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Graphs present a key concept of discrete mathematics. Informally, a graph consists of

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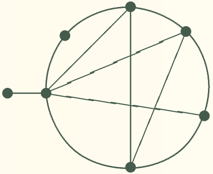
- and *edges* ("arcs") between pairs of vertices.



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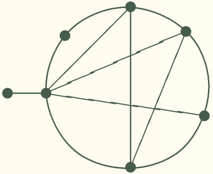

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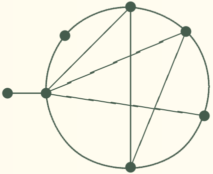

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Brief outline of this lecture

- What is a graph: definition of a graph and its basic terms, examples and trivial classes of graphs.
- Vertex degrees, degree sequence of a graph (score).
- Subgraphs and isomorphism, recognizing (non-)isomorphic graphs.
- Directed graphs and multigraphs.

1.1 Defining a graph

Definition 1.1. Graph (actually, a *simple undirected* graph)
is a pair $G = (V, E)$ where V is the *vertex* set and E is the *edge* set – a subset of pairs of vertices.

Notation: An edge between vertices u and v is denoted by $\{u, v\}$, or shortly uv.
The vertices joined by an edge are *adjacent* or *neighbours*.
The vertex set of a (known) graph G is referred to as $V(G)$, the edge set as $E(G)$.

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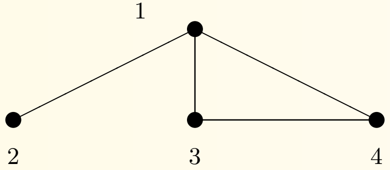
Remark: How can we describe a graph?
Either by listing the vertices and edges, or with a nice picture...

$$V = \{1, 2, 3, 4\}, \quad E = \Big\{\{1, 2\}, \{1, 3\}, \{1, 4\}, \{3, 4\}\Big\}$$



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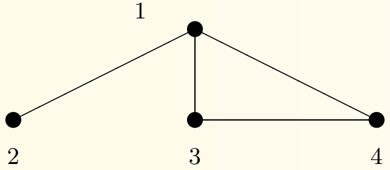


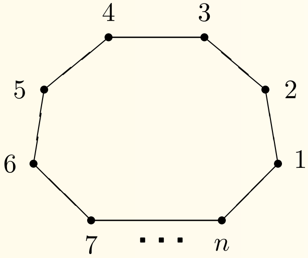
Which one do you like more?

Basic graph classes

It is a custom to refer to some basic special graphs by descriptive names such as the following.

The cycle of length n has $n \geq 3$ vertices joined in a cyclic fashion with n edges:

C_n



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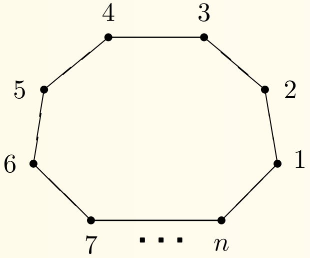
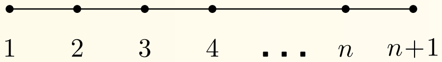


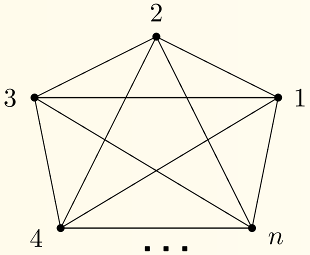
The path of length n has $n+1$ vertices joined consecutively with n edges:

P_n



The complete graph on $n \geq 1$ vertices has n vertices, all pairs forming edges (i.e. $\binom{n}{2}$ edges altogether):

$\boldsymbol{K_n}$



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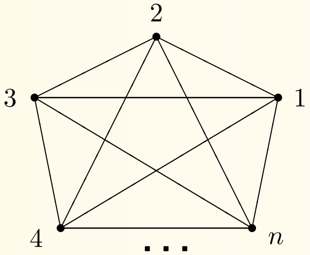
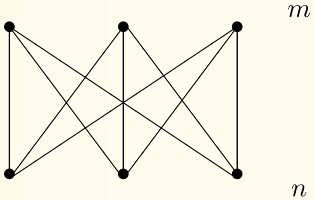


