

# PV181 Laboratory of security and applied cryptography



## Symmetric cryptography

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## Before we start

- Log into your account within IS
- Find and download provided materials
- Look where is the openssl folder
  - Start -> openssl -> version -a

# Goals of Cryptography

- Confidentiality (privacy) - preventing open access
  - **ciphers**
- Authentication:
  1. Entity – identity verification – various (password, MAC, ...)
  2. Data origin – identity of message originator – **MAC**
- Integrity - preventing unauthorized modification
  - **hash functions**
- Non-repudiation - preventing denial of actions
  - **digital signature**

# Crypto primitives

- **Ciphers** – encryption/decryption of data using **key**
  - Symmetric ciphers – **same** key for enc/dec
  - Asymmetric ciphers – **different** key for enc/dec
- **Random number generators (RNGs)**
  - Key generation
- **Hash functions** – “unique” fingerprint of data
- Based on previous: MAC, PBKDF, Digital signature

## Standards

Primitives are defined in various types of standards:

- FIPS PUB 197 – AES block cipher
- RFC1321 – md5 hash function
- NIST SP,...

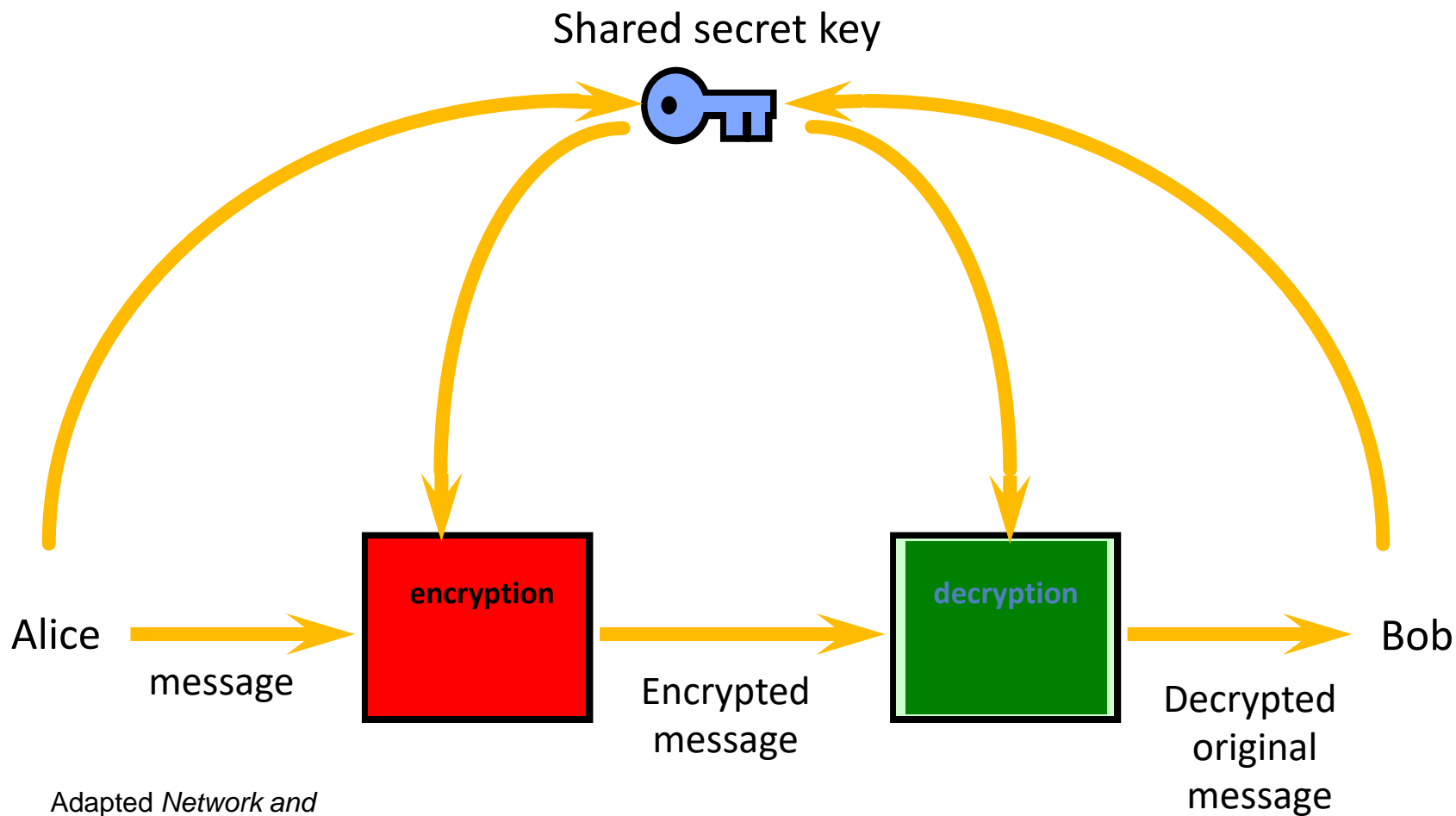
**Test vectors:** defined output to test implementation

