

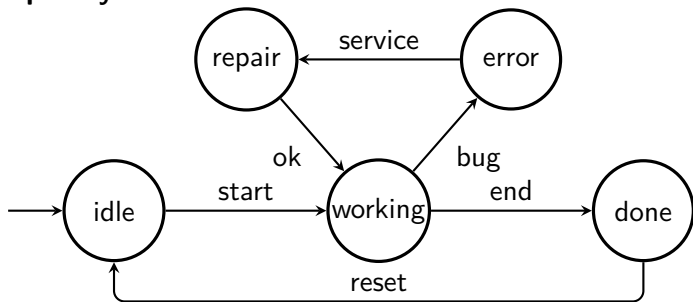
IA169 System Verification and Assurance

Verification of Systems with Probabilities

Vojtěch Řehák

Jiří Barnat

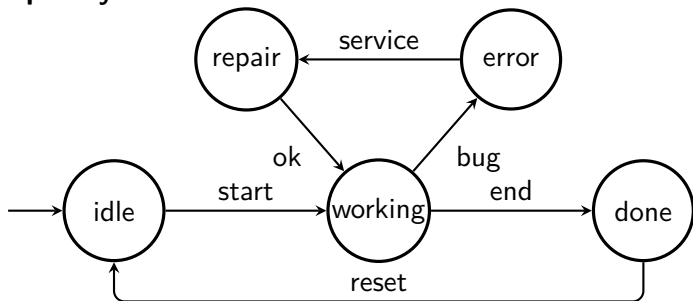
Fail-repair system



What are the properties of the model?

- $G(\text{working} \implies F \text{ done})$
- $G(\text{working} \implies F \text{ error})$
- $FG(\text{working} \vee \text{error} \vee \text{repair})$

Fail-repair system

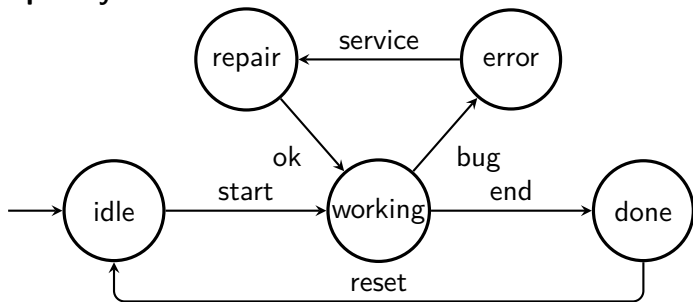


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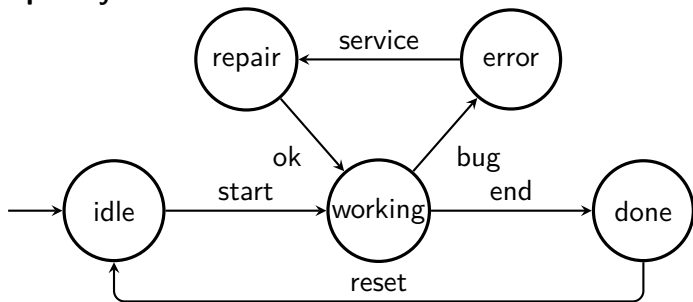
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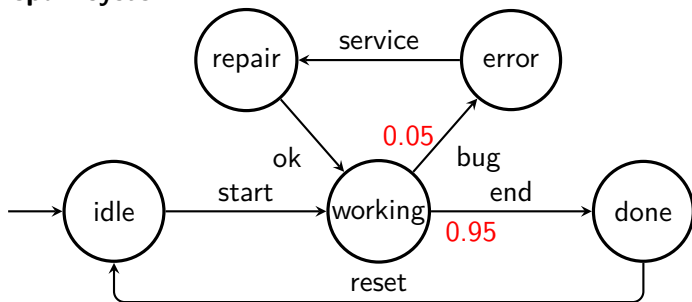
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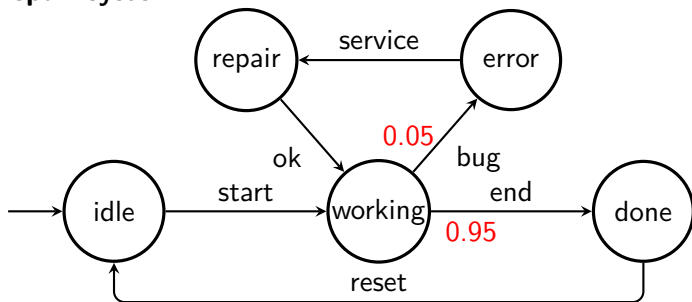
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What is the probability of reaching “done” from “working”

- with no visit of “error”?
- with at most one visit of “error”?
- with arbitrary many visits of “error”?

