

# Game theory 2

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# Mixed-strategy Nash equilibrium

# Matching pennies

- Two players
- Players choose heads or tails
- If players match heads/tails, I (Player 1) win both coins
- If players don't match heads/tails, opponent (Player 2) wins both coins

# Matching pennies

My pair

		My pair	
		Heads	Tails
Me	Heads	1, -1	-1, 1
	Tails	-1, 1	1, -1

# Matching pennies – mixed strategy

My pair

		My pair	
		Heads (0.5)	Tails (0.5)
Me	Heads (0.5)	1, -1	-1, 1
	Tails (0.5)	-1, 1	1, -1

# Calculation of mixed-strategy NE

# Game Y

		B	
		L (q)	R (1 - q)
A	U (p)	3, -3	-2, 2
	D (1 - p)	-1, 1	0, 0

# Game Y – Player A

- Player A plans to mix Up and Down strategy at a certain ratio  $p$
- Player B might play Left or Right
- Player A must find such a probability of playing U and D that makes Player B indifferent to selecting L or R
- Player B has to gain same utility from B's choice Left and Right
  - $EU_L = EU_R$
- Expected utility of Player B choosing Left:
  - $EU_L = f(p)$
- Expected utility of Player B choosing Right:
  - $EU_R = f(p)$



# Game Y - Player A's strategy

- $EU_L$ :
  - Some % of time ( $p$ ) gets B utility  $-3$
  - Rest of the time ( $1 - p$ ) gets B utility  $1$

- $EU_L = (p) * (-3) + (1 - p) * (1)$

- $EU_L = -3p + 1 - p$

- $EU_L = 1 - 4p$

B

		L ( $q$ )	R ( $1 - q$ )
A	U ( $p$ )	3, -3	-2, 2
	D ( $1 - p$ )	-1, 1	0, 0

# Game Y - Player A's strategy

- $EU_R$ :
  - Some % of time ( $p$ ) gets B utility 2
  - Rest of the time ( $1 - p$ ) gets B utility 0

- $EU_R = (p) * (2) + (1 - p) * (0)$

- $EU_R = 2p + 0 - 0p$

- $EU_R = 2p$

		B	
		L ( $q$ )	R ( $1 - q$ )
A	U ( $p$ )	3, -3	-2, 2
	D ( $1 - p$ )	-1, 1	0, 0

# Player A's strategy – making B indifferent

## Comparison of $EU_L$ with $EU_R$

- $EU_L = 1 - 4p$

- $EU_R = 2p$

- $EU_L = EU_R$

- $1 - 4p = 2p$        $+4p$

- $1 = 6p$        $/6$

- $p = 1/6$

- $1 - p = 1 - 1/6 = 5/6$

- We've found the ideal mixed strategy for Player A

- If Player A plays Up  $1/6$  of time and Down  $5/6$  of time, Player B is indifferent to choosing Left or Right

- We need to do the same for player B

# Game Y - Player B's strategy

- $EU_U$ :
  - Some % of time ( $q$ ) gets A utility 3
  - Rest of the time ( $1 - q$ ) gets A utility -2

- $EU_U = (q) * (3) + (1 - q) * (-2)$

- $EU_U = 3q - 2 + 2q$

- $EU_U = 5q - 2$

B

		L ( $q$ )	R ( $1 - q$ )
A	U ( $p$ )	3, -3	-2, 2
	D ( $1 - p$ )	-1, 1	0, 0

# Game Y - Player B's strategy

- $EU_D$ :
  - Some % of time ( $q$ ) gets **A** utility **-1**
  - Rest of the time ( $1 - q$ ) gets **A** utility **0**

- $EU_D = (q)*(-1) + (1 - q)*(0)$

- $EU_D = -1q + 0 - 0q$

- $EU_D = -q$

B

		B	
		L ( $q$ )	R ( $1 - q$ )
A	U ( $p$ )	3, -3	-2, 2
	D ( $1 - p$ )	-1, 1	0, 0

# Player B's strategy – making A indifferent

## Comparison of $EU_U$ with $EU_D$

- $EU_U = 5q - 2$

- $EU_D = -q$

- $EU_U = EU_D$

- $5q - 2 = -q$        $-5q$

- $-2 = -6q$        $/-6$

- $q = 1/3$

- $1 - q = 1 - 1/3 = 2/3$

- We've found the ideal mixed strategy for Player B

- If Player B plays Left 1/3 of time and Down 2/3 of time, Player A is indifferent to choosing Up or Down

Mixed strategy NE

(  $\frac{1}{6}$  U ,  $\frac{1}{3}$  L )

# Game Y - MSNE

B

		B	
		L (1/3)	R (2/3)
A	U (1/6)	3, -3	-2, 2
	D (5/6)	-1, 1	0, 0



# Battle of sexes

- Want to go out together but have no means of communication
  - Have 2 choices – ballet or box fight
  - Player A prefers box fight
  - Player B prefers ballet
  - Both prefer being together than being alone
- 
- Preferences for player A:  $F > B > A$
  - Preferences for player B:  $B > F > A$

# Battle of sexes

Battle of sexes

B

		b	f
B	1, 2	0, 0	
F	0, 0	2, 1	

A

Detailed description: A 2x2 payoff matrix for a 'Battle of sexes' game. The row player is labeled 'A' and has strategies 'B' and 'F'. The column player is labeled 'B' and has strategies 'b' and 'f'. The payoffs are given as (A's payoff, B's payoff). The matrix shows that if both players choose the same strategy (B,b or F,f), they both receive a positive payoff. If they choose different strategies (B,f or F,b), both receive a payoff of 0.