- MD5 ("") = **d41d8cd98f00b204e9800998ecf8427e**

# Ciphers: Kerckhoffs' principle

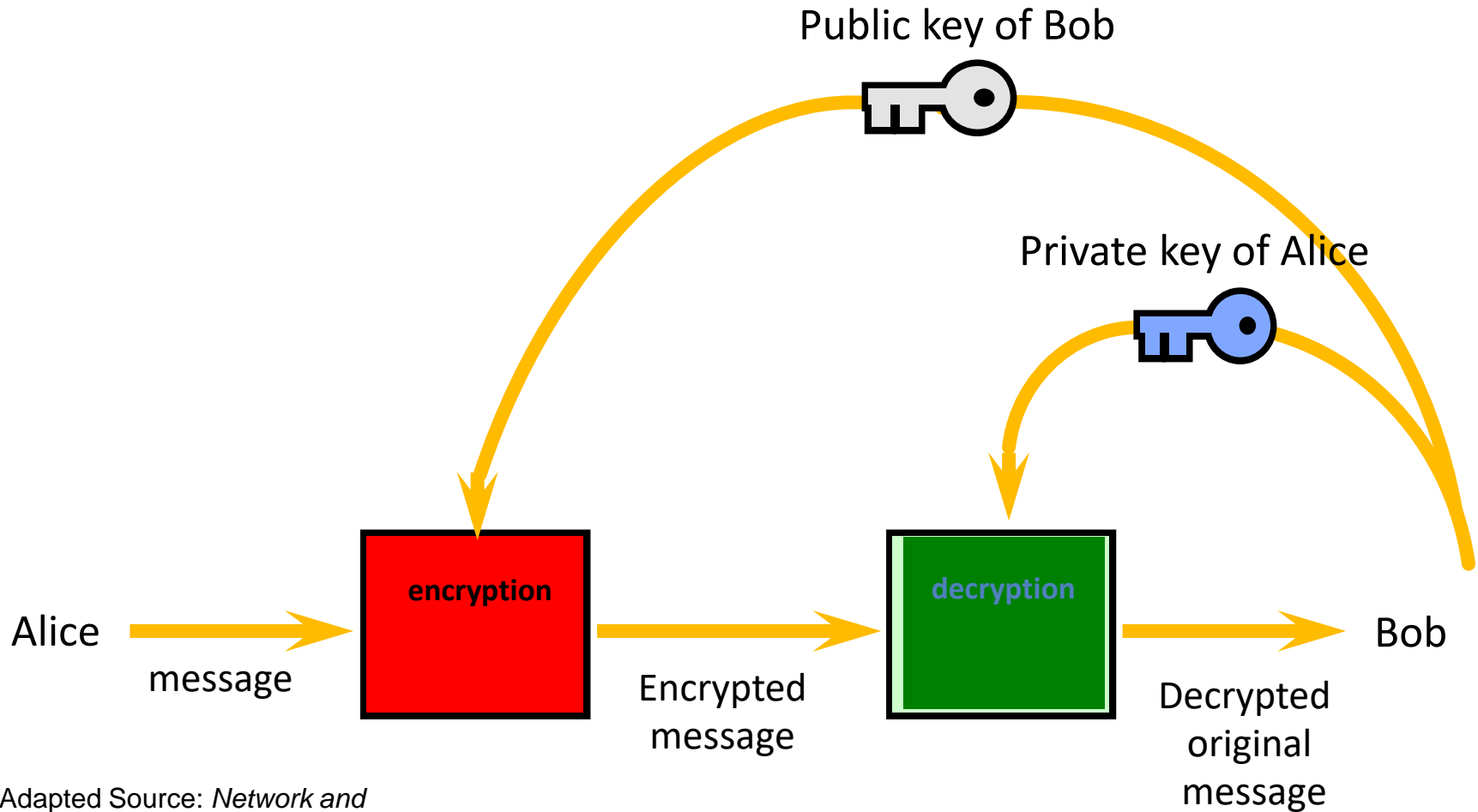
- A cryptosystem should be secure even if everything about the system, except the **key**, is public knowledge.
- I.e. only the **key** should be kept secret, not the algorithm.

