref5 contains sequences that require long-terminal gaps to be properly aligned. Total is the average of ref1–5.



Reconstructing Trees from Sequences Reconstructing the tree of life









Notations

How many possible trees are there?





How many possible trees are there?





$b(n) = \frac{(2n-5)!}{2^{n-3}(n-3)!}$
b(10) = 2027025
$b(55) = 2.9 \times 10^{84}$
$b(100) = 1.7 \times 10^{182}$
Finding the root of the tree





Finding the root of the tree











Three typical representations of the same tree



#NEXUS begin taxa; dimensions ntax=7; taxlabels DROME CIOIN **HYDMA** SACCE CAEEL ACRMI HOMSA ; end; begin trees; tree [&r] tree 1 = (((((CIOIN:0.4222,HOMSA:0.2777) [&label=97]:0.0575, (ACRMI:0.2611,HYDMA:0.37)[&label=100]:0.0745) [&label=100]:0.0764, DROME:0.42)[&label=100]:0.1034,CAEEL:0.6027):0.5804, SACCE:0.5832); end: begin figtree; set appearance.backgroundColour=#-1;

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(((((CIOIN:0.4222,HOMSA:0.2777)97:0.0575,(ACRMI:0.2611,HYD MA:0.3700)100:0.0745)100:0.0764,DROME:0.4200)100:0.1034, CAEEL:0.6027):0.5804,SACCE:0.5832);

end figtree;