# Battle of sexes – PS equilibriums

Battle of sexes

B

		B	
		b	f
A	B	1, 2	0, 0
	F	0, 0	2, 1

# Equilibriums

- 2 pure-strategies equilibriums
- How would they coordinate?
- Apart from pure strategies equilibriums there is one mixed strategy equilibrium for this game
- (  $\frac{1}{3} B$ ,  $\frac{2}{3} b$  )

# Battle of sexes – mixed strategy equilibrium

B

		b $\frac{2}{3}$	f $\frac{1}{3}$
A	B $\frac{1}{3}$	1, 2	0, 0
	F $\frac{2}{3}$	0, 0	2, 1

# Calculation of MS NE payoffs

# Battle of sexes – mixed-strategy NE payoffs

B

		B	
		b $\frac{2}{3}$	f $\frac{1}{3}$
A	B $\frac{1}{3}$	$1, 2$ $\frac{1}{3} * \frac{2}{3}$	$0, 0$ $\frac{1}{3} * \frac{1}{3}$
	F $\frac{2}{3}$	$0, 0$ $\frac{2}{3} * \frac{2}{3}$	$2, 1$ $\frac{2}{3} * \frac{1}{3}$

# Battle of sexes – mixed-strategy NE payoffs

B

		B	
		b $\frac{2}{3}$	f $\frac{1}{3}$
A	B $\frac{1}{3}$	1, 2 $\frac{2}{9}$	0, 0 $\frac{1}{9}$
	F $\frac{2}{3}$	0, 0 $\frac{4}{9}$	2, 1 $\frac{2}{9}$



# BoS – Payoffs for player A

- We simply multiply payoffs for player A and probabilities for each outcome and then sum them together

- Player A's payoffs:

- $u(B, b) = 1 * 2/9 = 2/9$

- $u(B, f) = 0 * 1/9 = 0$

- $u(F, b) = 0 * 4/9 = 0$

- $u(F, f) = 2 * 2/9 = 4/9$

- $EU(A) = 2/9 + 0 + 0 + 4/9$

- $EU(A) = 6/9$

- $EU(A) = 2/3$

		B	
		b 2/3	f 1/3
A	B 1/3	1, 2 2/9	0, 0 1/9
	F 2/3	0, 0 4/9	2, 1 2/9

# BoS – Payoffs for player B

- We simply multiply payoffs of player B and probabilities for each outcome and then sum them together

- Player A's payoffs:

- $u(B, b) = 2 * 2/9 = 4/9$
- $u(B, f) = 0 * 1/9 = 0$
- $u(F, b) = 0 * 4/9 = 0$
- $u(F, f) = 1 * 2/9 = 2/9$

- $EU(B) = 4/9 + 0 + 0 + 2/9$

- $EU(B) = 6/9$

- $EU(B) = 2/3$

		B	
		b 2/3	f 1/3
A	B 1/3	1, 2 2/9	0, 0 1/9
	F 2/3	0, 0 4/9	2, 1 2/9

# Battle of sexes NE

- Pure strategies NE

- ( **B** , **b** )
  - $EU(A) = 1$
  - $EU(B) = 2$
- ( **F** , **f** )
  - $EU(A) = 2$
  - $EU(B) = 1$

- Mixed strategies NE

- (  $\frac{1}{3}$  **B** ,  $\frac{2}{3}$  **b** )
  - $EU(A) = \frac{2}{3}$
  - $EU(B) = \frac{2}{3}$

B

		b	f
B	1, 2	0, 0	
F	0, 0	2, 1	

A

# FSS entrance game

- Two students meet at the main faculty entrance
- Both simultaneously decide whether to walk or stop
- If both walk, they collide and both get a bruise (payoff -5)
- If one stops and other walks
  - Student who stopped gets good karma for letting the other pass with payoff 1, but at the same time gets delayed, which is completely offsetting the value of the good karma
  - Student who walked gets to pass quickly and thus gets payoff 1
- If both stop, each would get good karma for letting the other pass, but each will get delayed

# FSS entrance game

B

		B	
		W	S
A	W	-5, -5	1, 0
	S	0, 1	0, 0

# FSS entrance game NE

B

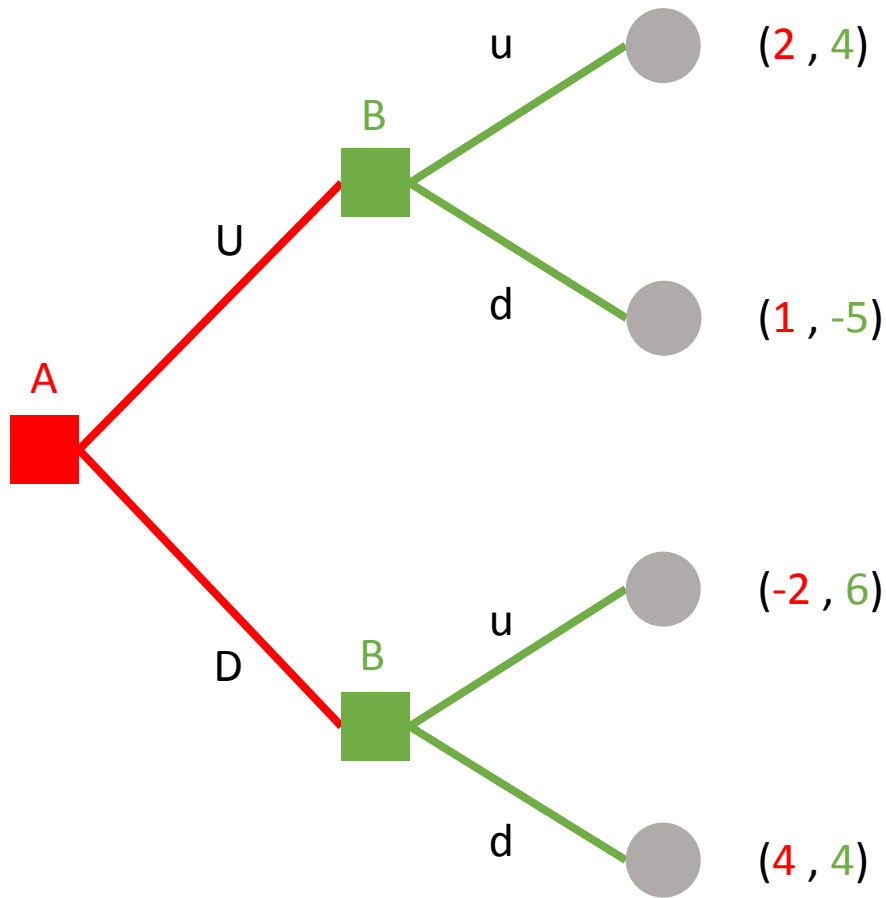
		B	
		W 1/6	S 5/6
A	W 1/6	-5, -5	1, 0
	S 5/6	0, 1	0, 0