# Symmetric cryptosystem



Adapted *Network and  
Internetwork Security* (Stallings)

# Asymmetric cryptosystem

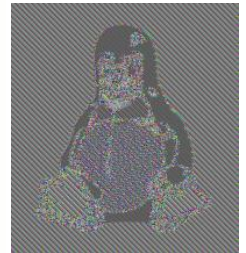


Adapted Source: *Network and Internet Security* (Stallings)



# Block cipher

- Input divided into blocks of fixed size (e.g 256 bits)
  - Padding - message is padded to complete last block
- Different modes of operation:
  - Insecure basic ECB mode – leaks info
  - Secure modes: CBC, OFB, CFB, CTR, ..
- CBC, OFB, CFB need initialization
  - Initialization vector (IV) – must be known



Source: [https://en.wikipedia.org/wiki/Block\\_cipher\\_mode\\_of\\_operation](https://en.wikipedia.org/wiki/Block_cipher_mode_of_operation)

# Block ciphers - padding

*Standard*

ANSI X.923

ISO 10126

PKCS7

ISO/IEC 7816-4

Zero padding

*method*

... | DD DD DD DD DD DD DD DD | DD DD DD DD **00 00 00 04** |

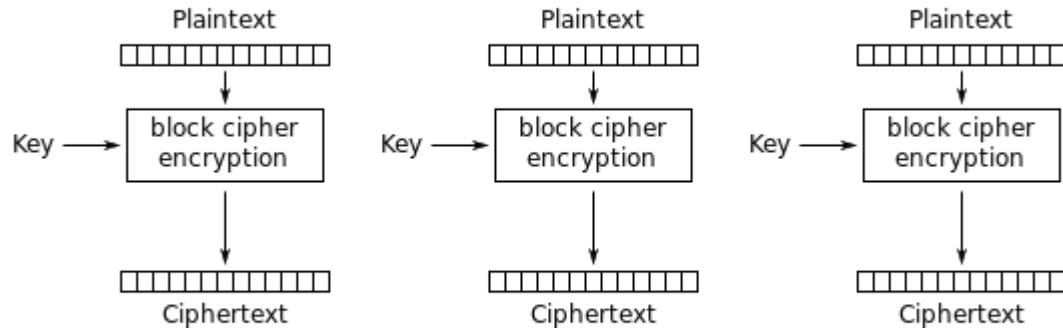
... | DD DD DD DD DD DD DD DD | DD DD DD DD **81 A6 23 04** |

... | DD DD DD DD DD DD DD DD | DD DD DD DD **04 04 04 04** |

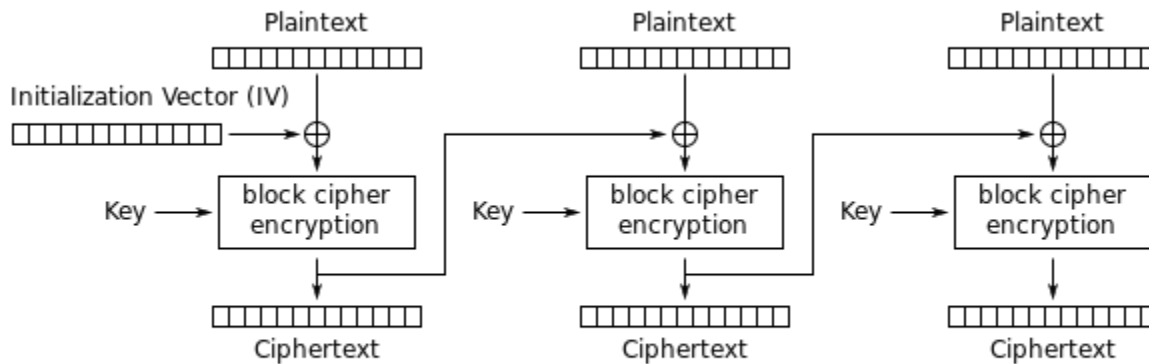
... | DD DD DD DD DD DD DD DD | DD DD DD DD **80 00 00 00** |

