



Key-value Stores II

Embedded, Distributed, and In-memory Stores

Lecture 5 of *NoSQL Databases* (PA195)

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Agenda



- **Embedded** local storages
 - LevelDB
 - **Local storage** for many systems, Log-structured Merge Tree
- **Distributed** key-value Stores - **representatives**
 - Riak
 - Basics, Riak **Links** & Indexes & Riak Search, Internal **features**
 - Infinispan
 - Basic **features**, example, advanced features, indexing & searching
- Memory **caches**
 - Memcached
- **Serialization**: Protocol Buffers, Apache Thrift

Embedded Stores: Representatives



- **Embedded** local storages

- LevelDB

- **Local storage** for many systems, Log-structured Merge Tree
- C++



- MapDB

- **Java** project, one-man show
- memory-mapped file storage



- RocksDB

- Embeddable persistent key-value store
- **Facebook**
- C++, but also connector from Java





levelDB

LevelDB: Basics

- **Embedded** key-value store
 - Using ideas from Google's **BigTable**
 - **Developers**: Jeffrey Dean and Sanjay Ghemawat from **Google**
- Initial release date: 2011
- License: New BSD Licence
- Language: **C++**
- LevelDB is a **backend** for Google **Chrome**'s IndexDB



LevelDB: Fundamental Features

- Basic **architecture** is a **LSM** Tree (see below)
- **Sorted** by keys
- Arbitrary byte **arrays**

- Basic **operations**: Get(), Put(), Del(), Batch()
- Bi-directional **iterators**



Log-structured Merge Tree

Log-structured Merge Tree (LSM Tree)

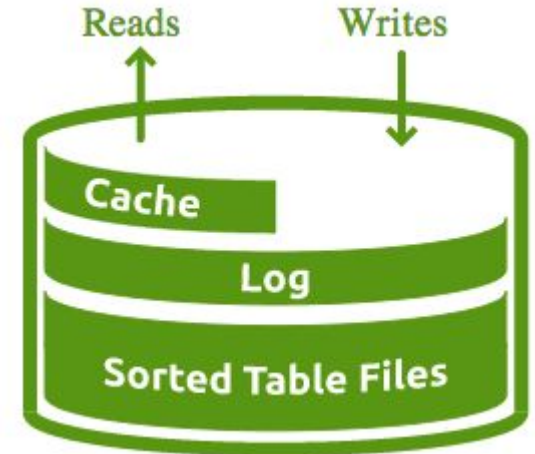
- data **structure** for indexed access to data files
 - can handle high **write frequency**
- **writes** applied to a sorted structure **in memory**
 - regularly **synchronized** to a sorted **disk** storage
- **read ops merge** data from memory & disk

O'Neil, Patrick E.; Cheng, Edward; Gawlick, Dieter; O'Neil, Elizabeth (June 1996). "The log-structured merge-tree (LSM-tree)". Acta Informatica 33 (4): 351–385.



LevelDB: Basic Architecture

- **Writes** go straight into a **log**
- Log is **flushed** to sorted table files (**SSTables**)
- **Reads merge** the log and the SSTable files
- **Cache** speeds up common **reads**

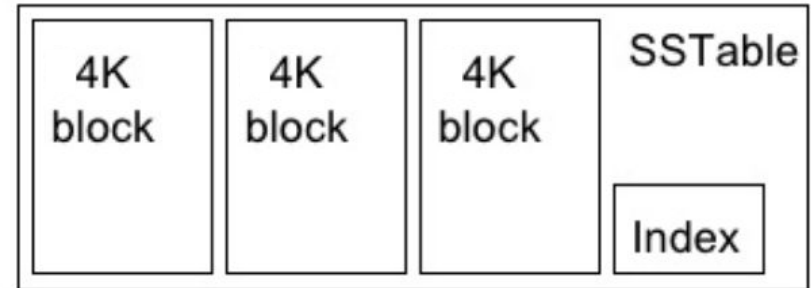




Basic Storage: SSTable Files

Sorted String Table (SSTable) Files:

- Limited to ~2MB each
- Divided into 4K blocks
- Final block is an index
- Bloom filter used for lookups