# Extensive form games

# Extensive form games

- Visualized as a game (decision) tree
- Players move sequentially
- Captures time in game
- Captures knowledge of agents – sometimes agents do not have information where they are located in game

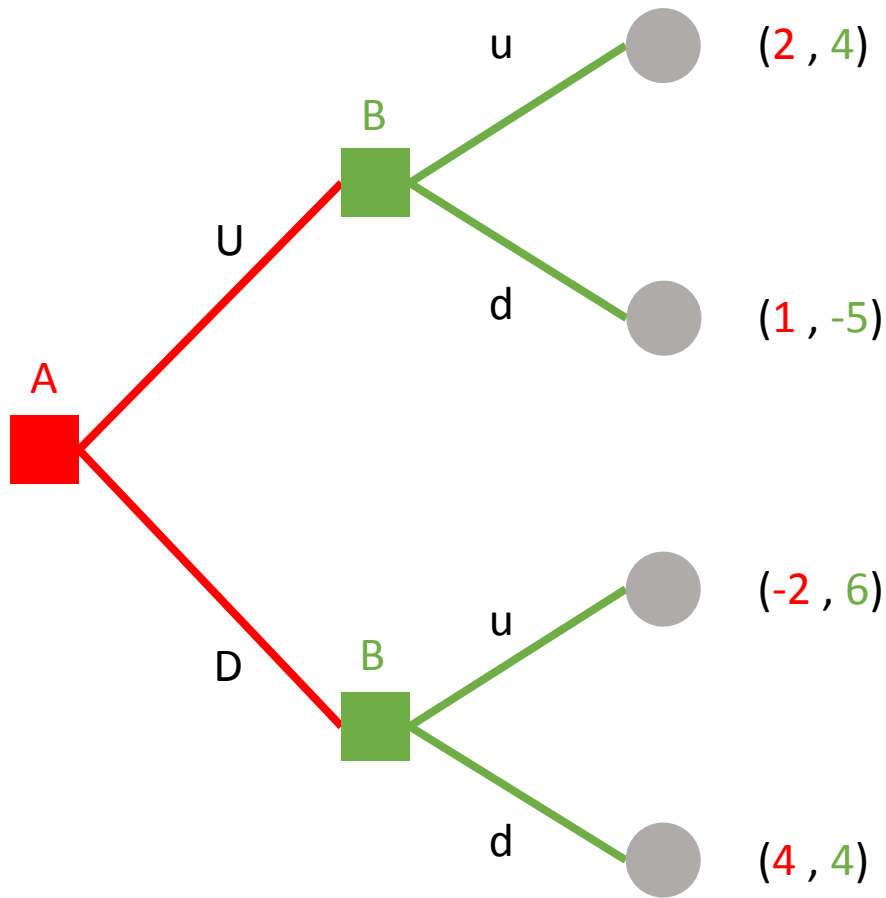




# Basic terminology

- Each square is called node
- Each line represents an action an owner of the node has at his disposal
- Nodes might either trigger other action or end
- Every circle is an end of the tree – it's called terminal node
- Every circle must yield payoffs for all the actors