... | DD DD DD DD DD DD DD DD | DD DD DD DD **00 00 00 00** |

# Block ciphers: ECB vs CBC mode



Electronic Codebook (ECB) mode encryption



Cipher Block Chaining (CBC) mode encryption

Source: [https://en.wikipedia.org/wiki/Block\\_cipher\\_mode\\_of\\_operation](https://en.wikipedia.org/wiki/Block_cipher_mode_of_operation)

# Random number generators

- Used to generate: keys, IV, ...
  1. Truly RNG - physical process
    - aperiodic, slow
  2. Pseudo RNG (PRNG) – software function
    - deterministic, periodic, fast
    - initialized by **seed** – fully determines random data
- Combination often used:
  - truly RNG used to generate **seed** for PRNG
  - dev/urandom, dev/random in Linux, **Fortuna** scheme

# Hash function

- **Cryptographic** hash function
- Input of arbitrary size
- Output of fixed size:  $n$  bits (e.g. 256 bits).
- Function is not injective (there are “**collisions**”).
- Hash is a compact representative of input (also called imprint, (digital) fingerprint or message digest).
- Hash functions often used to protect integrity. First the has is computed and then only the hash is protected (e.g. digitally signed).

# Hash function properties

- One-way property
  - It is easy to calculate  $\mathbf{h}(\mathbf{x})$  for arbitrary  $\mathbf{x}$ .
  - In a reasonable time it is not possible for the fixed  $\mathbf{y}$  to find  $\mathbf{x}$ , such that  $\mathbf{h}(\mathbf{x}) = \mathbf{y}$ .
- Collision resistance
  - (**weak**): In a reasonable time it is not possible for a given  $\mathbf{x}$  to find  $\mathbf{x}'$  ( $\mathbf{x} \neq \mathbf{x}'$ ) such that  $\mathbf{h}(\mathbf{x}) = \mathbf{h}(\mathbf{x}')$ .
  - (**strong**): In a reasonable time it is not possible to find any  $\mathbf{x}, \mathbf{x}'$  such that  $\mathbf{h}(\mathbf{x}) = \mathbf{h}(\mathbf{x}')$ .

# Cryptographic hash functions

- MD5: output 128 bits
  - Still used although not considered secure at all
  - Broken: efficient algorithm for finding collisions available
  - 128-bit output not considered secure enough
- RIPEMD
  - Output : 128, 160, 256 or 320 bits
  - Less frequently used
- Whirlpool
  - Output: 512 bits
  - Based on AES
  - Recommended by NESSIE project
  - Standardized by ISO

# Secure Hash Algorithm (SHA)

- **SHA-1**
  - NIST standard, collision found in 2016, 160 bits hash
- **SHA-2**
  - function family: SHA-256, SHA-384, SHA-512, SHA-224
  - defined in FIPS 180-2
  - Recommended
- **SHA-3**
  - New standard 2015
  - Keccak sponge function family: SHAKE-128, SHA3-224, ...
  - defined in FIPS 202, used in FIPS-202, SP 800-185
  - Recommended



# Hash functions - examples

- MD5
  - Input: „Autentizace“.
  - Output: 2445b187f4224583037888511d5411c7 .
  - Output 128 bits, written in hexadecimal notation.
  - Input: „Cutentizace“.
  - Output: cd99abbba3306584e90270bf015b36a7.
  - A single bit changed in input → big change in output, so called “Avalanche effect”
- SHA-1
  - Input: „Autentizace“.
  - Output: 647315cd2a6c953cf5c29d36e0ad14e395ed1776
- SHA-256
  - Input: „Autentizace“.
  - Output: a2eb4bc98a5f71a4db02ed4aed7f12c4ead1e7c98323fda8ecbb69282e4df584