The complete bipartite graph on $m \geq 1$ plus $n \geq 1$ vertices has $m + n$ vertices in two parts, and the edges join all the $m \cdot n$ pairs coming from different parts:

$K_{m,n}$



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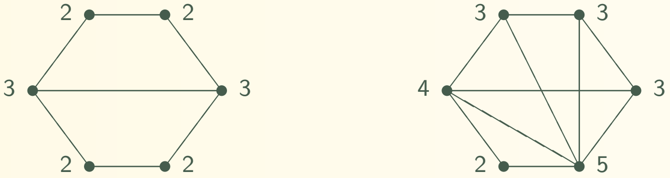

Definition: A graph is *d-regular* if all its vertices have the same degree d.

Notation: The *maximum* degree in a gr. G is denoted by $\Delta(G)$ and *minimum* by $\delta(G)$.

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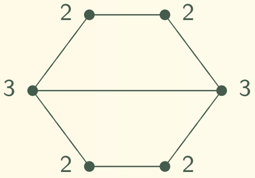

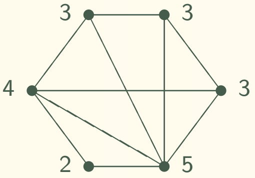

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Proof. When summing the degrees, every edge has two ends and hence it is counted twice. □

The degree sequence

Definition: The *degree sequence* (called also the score) of a graph G is the collection of degrees of the vertices of G, written in a sequence of natural numbers which is (usually) sorted as nondecreasing or nonincreasing.

In abstract graphs, their vertices usually have no names, and so we have to sort their degree sequence somehow. The particular custom is not important.

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Theorem 1.4. *Let $d_1 \le d_2 \le \cdots \le d_n$ be a sequence of natural numbers. There exists a simple graph on n vertices having a degree sequence*

$$d_1,\ d_2, \ldots,\ d_n$$

if, and only if, there exists a simple graph on $n-1$ vertices having the degree sequence

$$d_1,\ d_2, \ldots,\ d_{n-d_n-1},\ \ d_{n-d_n}-1, \ldots,\ d_{n-2}-1,\ d_{n-1}-1\,.$$

Example 1.5. *Is there a simple graph with degree sequence*

$(1, 1, 1, 2, 3, 4)$?

Using Theorem 1.4 we modify the sequence to $(1, 0, 0, 1, 2)$,

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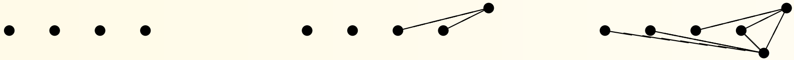
How can we construct such a graph? See. . .



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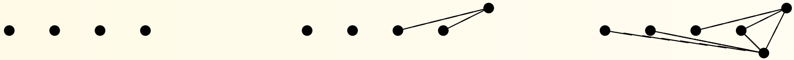

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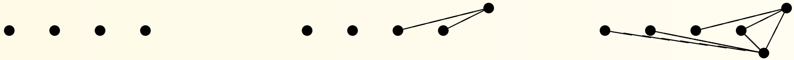

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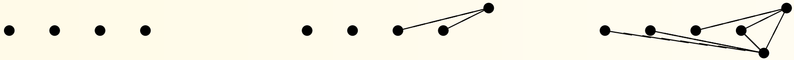

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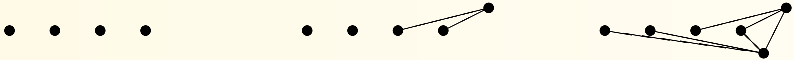

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Since the degrees of a graph cannot be negative, a graph with the last sequence does not exist, and neither does the original one!

□

1.3 Subgraphs and Isomorphism

Definition: A *subgraph* of a graph G is any graph H on a subset of vertices $V(H) \subseteq V(G)$, such that the edges of H form a subset of the edges of G and have both ends in $V(H)$.