# Backwards induction

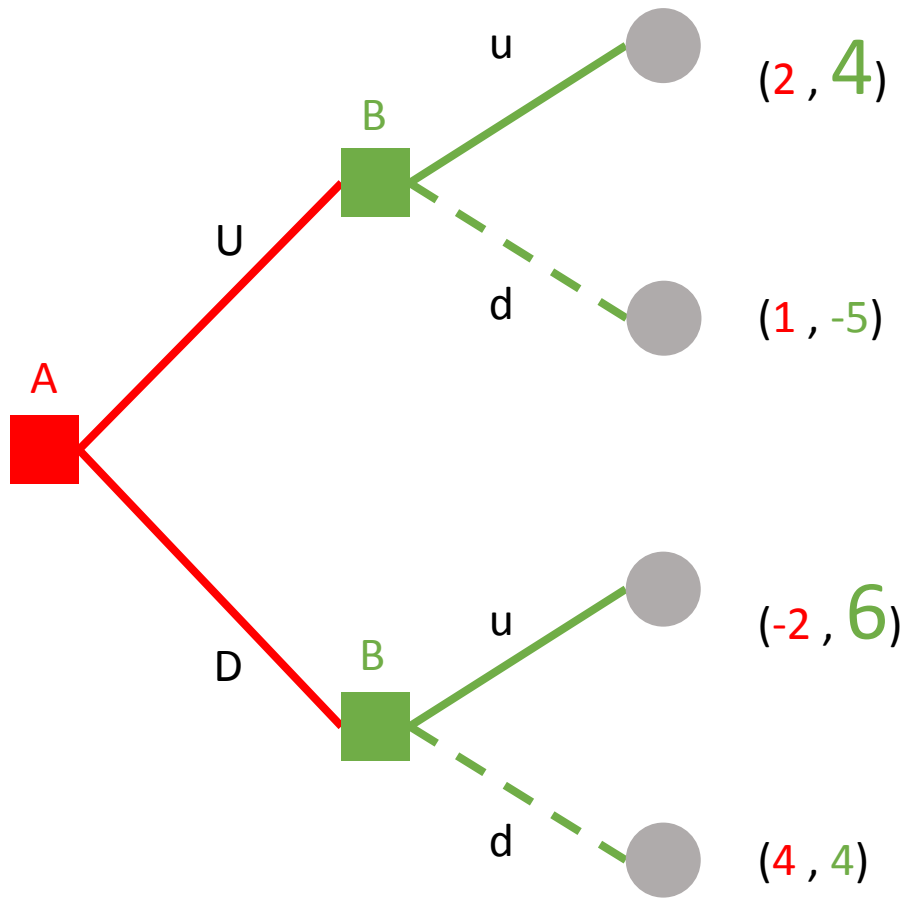


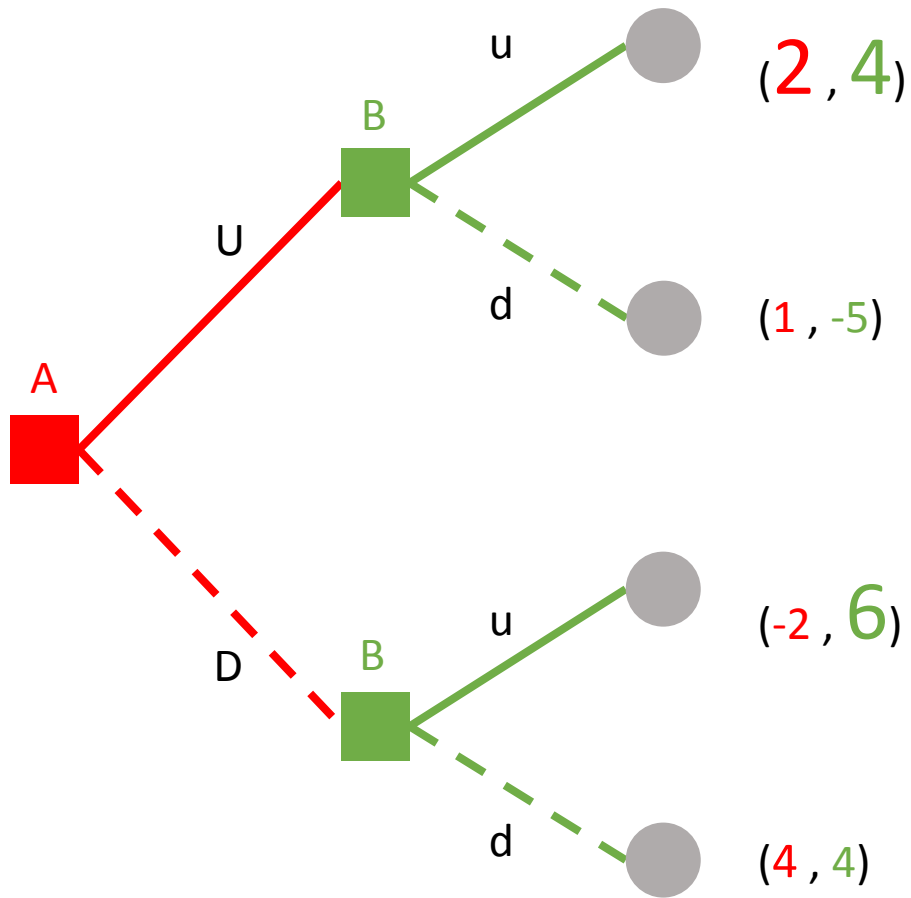
# Backwards induction

- Moves player make at nodes reached in an equilibrium are called behavior on the equilibrium path
- Moves player make at nodes that are not reached in an equilibrium are behavior off the equilibrium path
- Players can not commit themselves to make moves in future – decisions at nodes change decisions at later nodes

# Backwards induction

- Begin with decisions that lead only to terminal nodes
- Compare payoffs for decisions in each node leading to terminal node
- Find best reply to alternatives of player playing at the current node
- Work through nodes backwards and solve the outcomes of all nodes comparing payoffs for respective players







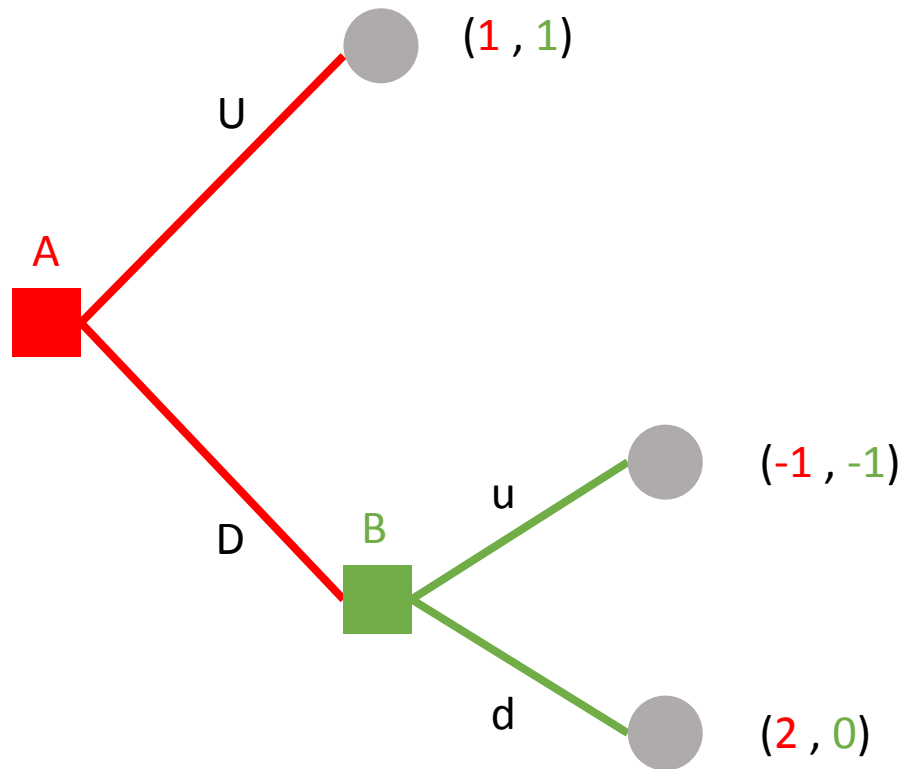
# Equilibrium of sequential game

- One Nash equilibrium in pure strategies
- ( U ; u , u )
- There may be more NE in pure strategies
- B decides – if A goes U than u yields better payoff in the upper node, if A goes D than u yields better payoff in the lower node
- A knows that B will choose u in both nodes, therefore compares payoffs in u for going U or D – U yields better payoff