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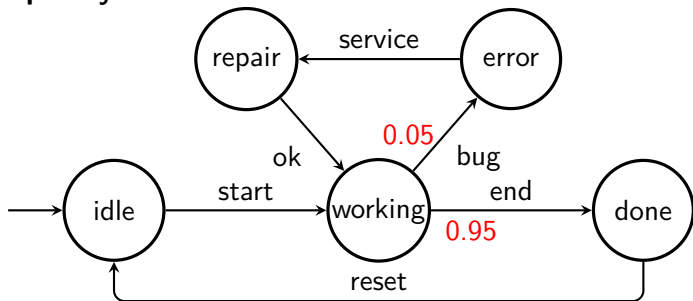


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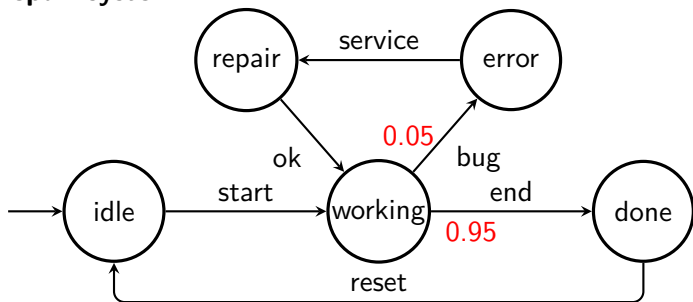
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Discrete-time Markov Chains (DTMC)

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- Standard modeling formalism for probabilistic systems.
- A finite diagram of states and state-changing transitions.
- Each transition is annotated with a probability p ($p \in [0, 1]$).
- The probabilities over transitions from a single state sum to 1. (They form discrete probability distribution.)

Observation

- Markov property (“memoryless structure”) — only the current state determines the successors (the past states are irrelevant).
- Each state has at least one outgoing edge (“no deadlock”).

Task: create DTMC modeling the following scenario

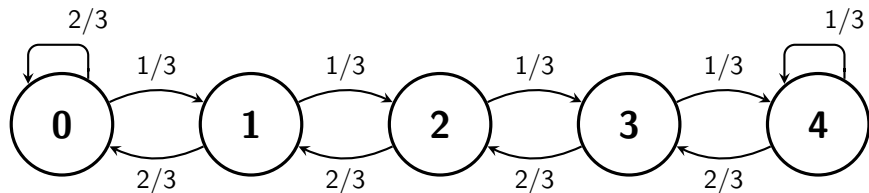
- A queue for at most 4 items.
- States of the graph encode how many items are enqueued.
- Every transitions encodes that either an item has arrived in the queue or one item has been consumed from the queue (exclusive or).
- Arrival of an item happens with the probability of $1/3$, while the dequeue operation happens with the probability of $2/3$.

Solution

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Task: create DTMC modeling the following scenario - continued

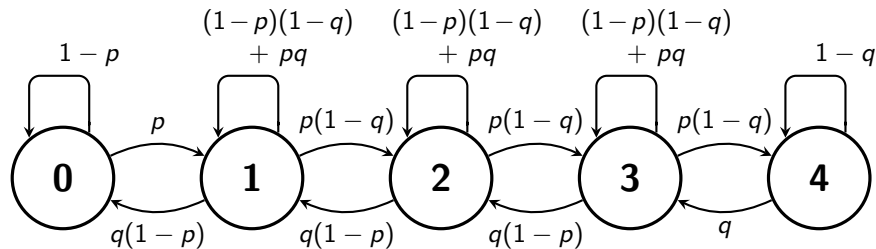
- If the actions of item arrival and item removal are independent, they both have their own probability of appearance with every time tick.
- A new item comes with probability $p = 1/2$, an item is removed with probability $q = 2/3$?
- With every time tick, one of the actions may occur, both actions may occur simultaneously, or none of them may occur at all.