Complex trees





(((((((Rhizopus oruzae:0.12338953118750640991.Cunninghamella elegans:0.14505462914970748689)100:0.15326253481223306441.(((((Clocophu) um_trabeuro:0.0724812458330700304.((Antrogramma_regramma_ 819373257)1000,0253950558505581476,(Paxillus_involutus:0,12507330174040107763,(Lentinula_edades:0,09578001724170833660,((Coprinopsis_cin erea:0.08955096563898440828,Hebeloma_cylindrosporum:0,10085378088044757994)100:0,01366340764780078071,Pleurotus_ostreatus:0,11634253457851 556623)100:0.01235187547388479599)100:0.01711017593272916998)100:0.01240279259156270436)100:0.01259472218900263042)100:0.10333051305697128 763.(Cryptococcus laurentii:0.11261447295515064626.(Cryptococcus neoformans.var.neoformans:0.00521778134981755806.(Filobasidiella neoforma /rss.(crgb/bcbcccus_laurencii)0,11201447/235150640625,(crgb/bcbcccus_nebrofmans;0*,nebrofmans;0*,000217/0134301750607,11100as101811a_nebrofmans;0*,000217/0134301750802170100;0*,0002122975341575844001100;0*,0002122975341575844001100;0*,000212297534157584400100;0*,000212297534157584400100;0*,000212297534157584400100;0*,000212297534157584400100;0*,000212297534157584400100;0*,000212297534157584400100;0*,00021229754157584400140035095205,U* omyces_apendiculatus:0*,01246074123125495556100;0*,01477257364870681322)100:0*,0646058920248333127100:0*,0686017462332440235,U*tilag_may dis:0*,25934482229560420327)100:0*,07030546047039644419.(((Saitopila_complicatio_1)226631741538455434, Laphrina_deformans:0*,570307905833770 0642)100:0.03883872443363502874,Schizosaccharomyces_pombe:0.33367484237741701358)100:0.03345088821164741988,(((Debaromyces_hansenii:0.1660 7000789678139085, (Saccharomyces_cerevisiae:0.07821550711422489699, Candida_glabrata:0.08167802849027459844)100:0.16114727512467122428)100:0 .09989437041082575852.Yarrowia_lipolytica:0.22680534204146188260)100:0.12819522989497894594.(Tuber_borchii:0.18713075564612405288.((((Para .035845704106237852,Tarrow1a_11po1gtCat0,2258054204146183250710010,12813522353437834534,(10Der_DorCnit10,18713073564612405285,(((Tran coccidioides_brasiliensis0,04452412585008645487,Ajellowyces_capsulatus10.0554615509552374164110010,0558247320008983514,(Trichophyton_ru brum:0,13478802145894250297,(Coccidioides_immitis0,0009861500455473006,Coccidioides_posadasii:0,00013275597489488007)10010,06751130311481 0437469)1001,01752185514248328643)10010,02570153410152175091,(Thermomyces_langinosus:0,14047726034864610467,(Aspergillus_flavus:0,136258 86602121553559,((Emericella_nidulans:0,019378923659861395,Aspergillus_nidulans:0,0117495551954378149)1001:0,0511031254373287538,4578149)1001 .07010364253857682970)100:0.01568149707407409318)100:0.03916392516601811119)100:0.02443387059988460716)100:0.05504556307901582041,(((Mycos phaerella_graminicola:0.11382805992322068966.Aureobasidium_pullulans:0.13479608922737240650)100:0.03222541190278756240.((Phaeosphaeria_nod orum:0.03601370297481657629.Leptosphaeria maculans:0.07294024018301276113)100:0.02604153023893227351.Alternaria brassicicola:0.32946258067 0rum:0.05013702374016376275316916376275129163771100:0.03664059956326781720.((((Verticillium_dahliae:0.1374302698027823891,Glomerella_cingulata:0.05 2662371278755867362)100:0.0265933622261522701100:0.03664059956326781720.(((Verticillium_dahliae:0.1374302698027823891,Glomerella_cingulata:0.05 2662371278755867362)100:0.02659336222615227021.((((ibberella_caee:0.00009861600456473006,Fusarium_org).00009861600456473006,Fusarium_00009861600456473006,Fusarium_000009861600456473006,Fusarium_000009861600456473006,Fusarium_000009861600456473006,Fusarium_000009861600456473006,Fusarium_000009861600456473006,Fusarium_000009861600456473006,Fusarium_000009861600456473006,Fusarium_000009861600456473006,Fusarium_000009861600456473006,Fusarium_000009861600456473006,Fusarium_000009861600456473006,Fusarium_000009861600456473006,Fusarium_000009861600456473006,Fusarium_000009861600456473006,Fusarium_000009861600456473006,Fusarium_000009861600456473006,Fusarium_00000986160045647 f.sp.cucumerinum:0.00473059578454169526)100:0.00357223653413873228)100:0.01712573793927751192,(Gibberella_moniliformis:0.00009861600456473 006.Fusarium verticillioides:0.00009861600456473006)100:0.00755012976207619646)100:0.01660747744541156770)100:0.04668007882194781522.(Meta 006;rusarlum_verclf1110106;0,0005306;0000;000;000;000;50125760/0015606;000;000;600647/443411367/07100;0,0400;0 rhizium_anisopliae:0,06747750588644643721,((Hypocrea_jecorina:0,04281188204986635084,((Trichoderma_stroviride:0,02788662551141205857, Trich oderma_asperellum:0,02793167656533112878)100:0,03084253786435780034,(Hypocrea_virens:0,01967361655791701955,(Hypocrea_lixii:0,000098616004 56473006,Trichoderma_harzianum:0,00009861600456473005)100:0,01640815227372549590)100:0,01031552333099211745)100:0,01407183499215602582)100 :0.03728293661804039927, (Beauveria_bassiana:0.00009861600456473006, Cordyceps_bassiana:0.00017480976326447167)100:0.11379015715330136060)10 0;0,01457733600595170070)100;0,01214032674197134364)100;0,03171283347589626267)100;0,01752107533795072003,Cryphonectria_parasitica;0,12585 408689869942434)100:0.01404062191232230070.((Magnaporthe grisea:0.14740839006608871808.((Neurospora crassa:0.09790422357055546254.Chaetomi um_cupreum:0.03622272647936440604)80:0.02088625388322065332.Corynascus_heterothallicus:0.08006458123554179018)100:0.02903048934692216074)8 00.0164997564221958331.(Ophiostoma_piliferum:0.09029085190095893776.Ophiostoma_clavigerum:0.05973185389895070452)100;0.06271238182555409 431)80:0.01044724186632970631)100:0.05716276100425224382.(((Sclerotinia_sclerotiorum:0.0387972320106123483,(Botrytis_cinerea:0.0000986160 0456473006,Botryotinia_fuckelianaC:0.00009861600456473006)100:0.01606152746805612069)100:0.06410552334506864079,(Amorphotheca_resinae:0.07 099850624761384899,Blumeria_graminisf.sp.hordei:0,12573210322159533714)100:0,01007185390375272237)100:0,01372048555964769899,Geomyces_pann orum:0.08942182088142704155)100:0.03434252931156196037)100:0.03464174207103632580)100:0.01879298933467762195)100:0.05538850743967564660)10 0:0.09299628620991588768)100:0.04867625580472154101)100:0.07498954862375889485)100:0.06386475531945318140)100:0.03914280476745674031,((Spi zellomyces punctatus:0.18931369968937947212.Neocallimastix patriciarum:0.28843028482622112829)100:0.04371345560798906016.(Allomyces macrog ynus:0,11459719607743417145,Blastocladiella_emersonii:0,11999224967666824448)100:0.20328647132686597510)100:0,03664269603944932119)100:0,1 0526374086763891358,Dictyostelium_discoideum:0,70552689158042469764)100:0,11685579381042721092,Caenorhabditis_elegans:0,421899305318440387 41)100;0.04434817203516049772, Drosophila_melanogaster:0.31104455613948034376)100;0.06514413900702456517, Ciona_intestinalis:0.2957740483821 7330463)100:0.09667503731183066384,Homo_sapiens:0.09667503731183066384)100;

Tree display is an unsolved problem





Welcome to iTOL!