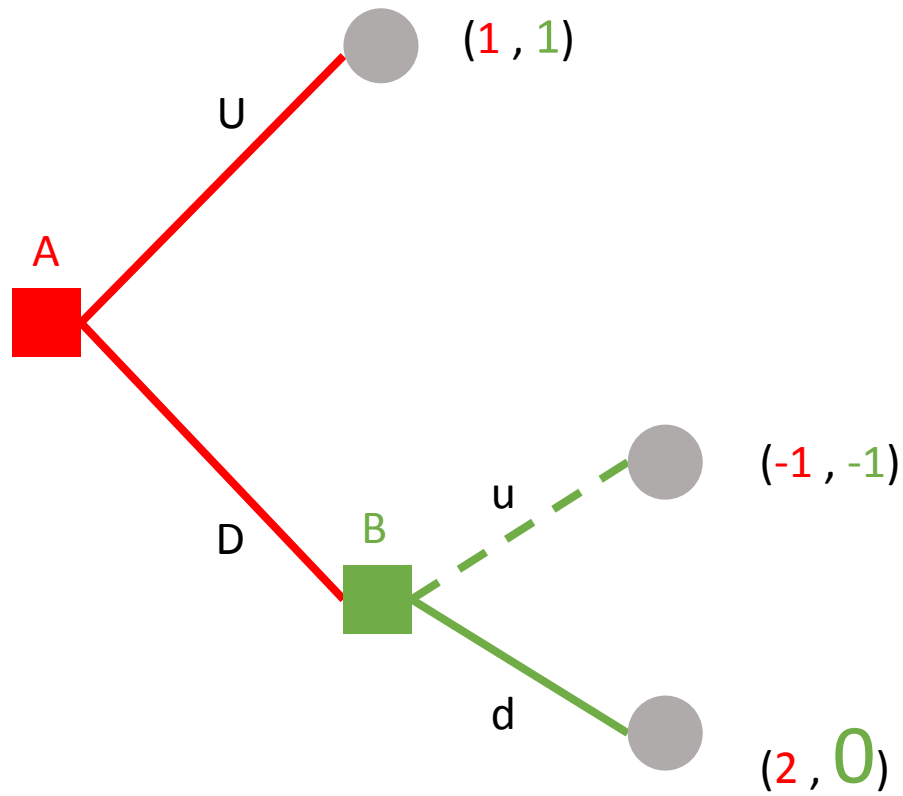
# Backwards induction

- Works easily in games of perfect information
- All actors are aware of all previous actions and can also anticipate, what actors will do based on their expected utilities over outcomes at subsequent nodes – actors have a perfect recall
- However, backwards induction assesses only rationality on the equilibrium path. NE found off the equilibrium path will not be found through backwards induction