# Password protection

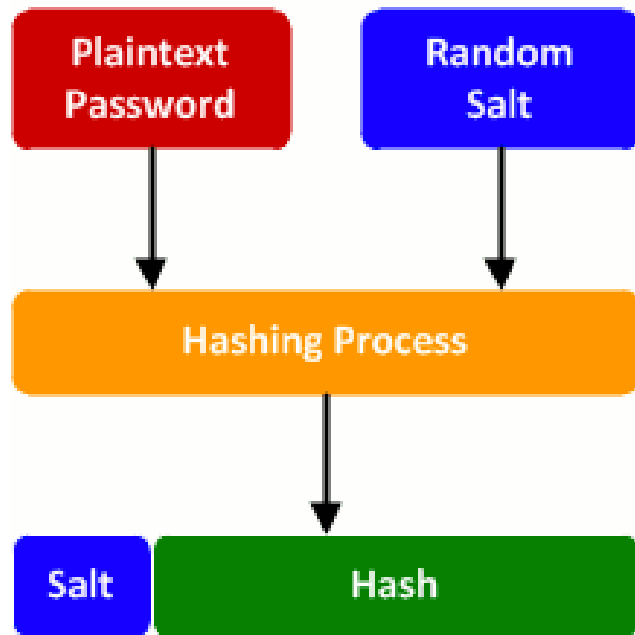
## password hashing & salting

1. Clear password could be stolen:
  - store hash of password  
 $\text{hash} = H(\text{password})$
  - Checking: password is correct if **hash** matches
2. Attack (brute force or dictionary)
  - trying possible passwords “aaa”, “aab” ... “zzz” – N tests
  - N test for single but also for 2,3,... passwords !!!
3. Slow down attack - increase password size:
  - random “**salt**” is added to password,