Levels in LevelDB

Log: Max size of 4MB then **flushed** into a set of **Level 0** SSTables

Level 0: Max of 4 SST files then **the files** compacted into **Level 1**

Level 1: Max total size of 10MB **then the files** compacted **into L2**

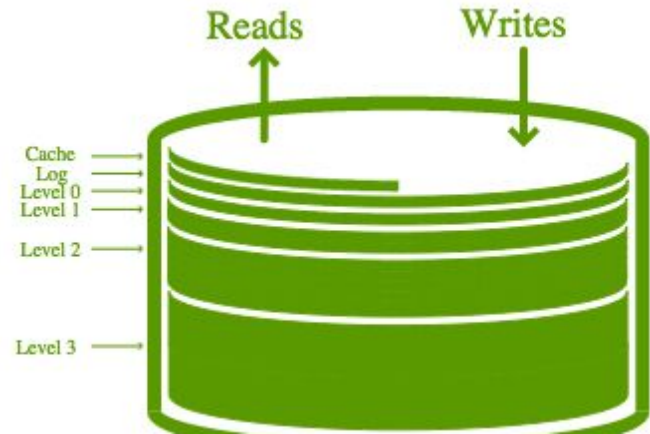
Level 2: Max total size of 100MB then the file **compacted** into L3

Level 3+: Max total size of 10x **previous** level size then the files compacted into **next level**.

0 → 4 SST, 1 → 10M, 2 → 100M,

3 → 1G, 4 → 10G, 5 → 100G, 6 → 1T,

7 → 10T, ...





levelDB

LevelDB: Universal Backend

- LevelDB is a **popular backend** storage for many (distributed) database systems
 - Web browser **IndexDB** (in Chrome)
 - Riak, Infinispan
 - LevelUp/LevelDown for **Node.js**
 - etc.

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- Memory caches
 - Memcached
- Serialization: Protocol Buffers, Apache Thrift

Distributed K-V Store: Riak

Riak: Basic Information

- Open source, distributed key-value database
 - Company Basho, first release: 2009
 - OS: Linux, BSD, Mac OS X, Solaris
- Language: Erlang, C, C++, some parts in JavaScript
- Built-in support for MapReduce
- Provides a full-text search engine on the data
 - “Riak search”

Riak: Basic Mission

● Availability

- Riak replicates and retrieves data intelligently so it is **available for read/write** operations, even in failure conditions

● Fault-Tolerance

- You can lose access to many nodes due to **network partition** or hardware failure **without losing** data

● Operational Simplicity


- **Add new machines** to your Riak cluster **easily** without incurring a larger operational burden

● Scalability

- Riak automatically distributes data around the cluster and yields a **near-linear performance increase** as you add capacity

Riak: Basics

Oracle	Riak	namespace of keys
database instance	Riak cluster	
table	bucket	Terminology in RDBMS vs. Riak
row	key-value	
rowid	key	



- Stores keys into **buckets** = a namespace for keys
 - Like tables in a RDBMS, directories in a file system, ...
 - **Bucket** has its own **properties**
 - `n_val` – **replication** factor
 - `allow_mult` – allowing **concurrent updates**
 - ...

Riak: Interaction with the DB

- Default: **HTTP Interface** (Web services)
 - GET (retrieve value), PUT (update), DELETE (delete), ...
 - example:
`http://localhost:8098/buckets/test/keys/mykey`
- **Native Erlang** interface
- **Connectors** from many (not) standard **languages**
 - C, C#, C++ , Clojure, Dart, Go, Groovy, Haskell, Java, JavaScript, Lisp, Perl, PHP, Python, Ruby, Scala, Smalltalk

Riak: Basic Operations

- Using `curl -X method URL -d data`
 - command line **tool** to **communicate** with server (HTTP(S),...)

```
curl -X PUT http://localhost:8098/buckets/authors/keys/David
-d '{"name": "David Novák", "affiliation": "MU"}'
```

```
curl -X GET http://localhost:8098/buckets/authors/keys/David
{"name": "David Novák", "affiliation": "MU"}
```

```
curl -X DELETE
http://localhost:8098/buckets/authors/keys/David
```


Riak: Additional Functionality

- Riak can have several types of local storage
 - typically referred to as **backends**
 - memory, **LevelDB**, etc.
- Riak has **additional** functions to work with values
 - Riak **links**
 - **Indexes**
 - Riak **search**