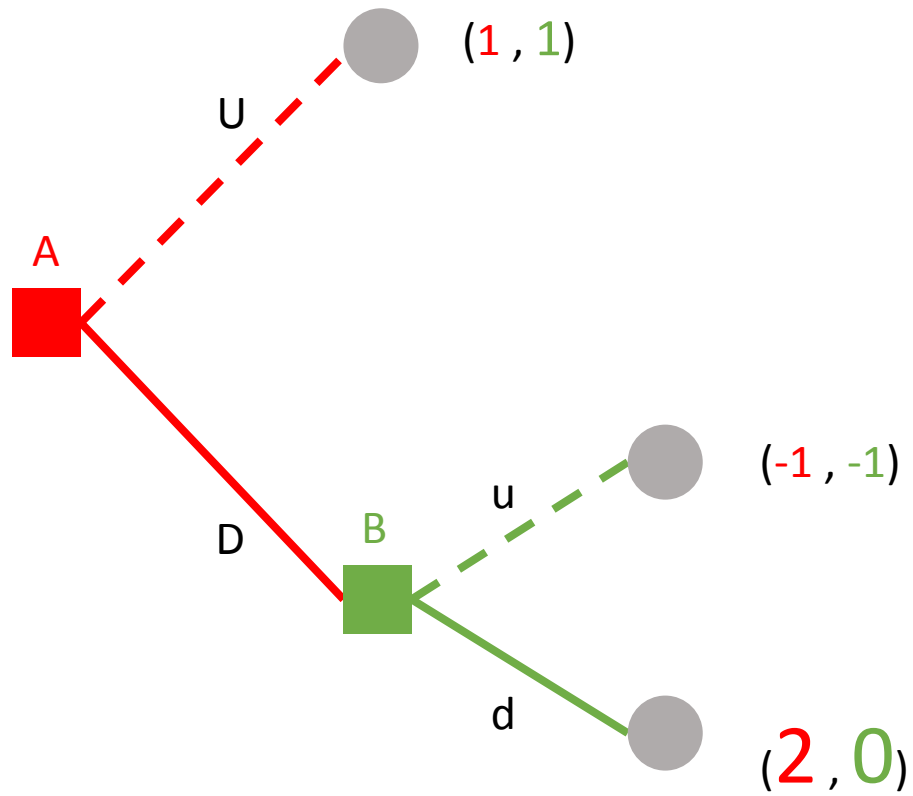
# Selten's game



# Selten's game



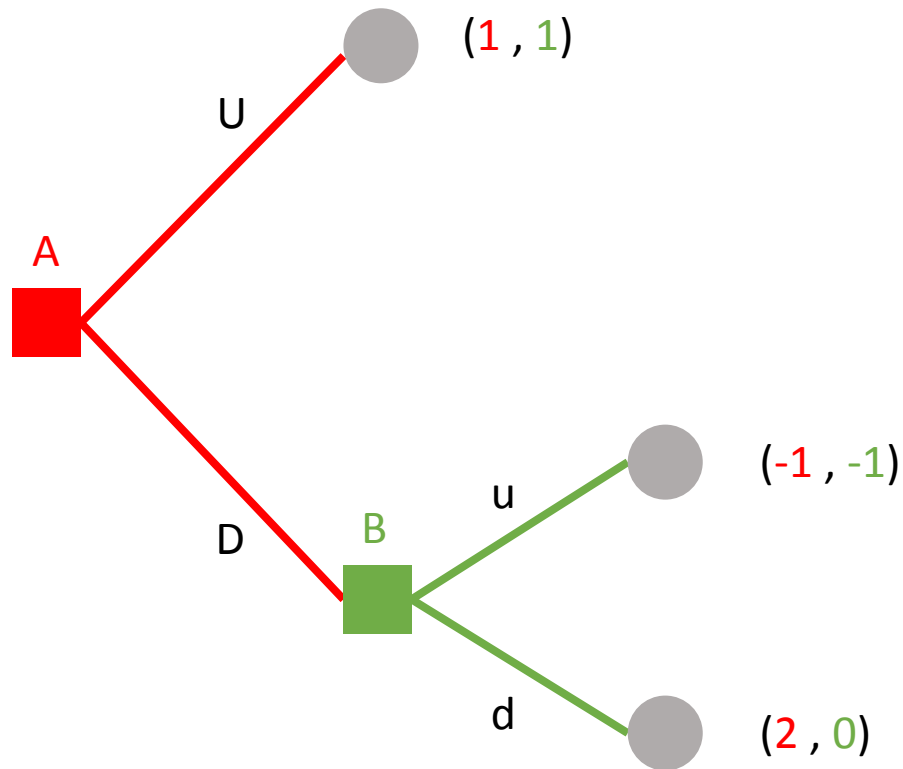
# Selten's game



# Rewrite Selten's game into matrix

- Player A has 2 moves
  - U, D
  - These constitute 2 rows in a matrix
- Player B has also 2 moves
  - u, d
  - These constitute 2 columns in a matrix
- We have 2x2 game in normal form