Interactive Tree Of Life is an online tool for the display and manipulation of phylogenetic trees. It provides most of the features available in other tree viewers, and offers a novel circular tree layout, which makes it easy to visualize mid-sized tree (up to several thousand leaves). Trees can be exported to several graphical formats, both bitmap and vector based. more...

NEW! If you are using iTOL to upload and display your own trees, you can create a personal iTOL account. It will allow you to access your trees from anywhere, organize them into workspaces and projects and easily manage datasets and other tree features. Detailed list and explanation of available features can be accessed here.

înis is	the new descrip	ition of the project.	
	Tree	Description	Datasets
: + +	Tree #1	Initial parameters, not modified	
	Nematodes	Uploaded tree	
	Metazoa	Uploaded tree	
	Tree 3	Uploaded tree	

If you already have an account, go to the login page.



Some more notations





Character based phylogeny reconstruction:

- A character has to be expressed in at least two states in the taxa under study.
 Taxa are grouped on the basis of shared character states.
- An evolutionary derived character (state) is called an *Apomorphy*
- Aut-Apomorphy: an evolutionary derived character (state) present only in a single taxon
- Syn-Apomorphy: an evolutionary derived character (state) shared by a group of taxa.
- Plesiomorphy: an ancestral character (state) shared by a group of extant taxa.
- Homoplasy: A derived character (state) that is shared for reasons other than common decent.

The Parsimony Principle





A rule in science and philosophy stating that entities should not be multiplied needlessly.

This rule is interpreted to mean that the simplest of two or more competing theories is preferable and that an explanation for unknown phenomena should first be attempted in terms of what is already known.

Also called law of parsimony. (Ockham's razor, ca 1285-1350)





The Parsimony Principle







Data	Method	Evaluation Criterion			
(Maximum Parsimony	Parsimony			
Characters (Alignment)	Statistical Approaches: Likelihood, Bayesian	Evolutionary Models			
Distances	Distance Methods	J			

The Criterion of Maximum Parsimony



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Find the tree τ that minimizes the following expression:

$$L(\tau) = \sum_{k=1}^{B} \sum_{j=1}^{L} \omega_j \cdot \operatorname{diff}(x_{k'j}, x_{k''j})$$

where diff measures the distance between two characters $\omega_{j}\,\text{is}$ an alignment specific weight factor

L alignment length

B number of branches in the tree

k['] and *k*^{''} are the two nodes connected by branch *k*

The Criterion of *Distance (Hamming Distance)*



а	g	С	t	t	а	С	С	t	g	t	t	а	С	t
2 C	g	t	а	а	a	t	t	t	С	С	С	g	а	t
3 C	g	С	а	а	g	t	t	t	С	С	С	g	а	t
l c	а	С	t	t	а	t	t	а	g	t	С	а	а	С
						↓ ↓								
		Se	eq	1	Se	eq 2	2	Sec	д 3	5	Seq	4		
Seq	1	0			11			11		8	3			
Seq 2		11		0			2		1	10				
Seq	3	11	1		2			0		g)			
	a c c c Seq Seq Seq	a g c g c g c a Seq 1 Seq 2 Seq 3	a g c c g t c g c c a c Seq 1 0 Seq 2 12 Seq 3 12	a g c t c g t a c g c a c a c t Seq 1 0 Seq 2 11 Seq 3 11	a g c t t c g t a a c g c a a c a c t t Seq 1 0 Seq 2 11 Seq 3 11	a g c t t a c g t a a a c g c a a g c a c t t a Seq 1 0 11 Seq 2 11 0 Seq 3 11 2	a g c t t a c c g t a a a t c g c a a g t c a c t t a t Seq 1 0 11 Seq 2 11 0 Seq 3 11 2	a g c t t a c c c g t a a a t t c g c a a g t t c a c t t a t t Seq 1 0 11 Seq 2 11 0 Seq 3 11 2	a g c t t a c c t c g t a a a t t t c g c a a g t t t c a c t t a t t a Seq 1 0 11 11 Seq 2 11 0 2 Seq 3 11 2 0	a g c t t a c c t g c g t a a a t t t c c g c a a g t t t c c a c t t a t t a g Seq 1 0 11 11 Seq 2 11 0 2 Seq 3 11 2 0	a g c t t a c c t g t c g t a a a t t t c c c g c a a g t t t c c c a c t t a t t a g t Seq 1 0 11 11 8 Seq 2 11 0 2 1 Seq 3 11 2 0 9	a g c t t a c t	a g c t t a c c t t t t t t a a t t t a a a t t t t a a a t t t t a a t t t t a a t t t t a a t t t t a a t t t a a t t t t t a a t t t t a a t t t t t t a a t t t t t t a a t t t a a t	a g c t t a c c t g t t a c c g t a a a t t t c c c g a c g c a a g t t t c c c g a c a c t t a t t a g t c a a $ \begin{array}{c} $