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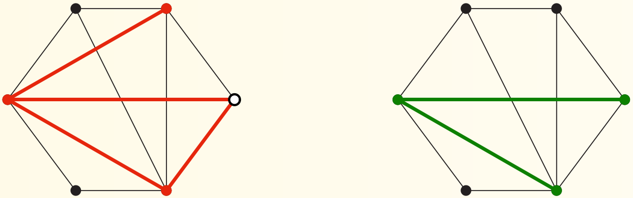

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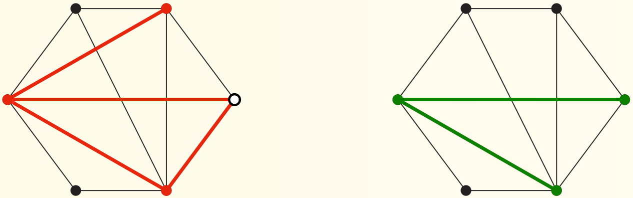

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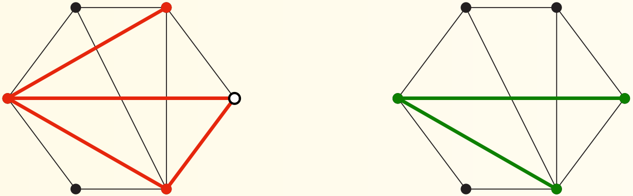

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Definition: An *induced subgraph* is a subgraph $H \subseteq G$ such that all edges of G between pairs of vertices from $V(H)$ are included in H.

Definition 1.6. An isomorphism between graphs G and H is a bijective (*one to one*) mapping $f : V(G) \to V(H)$ such that, for every pair of vertices $u, v \in V(G)$, it holds that $\{u, v\}$ is an edge of G if and only if $\{f(u), f(v)\}$ is an edge of H.

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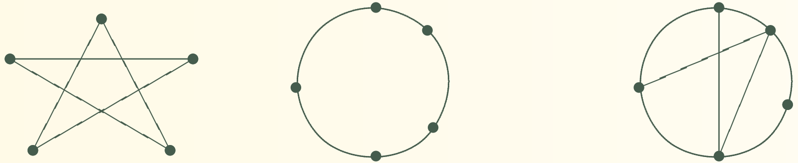

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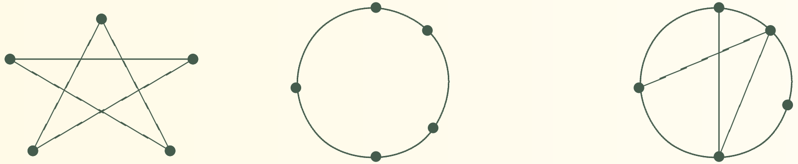

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Fact: If a bijection f is an isomorphism, then f must map the vertices of same degrees, i.e. $d_G(v) = d_H(f(v))$. The converse is not sufficient, however!

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We shall first compare the numbers of vertices and of edges. (Agree.)

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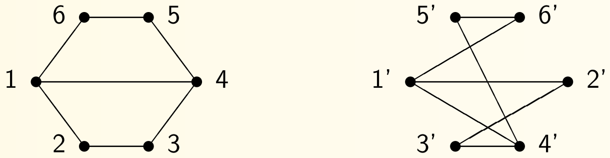
In this particular case, it helps to notice that the two degree-3 vertices on the left are symmetric, and so we may choose any one of them to map it to the leftmost vertex of the second graph. Numbering the vertices by $1, 2, 3, 4, 5, 6$, we can construct our mapping (already have $1 \to 1'$) as follows, see on the picture below.



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Theorem 1.8. *The relation "to be isomorphic" $\simeq$ is an equivalence on the class of all graphs.*

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The important corollary of this claim is that the class of all graphs is partitioned into the *isomorphism classes*.
Hence when we speak about a "graph", we (usually) actually mean its whole isomorphism class. A particular presentation of the graph does not matter.

Additional graph notation

Notation: Consider an arbitrary graph G.

- A subgraph $H \subseteq G$ isomorphic to a cycle is called a *cycle in G*.
- Specially, a *triangle* in G is a subgraph isomorphic to the cycle of length 3.