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Task: create DTMC modeling the following scenario - continued

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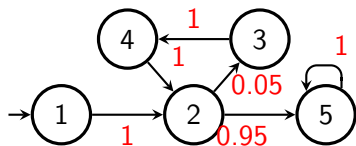


Discrete-time Markov Chain is given by

- a set of states S ,
- an initial state s_0 of S ,
- a probability matrix $P : S \times S \rightarrow [0, 1]$, and
- an interpretation of atomic propositions $I : S \rightarrow AP$.

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$$P = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0.05 & 0 & 0.95 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

Property Specification

Recall some non-probabilistic specification languages:

LTL formulae

$$\varphi ::= p \mid \neg\varphi \mid \varphi \vee \varphi \mid X\varphi \mid \varphi U \varphi$$

CTL formulae

$$\varphi ::= p \mid \neg\varphi \mid \varphi \vee \varphi \mid EX\varphi \mid E[\varphi U \varphi] \mid EG\varphi$$

Syntax of CTL*

state formula	$\varphi ::= p \mid \neg\varphi \mid \varphi \vee \varphi \mid E\psi$
path formula	$\psi ::= \varphi \mid \neg\psi \mid \psi \vee \psi \mid X\psi \mid \psi U \psi$

We need to quantify probability that a certain behaviour will occur.

Probabilistic Computation Tree Logic (PCTL)

Syntax of PCTL

state formula	$\varphi ::= p \mid \neg\varphi \mid \varphi \vee \varphi \mid P_{\bowtie b}\psi$
path formula	$\psi ::= X\varphi \mid \varphi U\varphi \mid \varphi U^{\leq k}\varphi$

where

- $b \in [0, 1]$ is a probability bound,
- $\bowtie \in \{\leq, <, \geq, >\}$, and
- $k \in \mathbf{N}$ is a bound on the number of steps.

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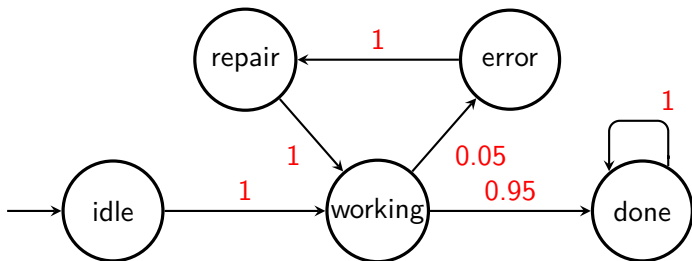
- $b \in [0, 1]$ is a probability bound,
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- $k \in \mathbf{N}$ is a bound on the number of steps.

A PCTL formula is always a state formula.

$\alpha U^{\leq k} \beta$ is a bounded until saying that α holds until β within k steps. For $k = 3$ it is equivalent to $\beta \vee (\alpha \wedge X\beta) \vee (\alpha \wedge X(\beta \vee \alpha \wedge X\beta))$.

Some tools also supports $P_{=?}\psi$ asking for the probability that the specified behaviour will occur.

We can also use derived operators like G , F , \wedge , \Rightarrow , etc.



Probabilistic reachability $P_{\geq 1}(F \text{ done})$

- probability of reaching the state *done* is equal to 1

Probabilistic bounded reachability $P_{>0.99}(F^{\leq 6} \text{ done})$

- probability of reaching the state *done* in at most 6 steps is > 0.99

Probabilistic until $P_{<0.96}((\neg \text{error}) U \text{ done})$

- probability of reaching *done* with no visit of *error* is less than 0.96

Qualitative PCTL properties

- $P_{\bowtie b} \psi$ where b is either 0 or 1

Quantitative PCTL properties

- $P_{\bowtie b} \psi$ where $b \in (0, 1)$

In DTMC where zero probability edges are erased, it holds that

- $P_{>0}(X \varphi)$ is equivalent to $EX \varphi$
 - there is a next state satisfying φ
- $P_{\geq 1}(X \varphi)$ is equivalent to $AX \varphi$
 - the next states satisfy φ
- $P_{>0}(F \varphi)$ is equivalent to $EF \varphi$
 - there exists a finite path to a state satisfying φ

but

- $P_{\geq 1}(F \varphi)$ is **not** equivalent to $AF \varphi$
(see, e.g., *AF done* on our running example)

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- $P_{\geq 1}(F \varphi)$ is **not** equivalent to $AF \varphi$
(see, e.g., *AF done* on our running example)

There is no CTL formula equivalent to $P_{\geq 1}(F \varphi)$,
and no PCTL formula equivalent to $AF \varphi$.