Find branch lengths L(b) such that the sum of the branch lengths connecting any two leaves gets close to the measured distances between all pairs of leaves. That is

$$D_{\text{measured}} (A,B) = L(1) + L(2) + L(3) + L(4)$$



UPGMA = Unweighted Pair Group Methods using Arithmetic means.



The ultrametric condition implies the molecular clock

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Clustering methods work well, if sequences evolve according to a molecular clock



or equivalently: if the ultrametric inequality is holds:

$$d(A,B) \le \max\left\{d(A,C), d(B,C)\right\}$$

for each triple (A,B,C)

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Theorem: Four-Point-Condition

A distance matrix $(d_{i,j})_{i,j=1...,n}$ is representable as a tree, if and only if $d(u,v) + d(x,z) \le \max\{d(u,x) + d(v,z), d(u,z) + d(v,x)\}$ for all $u, v, x, z \in \{1, 2, ..., n\}$



Theorem: Four-Point-Condition

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1. begin with star tree:



2. compute for each pair (1,2) the **net-divergence** $\frac{1}{2(N-2)} \sum_{k=3}^{N} (D_{1k} + D_{2k}) + \frac{1}{2} D_{12} + \frac{1}{N-2} \sum_{3 \le i < j} D_{ij}.$ (1)

3. take the pair (A,B) that minimizes Eq. (1)



4. cluster (A,B) and define an interior node W



5. compute branch lengths for the external edges:

$$L(A,W) = \frac{1}{2} \left(D(A,B) + \frac{1}{m-2} \sum_{k=1}^{m} D(A,k) - D(B,k) \right)$$
$$L(B,W) = \frac{D(A,B)}{2} - L(A,W)$$



6. compute distance W to the remaining m-2 leaves:

$$D(W,k) = \frac{1}{2} (D(A,k) + D(B,k) - D(A,B))$$

7. continue with step 1 with the reduced set of leaves

The Neighbour Joining Algorithm









Find a tree $\boldsymbol{\tau}$ that minimizes

$$S(\tau) = \sum_{i,k} \left(\rho(i,k) - D(i,k) \right)^2$$

where $\rho(i,k)$ is the length of the unique path connecting leaves *i* and *k* in the tree.

Distance Correction







$$obs(d) = \frac{3}{4} - \frac{3}{4} Exp[-4d/3]$$

obs(d) can be estimated from the number of observed different pairs of positions n_1 between two aligned sequences of length *I*. Solving

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$$\frac{n_1}{l} = \frac{3}{4} - \frac{3}{4} Exp[-4d/3]$$

leads to Jukes Cantor correction:

$$d = -\frac{3}{4} Log \left[1 - \frac{4}{3} \frac{n_1}{l} \right]$$

Distance Correction





The Problem: Different alignments, different trees



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The Problem: Different alignments, different trees



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The alignment strategy may have more impact on the reconstructed tree than does the type of tree building method. Morrison and Ellis (1997) Mol. Biol. Evol. 14:428-441 Focussing on stable parts of the alignment



Gblocks (Castresana (2000) Mol. Biol. Evol. 17:540-552 Objective: Define a set of conserved blocks from an alignment to be used in phylogeny reconstuction

Approach:

1) Classification of Columns

- > non-conserved : <n/2 + 1 identical residues, or a gap
- > conserved : $\geq n/2 + 1$ and < 85% identical residues
- highly conserved :>85% identical residues

2) discard contiguous stretches of non-conserved positions (default I = 8)
3) from remaining blocks: remove flanking positions until blocks begin and end with highly conserved positions, i.e. selected blocks are anchored by positions that can be aligned with high confidence