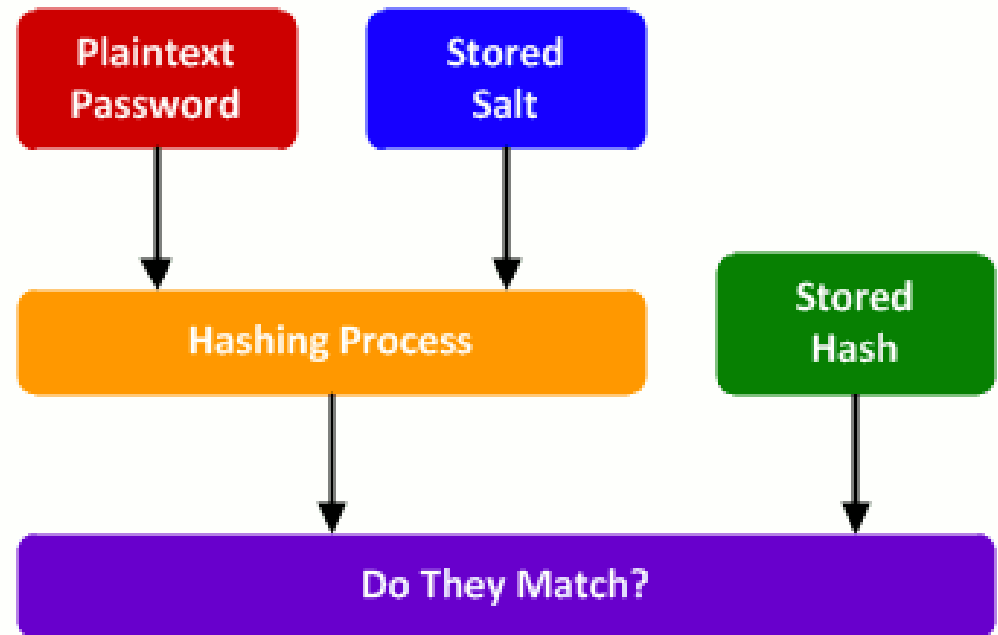
# Password protection

## password hashing & salting

### Password Creation



### Password Verification

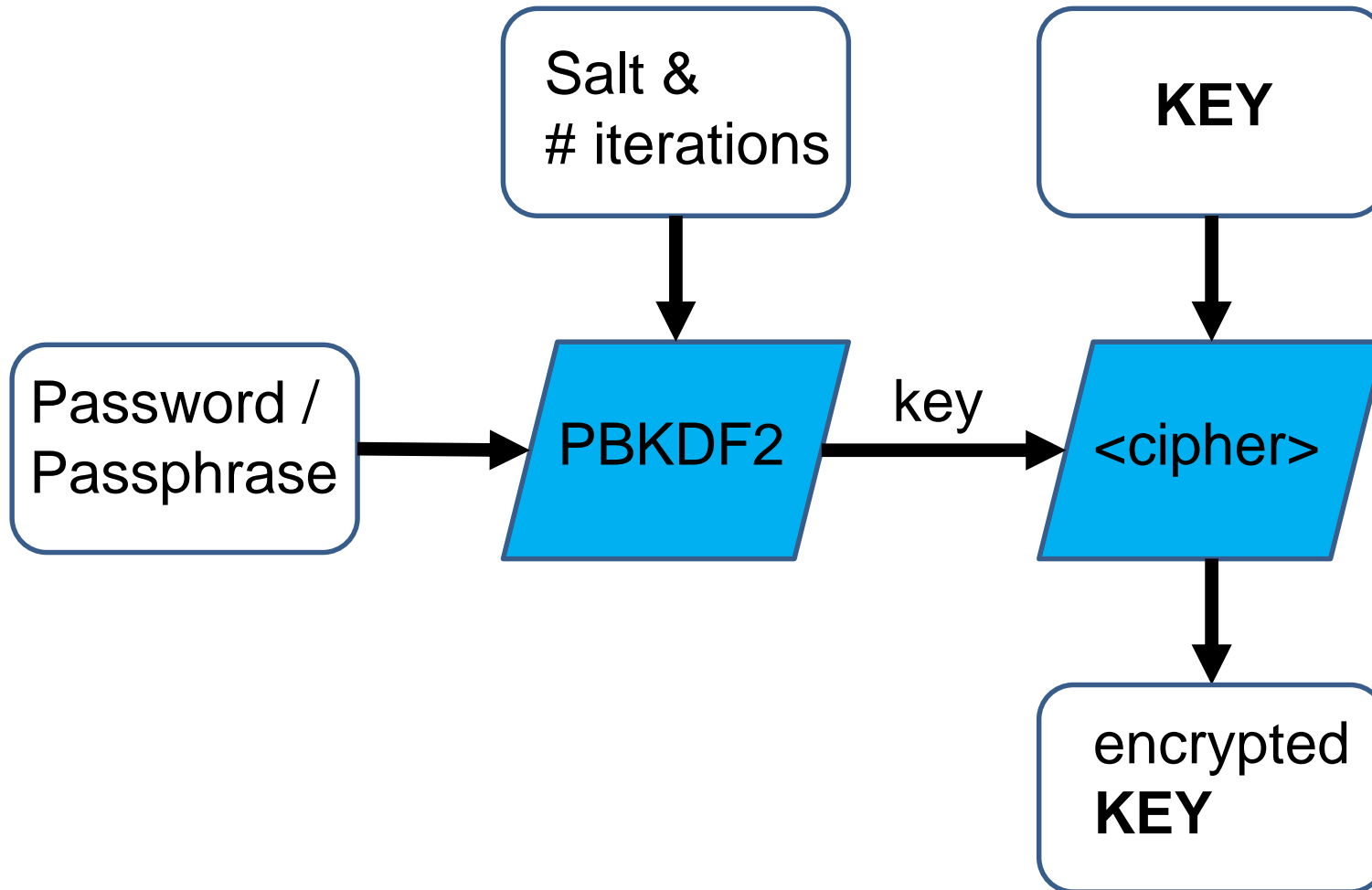


Source: <http://blog.conviso.com.br/worst-and-best-practices-for-secure-password-storage/>