Analysis of Discrete-time Markov Chains

Transient analysis

- probability distribution after k -steps
- probability of reaching a state within k -steps

Long run analysis

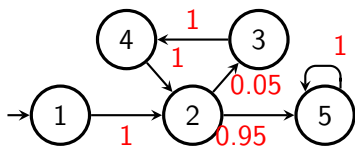
- states visited infinitely often with probability one
- stationary (invariant) distribution

Model Checking

- model checking DTMCs
- model checking MDPs

Transient Analysis

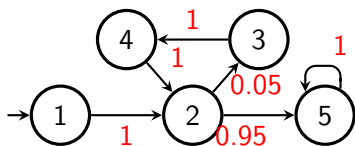
Quantitative - forward reachability



$$P = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0.05 & 0 & 0.95 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

Probability distribution after k steps when starting in 1

Quantitative - forward reachability

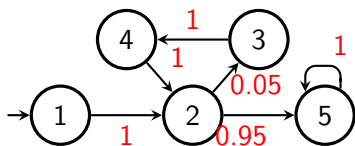


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Probability distribution after k steps when starting in 1

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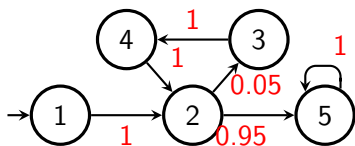
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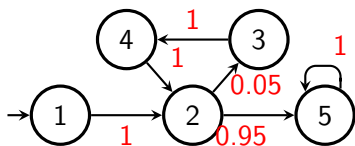
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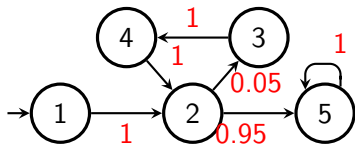
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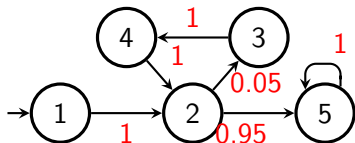
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$$\begin{bmatrix} 1 & 0 & 0 & 0 & 0 \end{bmatrix} \times P^5 = \begin{bmatrix} 0 & 0 & 0.0025 & 0 & 0.9975 \end{bmatrix}$$

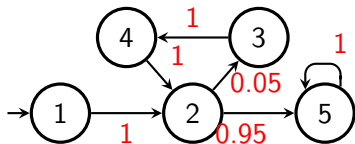
Quantitative - backward reachability



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Prob. of being in states 2 or 5 after k steps, i.e. $P_{=?} F^{=k}(2 \vee 5)$

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Prob. of being in states 2 or 5 after k steps, i.e. $P_{=?} F^{=k}(2 \vee 5)$

$$P \times [0 \ 1 \ 0 \ 0 \ 1]^T = [1 \ 0.95 \ 0 \ 1 \ 1]^T$$

$$P^2 \times [0 \ 1 \ 0 \ 0 \ 1]^T = [0.95 \ 0.95 \ 1 \ 0.95 \ 1]^T$$

$$P^3 \times [0 \ 1 \ 0 \ 0 \ 1]^T = [0.95 \ 1 \ 0.95 \ 0.95 \ 1]^T$$

$$P^4 \times [0 \ 1 \ 0 \ 0 \ 1]^T = [1 \ 0.9975 \ 0.95 \ 1 \ 1]^T$$

$$P^5 \times [0 \ 1 \ 0 \ 0 \ 1]^T = [0.9975 \ 0.9975 \ 1 \ 0.9975 \ 1]^T$$

Unbounded reachability

- Let $p(s, A)$ be the probability of reaching a state in A from s .

Observation: It holds that:

- $p(s, A) = 1$ for $s \in A$
- $p(s, A) = \sum_{s' \in \text{succ}(s)} P(s, s') * p(s', A)$ for $s \notin A$

where $\text{succ}(s)$ is a set of successors of s and $P(s, s')$ is the probability on the edge from s to s' .

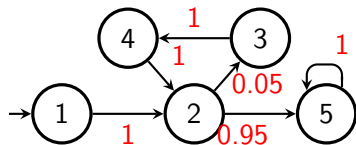
Theorem

- The minimal non-negative solution of the above equations equals to the probability of unbounded reachability.

"Up to" reachability

Task

- For the given DTMC compute the probability of reaching state 3 within 6 steps.
- Compute $P_{=?} F^{\leq 6} 3$.