Riak: Links

- A way to create **relationships** between objects
 - Like **foreign keys** in RDBMS or **associations** in UML
- Attached to objects via **HTTP header** “Link”
- Add a **book** and link to its **author**:

```
curl -X PUT http://localhost:8098/buckets/books/keys/NoSQL
-d '{"title": "Big Data a NoSQL databáze", "year": "2015"}'
-H 'Link: </buckets/authors/keys/David>; riaktag="wrote"'
```

Riak: Link Walking

- Locate a **key** and then **continue by link(s)**
 - target specification: `/bucket,linktype,[0/1]`
- Find the **authors** who **wrote** book NoSQL

```
curl -i http://localhost:8098  
/buckets/books/keys/NoSQL/authors,wrote,1
```

- Restrict to bucket **authors**
- Restrict to tag **wrote**
- *1 = include this step to the result*

Riak: Indexes

- **Secondary** indexes on the values
 - **Search** key-value pairs based on the **content**
- Indexes kept **locally** on every virtual **node**
- **Types** of indexes:
 1. **integer** index (search by **value** or **interval** of values)
 2. **binary** index (search by any type of **value**)
 3. **fulltext** index (Riak search)

Riak: Indexes (2)

- Indexes **cannot** be managed **automatically**
 - Because there is **no schema** on the values
- When inserting a value, one **can use** index
 - In HTTP API, use special HTTP **headers**

```
curl -X PUT http://localhost:8098/buckets/authors/keys/David
-H 'x-riak-index-surname_bin: Novak'
-H 'x-riak-index-phone_int: 5062'
-d '{"name": "David", surname "Novák", "phone ext": 5062 }'
```

Riak Search: Fulltext via Solr

- Riak provides a distributed, **full-text search** engine
 - Implemented using Solr (Lucene)
 - Inserted **values** are indexed automatically
 - ...and **then search** the data by “terms”
- Key features:
 - Different **parsers** for different mime types
 - JSON, XML, plain text, ...
 - **Exact match** queries: “Bus”
 - **Wildcards**: “Bus*”, “Bus?”
 - **Prefix matching**, proximity searches, range queries...

Java Client

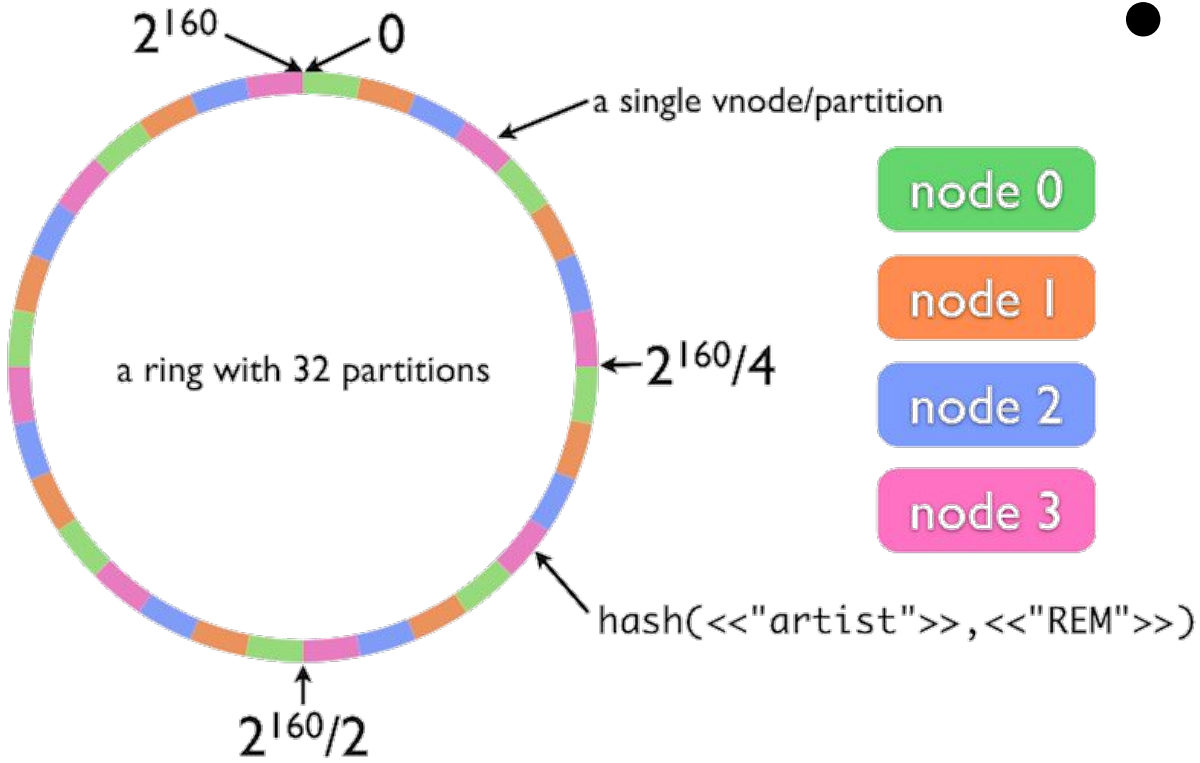
- **Java** Library for communication with Riak
 - Uses Protocol Buffers for communication
 - A way to (de)serialize objects to text (see below)

```
RiakClient client = RiakClient.newClient("168.0.0.1");
Namespace bucket = new Namespace("authors");
Location location = new Location(bucket, "David");
FetchValue fv = new FetchValue.Builder(location).build();
FetchValue.Response response = client.execute(fv);
String obj = response.getValue(String.class);
```