4) discard blocks with I < 15

5) remove all positions with gaps together with adjacent positions until a

conserved position is reached

6) discard blocks with I < 10

Focussing on stable parts of the alignment

	70	80	90	100	110	120
nad2 nardo				CENCIC DV	NEACT MUEL NU	
nad3_acaca	E-DTUSK NV VVI		TOTMVT O WA	ISTABLS-DV	AP OLIVE LAV	
nad3_allma	G-DAROKEDVSEVIJ		TEVVETLEFAS	SVTHNVS_LL	GOTTTTFLV	TAPT
nad3_apec	N-SARVPESERE	VATLELIED	DOTAL DOPLPI	FSVFFHP-	-IHTPI	TVCLT
nad3 arath	G-DARSREDIREYIN	VSILELIPD	DEVTERSWAY	VPPNKID-LF	GFOSMMAFLE	TTOT
nad3 balca	A-SARLPESLREFL	VAILFLEFD	LPTALLLPFP/	ALSARDPOL	SFTLAFLILL	TIGLI
nad3 chocr	D-DARATEDIREY	VAILFLIFD	LEISFLEPWSI	LVLGEIS-ĨI	GF <mark>@</mark> SMIVFLV	TIGFI
nad3_drome	S-SSRLPESLREFL1	IT <mark>II<mark>FLIF</mark>D</mark>	VE <mark>T</mark> ALIL <mark>P</mark> MI	IIMKYSNIMI	WTITSIIFIL	LICLY
nad3_human	S-PARVPRSMK FL	VAI TELLED	LPI ALLLPI	VALQTTNLPL	MVMSSLLLII	ALSLA
nad3_ktun	S-S <mark>AR</mark> LP <mark>R</mark> SM <mark>RF</mark> FL]	ITVV <mark>EL</mark> VED	VETVLLPYLI	FSSGWSIDVF	SLVGSMMILV	IICVL
nad3_lter	S-TARIPESTREFLI	LAIIFIVED	IDTVLLMPLP?	TILHTSDVFT	TVTTSVLFLM	LICLI
nad3_marpo	D-D-RSREDIESYL	VSIL:II:D	ID VTFIDE WAY	VSLNKIG-LF	GF <mark>O</mark> SMMVFLF	VELOPV
nad3_metse	G-TPGRPISSI SFL	IGUINEILIED	IN SFINE WC	VVCNQVF-PF	GY <mark>M</mark> TMIVFLAV	LCLV
nad3_picca	R-QSTTYSIKEI		INCLUSIL YT	LSMYNTN-IY	GLFILLYFLLP	11GF1
nad3_podan	LGON TO GVK F11		LETEPFA	VSEIVNN-II	GLIILLGFITII	
nad3_prowi		VALLEITED	IN VALUE WAL	ALCOVE_TE	GP MSMMLFLF	
nau5_recam	D-D-GALDVILLI		WAVAP IN TWA	VALSDVI-IF		
	130	140	150			
nad3 parde	VENKKGATENA					
nad3 acaca	YEWOKGALEWD					
nad3_allma	YEFVSGAITDSF					
nad3 apec	FENVOCGUDNAE					
nad3_arath	YEWKR CASDRE					
nad3_balca	YENME <mark>C</mark> GLENAE					
nad3_chocr	YEWYK <mark>GAL</mark> EWE					
nad3_drome	H <mark>BNNQCMLNNSN</mark>					
nad3_human	YENLOKGLOWTE					
nad3_ktun	HENSEESLENFSSSI	N				
nad3_lter	HERKECSIDOSS					
nad3_marpo						
nad3_metse		IDNUKCHEG	VUVVINUT			
nad3_picca	THINTKOI INTEMPT	UT MM M D D N V	CONTRVICET			
nad3_pouan	VERREAT	VIIMIRENI	SSTIETLGKI			
nad3_prowi	VISOKKCAT DOF					
naa5_reeam						

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A typical variant: Weighted Sum of Pairs



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A typical variant: Weighted Sum of Pairs



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MFP



Weighting of sequences: one variant





A typical variant: Weighted Sum of Pairs



$$\sigma_{wsop}(\alpha) = \sum_{i < j} \omega_i \omega_j S(\alpha_i, \alpha_j)$$


Orthologous Sequences, Please!!





Arguments for orthology assumption:

- a sequence tree that is congruent to the species tree
- conservation of genomic position
- sequence similarity (typically, reciprocal best blast hit)
- similarity of function

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