Wrong Solution

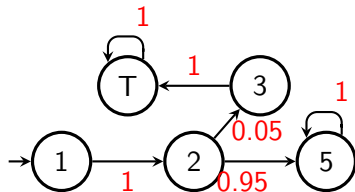
- BEWARE!
- We cannot sum the probabilities of repeated visits!

$$P_{=?} F^{\leq 6} 3 \neq \sum_{i=0}^6 P_{=?} F^{=i} 3$$

"Up to" reachability – continued

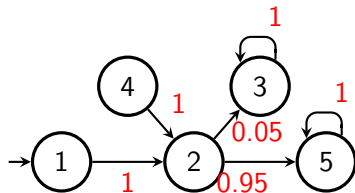
Possible Solution 1

- We may only sum the probabilities if we make sure, that no revisit of a state is possible.
- We have to modify the DTMC.
- $P_{=?} F^{\leq 6} 3 = \sum_{i=0}^6 P_{=?} F^{=i} 3$



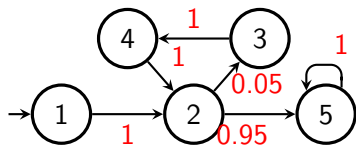
Possible Solution 2

- Alternatively, we can make the target state absorbing.
- $P_{=?} F^{\leq 6} 3 = P_{=?} F^{=6} 3$



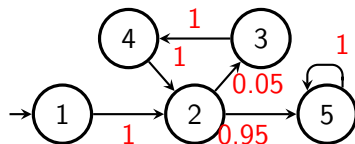
Long Run Analysis

Long run analysis

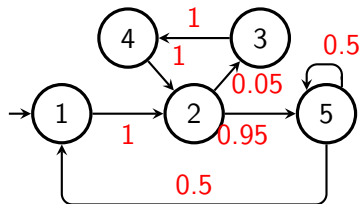


Recall that we reach the state 5(*done*) with probability 1.

Long run analysis



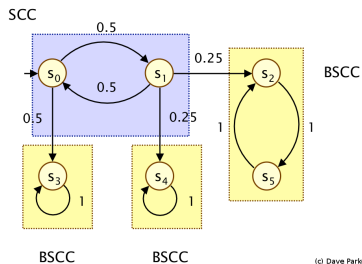
Recall that we reach the state 5 (*done*) with probability 1.



What are the states visited infinitely often with probability 1?

States visited infinitely often

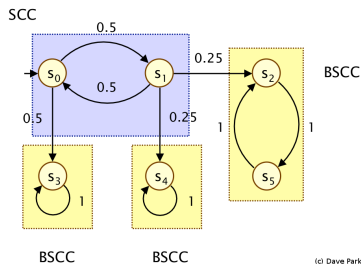
Decompose the graph representation onto strongly connected components.



¹This holds only in DTMC models with finitely many states.

States visited infinitely often

Decompose the graph representation onto strongly connected components.

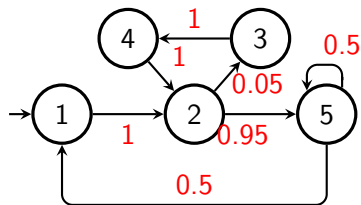


Theorem ¹

- A state is **not visited** or **visited infinitely often** with probability 1 if and only if it is in a **bottom strongly connected component**.
- All other states are **visited finitely many times** with probability 1.

¹This holds only in DTMC models with finitely many states.

How often is a state visited among the states visited infinitely many times?



Theorem

$$\lim_{n \rightarrow \infty} E \left(\frac{\# \text{ visits of state } i \text{ during the first } n \text{ steps}}{n} \right) = \pi_i$$

where π is a so called **stationary** (or **steady-state** or **invariant** or **equilibrium**) **distribution** satisfying $\pi \times P = \pi$.

DTMC Extensions

Markov Decision Processes (MDP)

- Extends DTMC with non-determinism.
- For a given state, there is a choice of probability distribution we may use to proceed to the next state (non-deterministic choice of action, every action represents one probability distribution over the successors).

Model Checking MDPs

- Satisfaction of a property ranges between **Pmin** and **Pmax** depending on the resolution of the non-determinism.
- By resolving the non-determinism in MDP we get DTMC.
- PRISM – Probabilistic model checker

Other DTMC, MDP Extensions

- Rewards
- Partial observability