# Selten's game

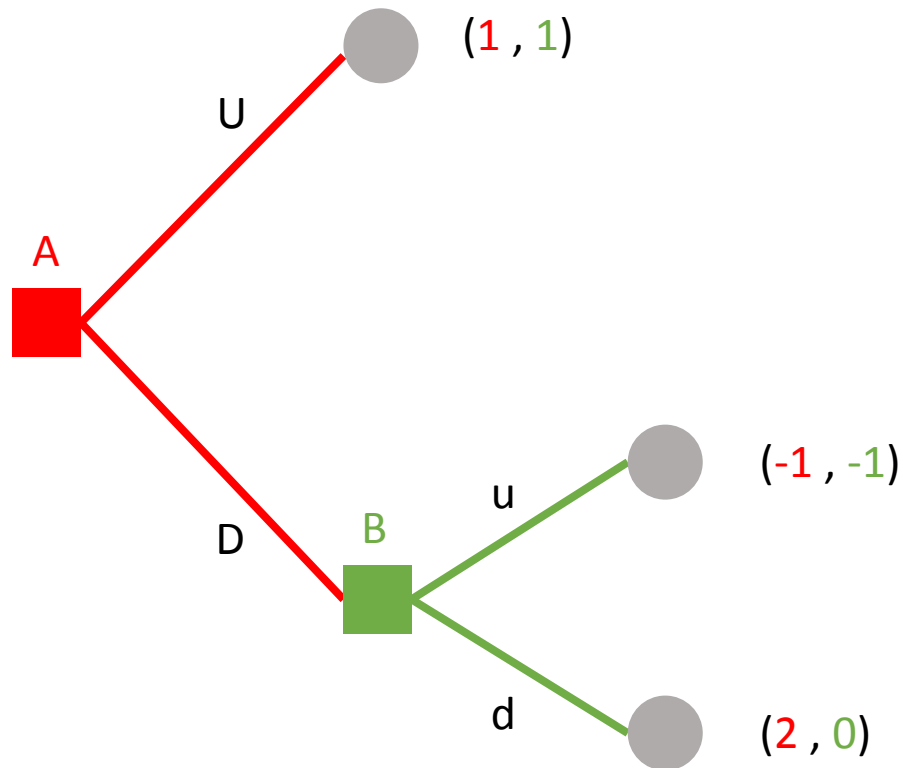


# Strategic form of Selten's game

		B	
		u	d
A	U		
	D		



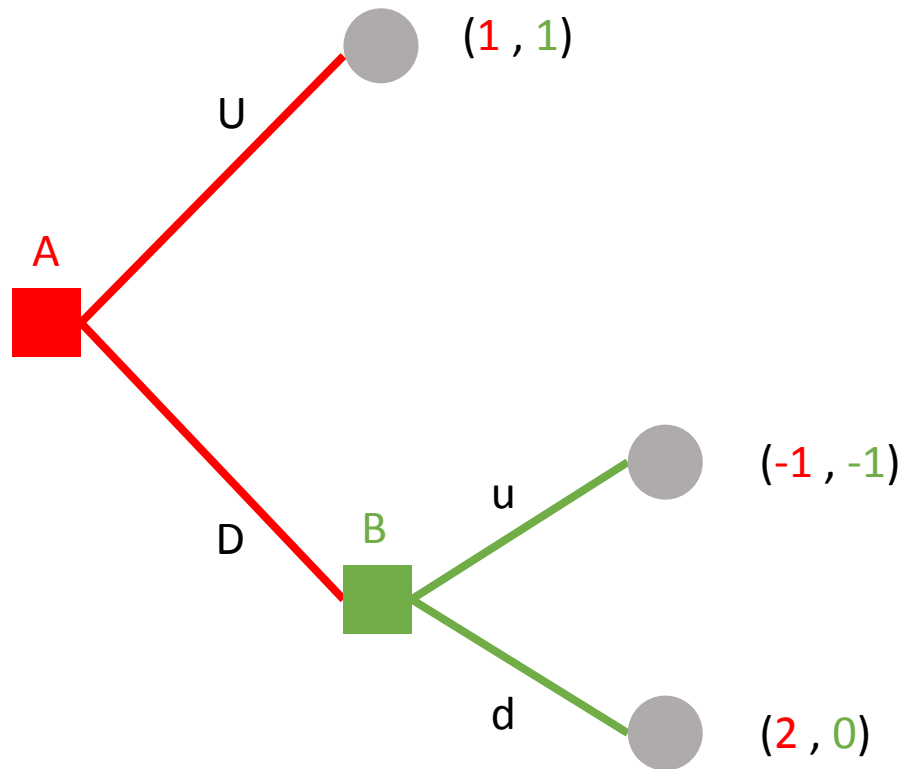
# Selten's game



# Strategic form of Selten's game

		B	
		u	d
A	U		
	D	-1, -1	2, 0

# Selten's game



# Strategic form of Selten's game

		B	
		u	d
A	U	1, 1	
	D	-1, -1	2, 0

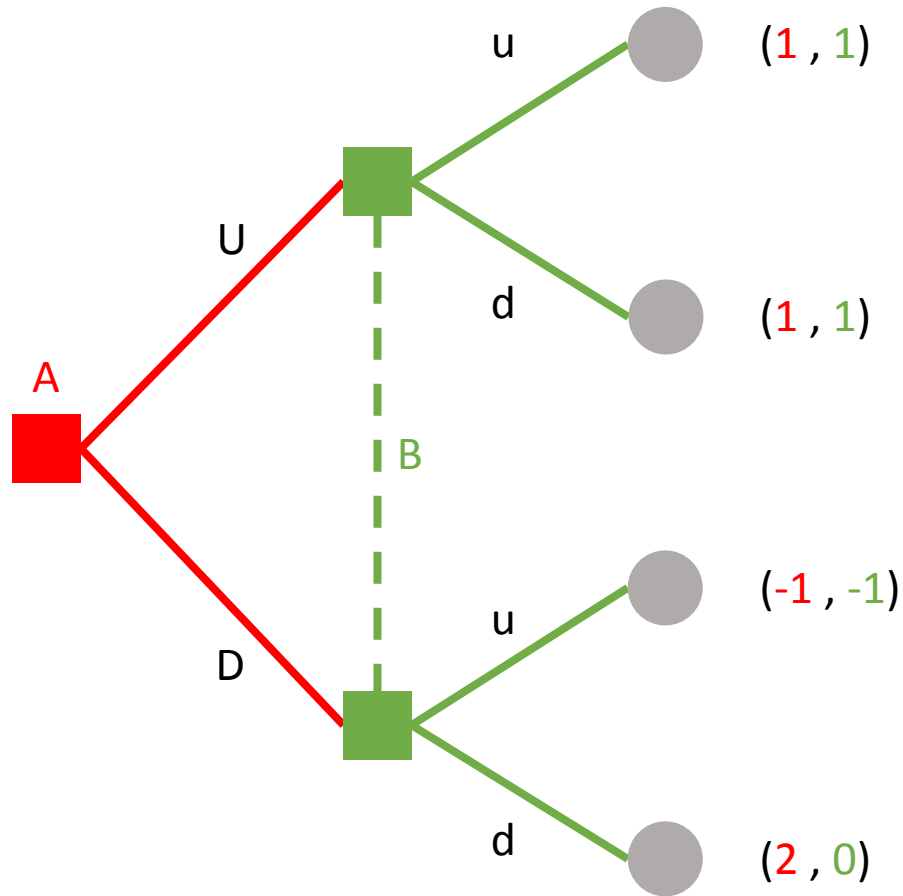
# Strategic form of Selten's game

		B	
		u	d
A	U	1, 1	1, 1
	D	-1, -1	2, 0

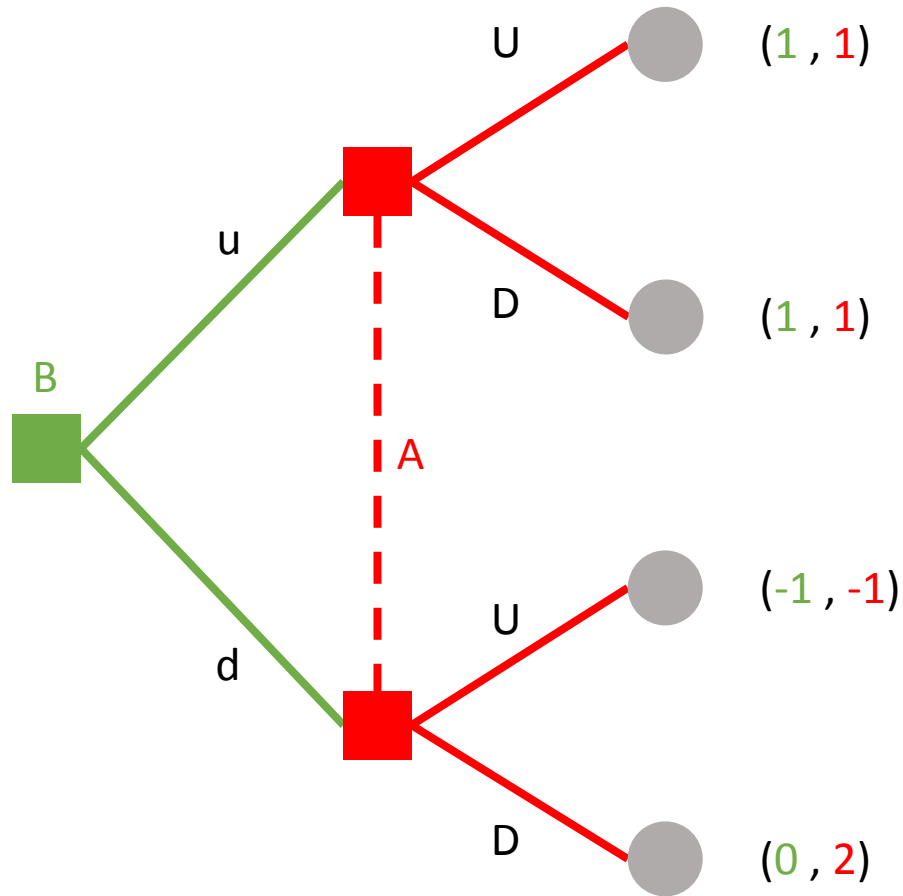
# Rewriting extensive form into normal form