Riak: Internal Features

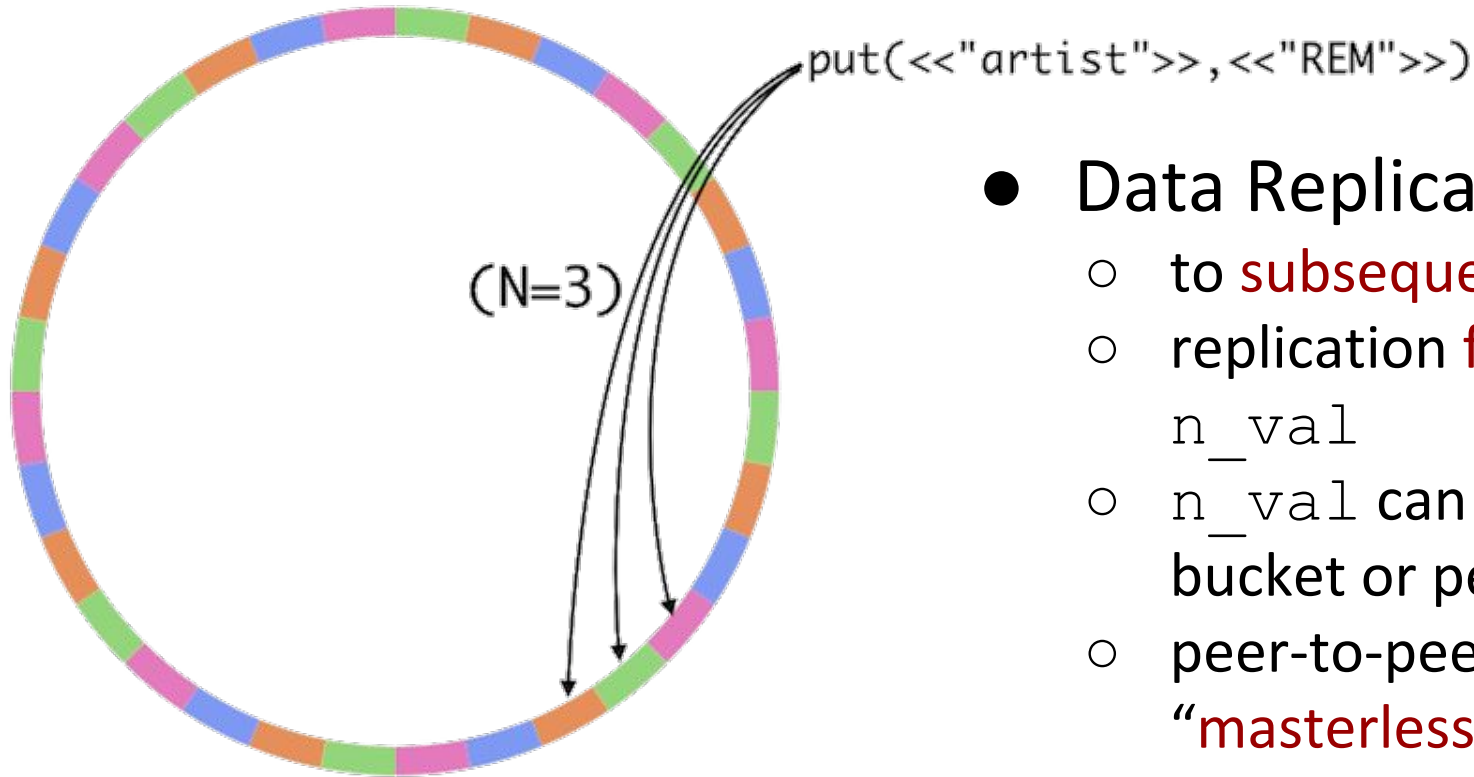
- Let us have a look **behind the scene** of Riak
 - Consistent **hashing**
 - and virtual nodes
 - Peer-to-peer (masterless) data **replication**
 - Read/Write **Quorums**
 - Hinted **handoffs**
 - High availability
 - **Vector clocks**
 - Riak siblings
 - **Gossip** protocol
 - Query processing
 - Riak Enterprise