## Key protection

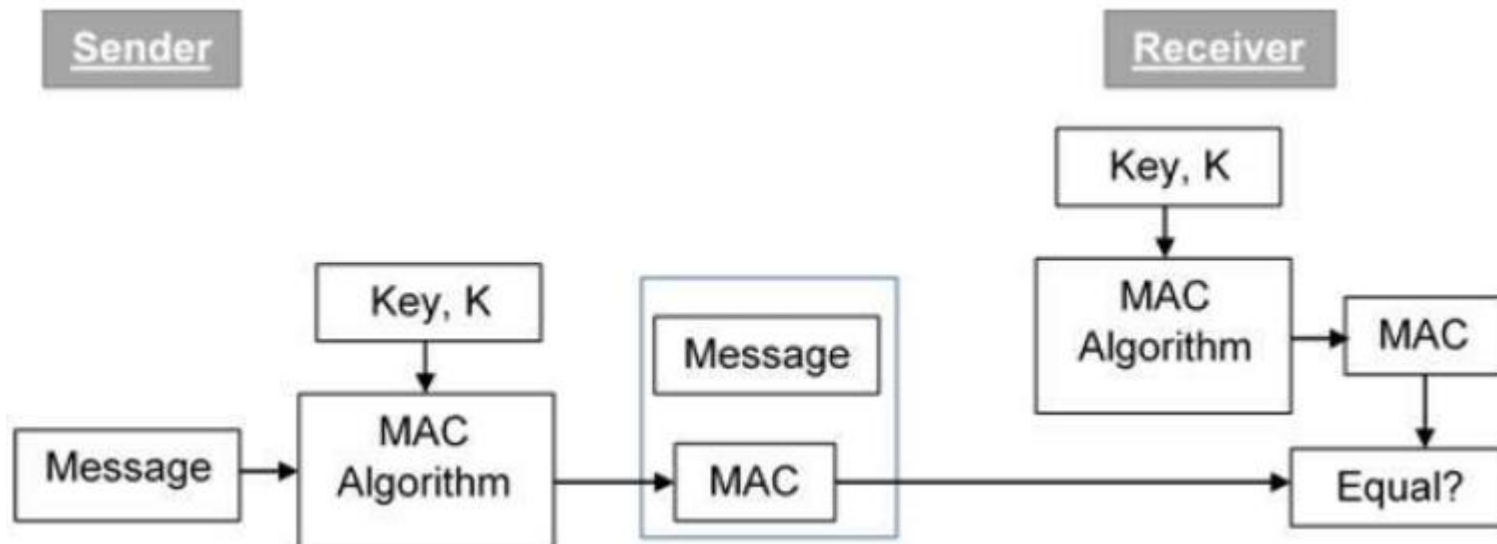
- Encrypt **key**  $K$  (using cipher and other **k**):
  - Key **k** typically derived from password
- Insufficient entropy of passwords
  - E.g. only 17000 guesses for \*\*\*
  - **salt** protects many passwords not single (is stored)
- Password Based Key Derivation Function (PBKDF):
  - 2 types PBKDF2 is newer (PKCS#5 )
  - slow down hash function
  - iterate hash function  $c$  times -  $\text{Key} = H^c(\text{pwd} \mid \text{salt})$ :

# PBKDF2



# Message authentication code (MAC)

- Based on block cipher (MAC) or hash function (HMAC)
- Key + message → algorithm → fixed size block MAC



Source: [https://www.tutorialspoint.com/cryptography/message\\_authentication.htm](https://www.tutorialspoint.com/cryptography/message_authentication.htm)

# RSA: mathematics

- Prime multiplication is simple & Factorization of integers is computationally intensive.
- We choose randomly 2 primes and compute  $n$  and  $\varphi(n)$  :
  - $p, q$
  - $n = p \cdot q$
  - $\varphi(n) = (p-1)(q-1)$ .
- $e$  is chosen such that  $\gcd(e, \varphi(n)) = 1$ .
- We compute  $ed = 1 \pmod{\varphi(n)}$ .
  
- Public key:  $n, e$ .  
Private parameters:  $p, q, d$ .  
Private key:  $d$ .
  
- For RSA with 1024 bit  $n$ , the encrypted message will be 1024 bit long.

## RSA: example

- Intentionally small numbers (such cryptosystem is **not** secure).
- We generate parameters:
  - $p = 17, q = 7,$
  - $n = pq = 119,$
  - $\varphi(n) = 16 \times 6 = 96.$
- Public exponent is selected  $e = 3,5$ , equation is solved (**3 can not be**)  $ed = 5d = 1 \pmod{96}$  to have  $d = 77.$
- The public key:  $(n = 119, e = 5),$   
The private key:  $d = 77.$
- Encryption/decryption:
  - Message  $m = 'C' = 65$
  - Encryption  $m' = 65^5 \pmod{119} = 46.$
  - Decryption  $m = 46^{77} \pmod{119} = 65$



## Links

- SHA1 collision:
  - <https://shattered.io>
- Salting password:
  - <https://crackstation.net/hashing-security.htm>
- OpenSSL
  - Manual: <https://www.openssl.org/docs/man1.0.2/>
  - [https://wiki.openssl.org/index.php/Command\\_Line\\_Uutilities](https://wiki.openssl.org/index.php/Command_Line_Uutilities)
  - <https://www.madboa.com/geek/openssl/>