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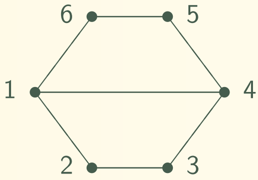
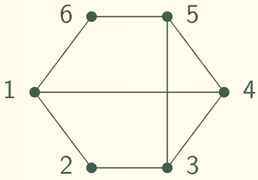
- An induced subgraph $H \subseteq G$ isomorphic to a cycle is called an *induced cycle in G*. Analogously an induced path...

What subgraphs and induced subgraphs can you find in the following graphs?





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Fact: No universal and efficient way of deciding an isomorphism between two graphs is currently known (the problem is not known in P). On the other hand, however, the problem is neither NP-hard to our knowledge, and so a general polynomial algorithm might emerge in the future...

1.4 Directed graphs and Multigraphs

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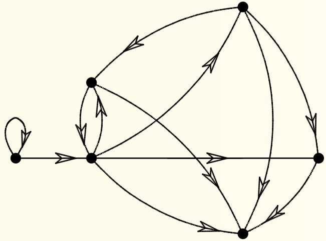

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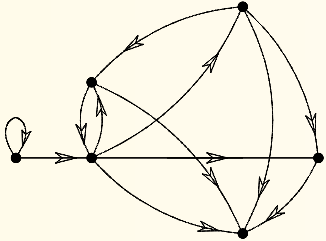

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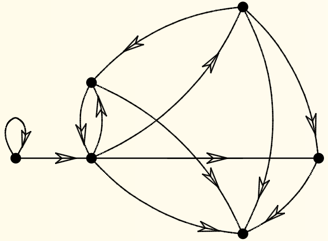

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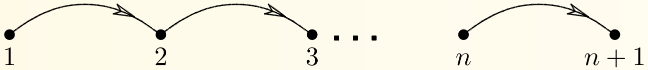
Notation: An edge–arc $e = (u, v)$ in a digraph D has its *tail* in u and the *head* in v, or e joins "from u to v". The opposite arc (v, u) is distinct from (u, v)!

For example, a *directed path* of length $n \geq 0$ looks as follows



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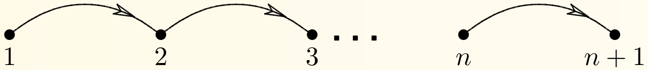

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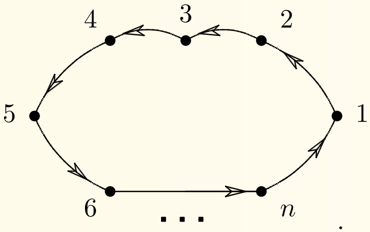

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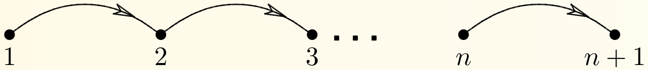

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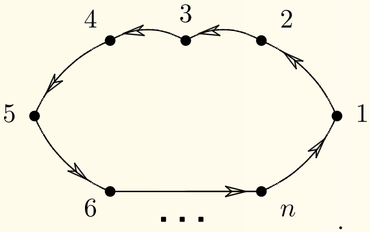

Definition: The number of arcs from u in a digraph D is called the *outdegree* $d_D^+(u)$, and the number of those to u the *indegree* $d_D^-(u)$.

Generalizing even more...

On some other occasions one may want to speak about structures in which more than one edge exist between one pair of vertices, and the edges might have mixed types (undirected or directed, loops).

This leads to so called *incidence model* of a (multi)graph in which edges are elements on their own, along with the vertices; as oposed to our default *adjacency model* where only vertices are considered as core entities.

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Definition: A *mixed multigraph* is a triple $M = (V, F, \varepsilon)$ where $V \cap F = \emptyset$ and $\varepsilon : F \to \binom{V}{2} \cup V \cup (V \times V)$ is an *incidence mapping* of the (multi)edges.

In the definition,

- $\binom{V}{2}$ represents unoriented edges,
- V unoriented loops, and
- $V \times V$ oriented edges and loops.