Consistent Hashing



- Data **Partitioning**
 - consistent hashing into $[0, 2^{160}]$
 - **data balancing** achieved by virtual nodes (vnode)

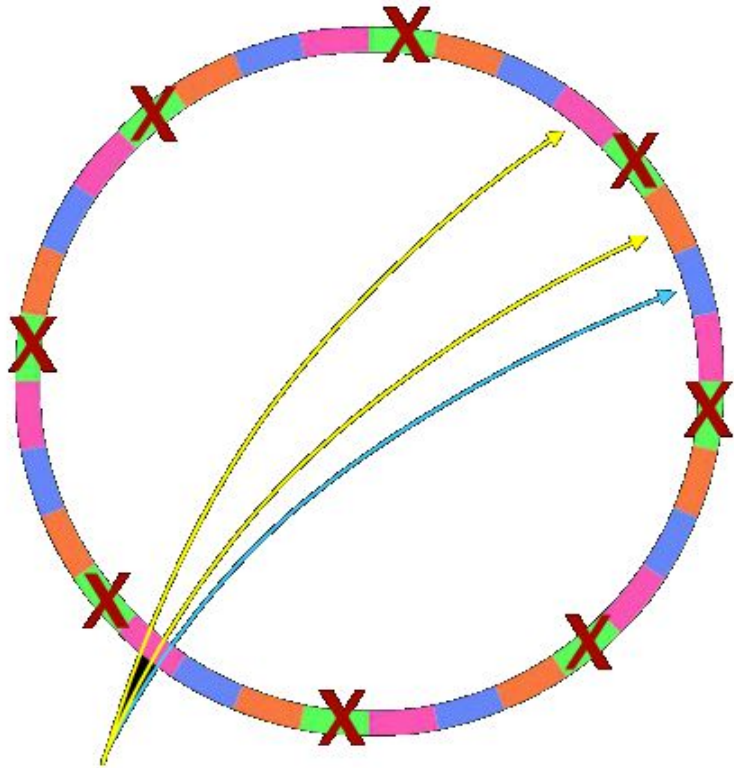
P2P Replication



- Data Replication

- to subsequent nodes
- replication factor `n_val`
- `n_val` can be set per bucket or per object
- peer-to-peer “masterless” replication

Hinted Handoffs



- Goal: High **availability**
- Hinted handoff
 1. In **case** of node **failure**
 2. **Neighboring** nodes temporarily **take over** storage operations
 3. When the failed node **returns**, the **updates** received by the neighboring nodes are **handed off** to it

```
put(<<"artist">>, <<"REM">>)
```

Vector Clocks

- **Any node** is able to receive any **request**
 - We need to know **which version** of a value is **current**
- When a **value** stored, it is **tagged** with a *vector clock*

```
curl http://localhost:8098/raw/plans/dinner  
-X PUT --data "Wednesday"
```

```
curl -i http://localhost:8098/raw/plans/dinner  
HTTP/1.1 200 OK  
X-Riak-Vclock: a85hYGBgzGDKBVISrLnh3BlMiYx5rAzLJpw7wpcFAA==  
Content-Type: text/plain  
Content-Length: 9
```

Wednesday

Vector Clocks (2)

- For each update, Riak **can determine**:
 - Whether one object is a **direct descendant** of the other
 - Whether the objects are descendants of a **common parent**
 - Whether the objects are **unrelated** in recent heritage
- If the objects are unrelated then Riak can:
 - **Auto-repair** data
 - Provide the data to the **user to decide**

```
curl -X PUT -H "X-Riak-ClientId: Ben"  
-H "X-Riak-Vclock:  
a85hYGBgzGDKBVISrLnh3BlMiYx5rAzLJpw7wpcFAA=="  
http://localhost:8098/raw/plans/dinner --data "Tuesday"
```

Riak: Siblings

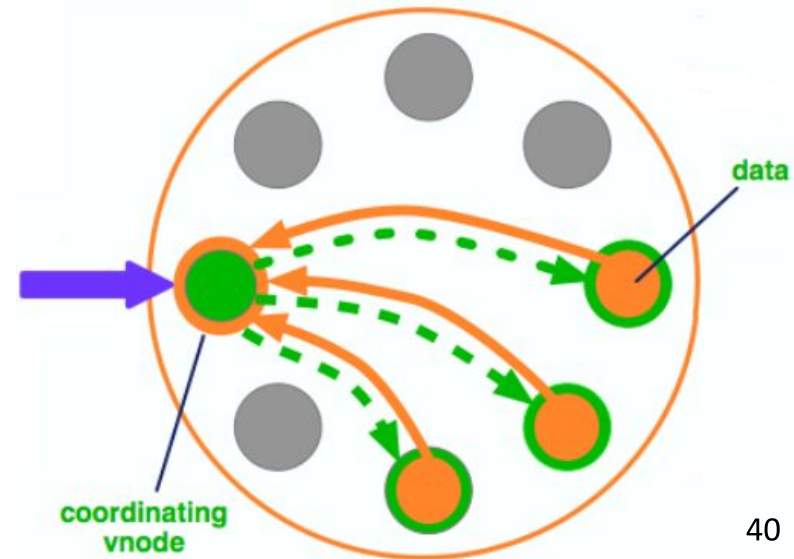
- **Siblings** of objects are **created** in case of:
 - **Concurrent writes** – two **writes** occur simultaneously with **same** vector clock value
 - **Stale vector clock** – **stale** v. clock **value** provided by client
 - **Missing vector clock** – write without a vector clock
- When **retrieving** an object we can:
 - Retrieve **all siblings**
 - **Resolve** the inconsistency

Gossip Protocol

- Gossip protocol
 - To share and communicate **ring state** and bucket properties **around the cluster**
- Each node **periodically** sends **its** current **view** of the ring state
 - To a **randomly-selected** peer

Riak: Request Sharing

- Each node can be a **coordinating vnode** = node responsible for a request
 - Finds the vnode for the key **according to hash**
 - Finds vnodes where **other replicas** are stored – next N-1 nodes
 - Sends a **request to all** vnodes
 - **Waits until** enough requests **returned** the data
 - To fulfill the read/write quorum
 - Returns the result to the **client**



Riak Enterprise

- **Commercial extension** of Riak
- Adds support for:
 - **Multi-datacenter replication**
 - Using more clusters and replication between them
 - Real-time replication – incremental synchronization
 - Full-sync replication – entire data set is synchronized
 - **SNMP monitoring**
 - Simple Network Management Protocol
 - **JMX monitoring**
 - Java Management Extensions

Distributed K-V Store: Infinispan

Infinispan



- **Java**-based key-value **store**
- **Originally**: a distributed memory **cache**
 - for Red Hat JBoss platform
- **Now**: fully-fledged **data-grid**
 - uses libraries from JBoss, Hibernate, etc.

Basics



- Developer: **Red Hat, open source** community
 - Originally developed as a memory-based cache for JBoss
- Initial release date: 2009, current version 10.0
- License: Apache version 2
- Language: Java
 - **embedding** to **Java** application OR
 - external service via various **APIs** (REST service, Memcached protocol, Hot Rod)
 - connectors: Groovy, Scala

<http://infinispan.org/>

<http://db-engines.com/en/system/Infinispan>

Infinispan: Hello World



```
public static void main(String args[]){
    Cache<String, Object> store =
        new DefaultCacheManager().getCache();
    store.put("key1", new MyClass("value1"));
    store.put("key2", "value2");

    if (store.containsKey("key1")) {
        Object result = store.get("key2");
        store.removeAsync("key2");
    }
    store.replaceAsync("key2", "value3");
    store.clear();
}
```

Infinispan: Features (1)



- Running in cluster
 - auto-sharding (distribution mode)
 - basically “consistent-hashing” (customizable)
 - fixed number of “segments” (like “vnodes” in Riak)
 - replication - master/slave (primary/backup owners)
 - synchronous (write through), asynchronous (write back/behind)
- Persistence
 - originally only memory-based, now fully configurable
 - file system store, JDBC store, LevelDB, JPA cache store,...
 - JBoss marshalling (serialization) of Java objects

Infinispan: Features (2)



● Cache features

- **eviction**/expiration (remove objects automatically)
 - either when the cache is full (**LRU**)
 - or after some time (**lifespan** of an entry)
- **invalidation** mode
 - a special type of cluster mode
 - when a value changes, other nodes are informed that their data is stale
- **L1 cache**
 - each **node** keeps a **local cache** of key/values retrieved from other nodes

● MapReduce

- full **support** of MapReduce processing
 - very efficient since version 7.0

Concurrent Operations



- **Full transactional** processing
 - Java Transaction API (**JTA**)
 - X/Open Extended Architecture (**X/Open XA**)
 - optimistic vs. pessimistic transactions
 - deadlock detection
 - Two-phase commit protocol (**2PC**)
- **Distributed Execution** Framework
 - executing a “Callable” on “**nodes storing** given set of keys”
 - compatible with standard Java Execution Framework

Concurrent Operations (2)



- Multi-version Concurrency Control (MVCC)
 - a technique to solve **concurrent access** to data
 - **faster** than strict use of r/w locks
 - popular in many (RDBMS) databases
- For transactions, user can choose **isolation levels**:
 - READ_UNCOMMITTED
 - **don't** use **transactions** at all
 - READ_COMMITTED (default)
 - any transaction does **see new value** immediately **after** its **commit**
 - REPEATABLE_READS
 - using MVCC, the **transaction** does **see** the **same values** all the time

Infinispan: Querying



- Additional indexes
 - to provide **search** over stored **values**
 - using **Hibernate Search** technology
 - ...and **Lucene**
- Vice **versa**:
 - **Infinispan** can serve as a distributed **storage** for **Lucene**

Example: Indexing



```
// A class to be indexed is annotated with @Indexed
// then you pick which fields and how to index them
@Indexed
public class Book {
    @Field String title;
    @Field String description;
    @Field @DateBridge(resolution=YEAR) Date publicationYear;
    @IndexedEmbedded Set<> authors = new HashSet<Author>();
}

public class Author {
    @Field String name;
    @Field String surname;
}
```

Example: Searching



Task: Find books on "book on scalable query engines"

```
SearchManager searchManager = Search.getSearchManager(store);  
// create a query via Lucene APIs or using builder  
QueryBuilder qBuilder =  
    searchManager.buildQueryBuilderForClass(Book.class).get();  
Query luceneQ = qBuilder.phrase()  
    .onField("description").andField("title")  
    .sentence("book on scalable query engines").createQuery();  
CacheQuery res = searchManager.getQuery(luceneQ, Book.class);  
// and there are your results!  
List objectList = res.list();
```




Memcached: Basic Info

- **In-memory distributed** key-value store
- Initial release date: **2003**
 - by Brad Fitzpatrick for LiveJournal
- License: New BSD Licence
- Language: **C**
- Used by:
 - LiveJournal, Wikipedia, Flickr, WordPress.com, Craigslist



Memcached: Features

- Memcached
 - store small chunks of **arbitrary** data (strings, objects)
 - keys up to 250 bytes, values up to 1MB
- Typical usage
 - **cache results** of database calls, API calls, or page rendering
- **API** is available for most popular **languages**

Memcached: Architecture



- Client-server architecture
 - Client-side **libraries** to contact the servers
 - Each **client knows all** servers
 - **Servers do not communicate** with each other
- **Static** sharding
 - The **client** computes a **hash(key)** to determine the server
 - Scalable **shared-nothing** architecture across the **servers**

Apache Thrift: Example

```
enum PhoneType {  
    HOME,  
    WORK,  
    MOBILE,  
    OTHER  
}  
  
struct Phone {  
    1: i32 id,  
    2: string number,  
    3: PhoneType type  
}
```



Lecture Summary



- **Key-value** stores are popular for its **simplicity** and **efficiency**
- Most of the real key-value stores provide **additional** functionality to **search** on the **values**
- Besides **distributed** systems, there are local **embedded** stores and in-memory caches
- There are general **frameworks** to provide **serialization** of objects into **binary** data