- If we rewrite extensive form game into normal form, only one matrix will emerge as a representation
- This does not work the other way around
- Since matrixes do not hold information about sequence of actions, one normal-form game might have multiple extensive-form representations that would completely alter outcomes of the game

# Selten's game from matrix



# Selten's game from matrix





# Information sets

- Player's information/knowledge about moves, which have been played in the game previously
- In games of perfect information, each node is information set – each player in each moment of the game knows, which moves were played before
- If player is uncertain about previous moves, information set contains more than one node
- Information sets capture player's ignorance about previous moves in game

# Strategic form of Selten's game

		B	
		u	d
A	U	1, 1	1, 1
	D	-1, -1	2, 0

# Strategic form of Selten's game

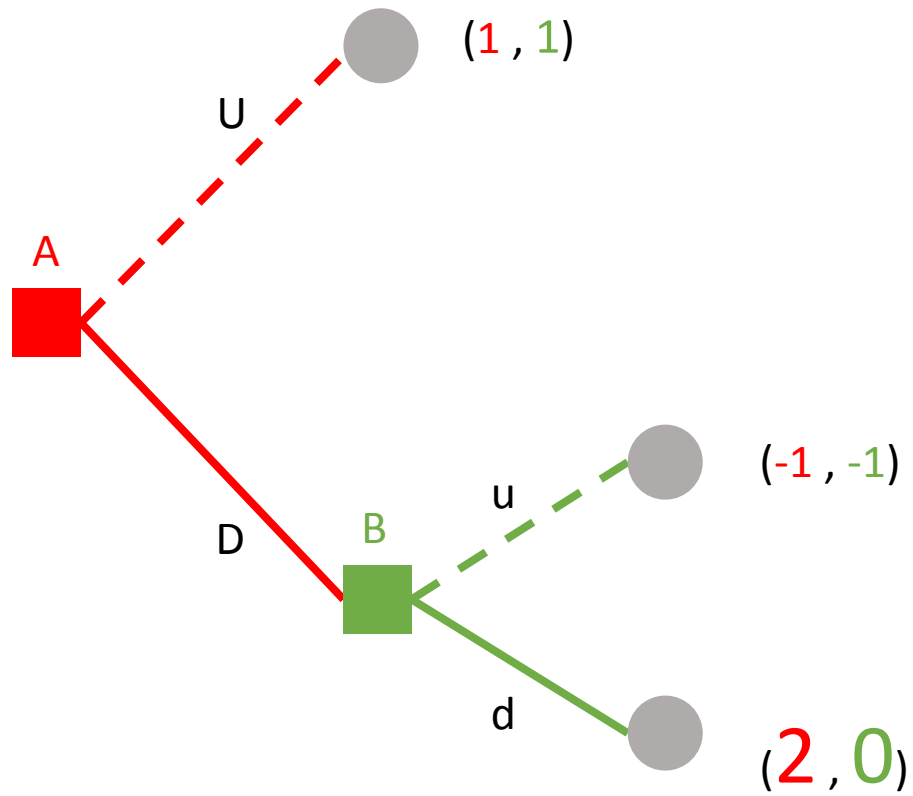
		B	
		u	d
A	U	1, 1	1, 1
	D	-1, -1	2, 0

# Strategic form of Selten's game

B

		B	
		u	d
A	U	1, 1	1, 1
	D	-1, -1	2, 0

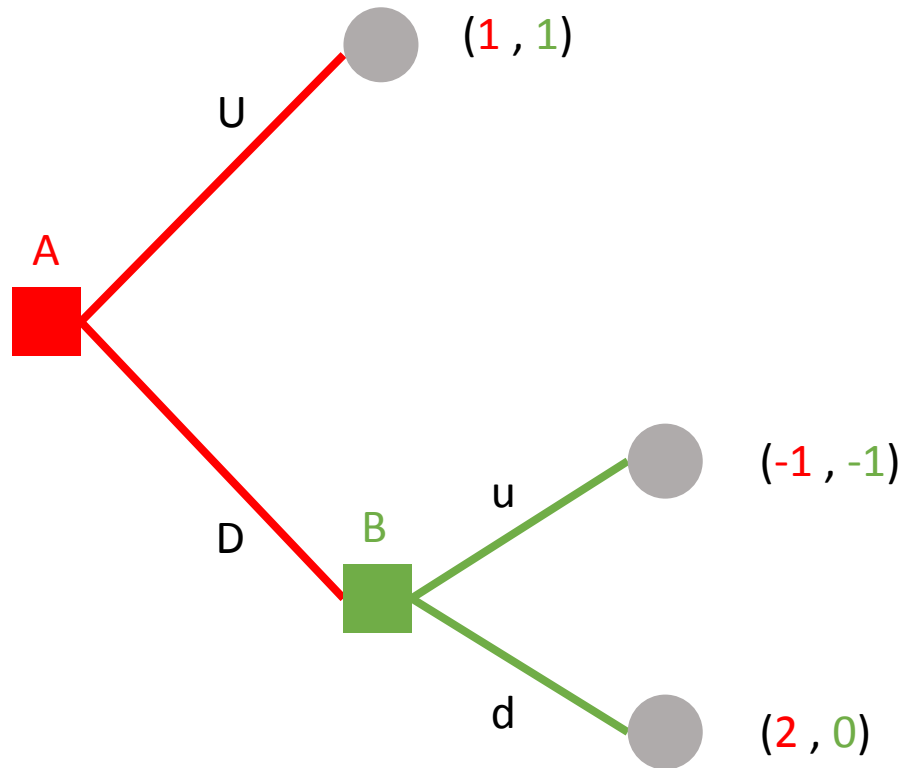
# Selten's game



# NE off the equilibrium path

- This game has another NE which is represented by action U which is not revealed in extensive form
- This equilibrium (U, u) is called a non-credible threat
  - B is willing to get its largest payoff, which will result from A playing U (deciding not to play the game)
  - B can change its payoff from A's decision about U and D only by threatening that it will play u if A plays D
  - But B will never play u, since it brings lower payoff than playing d

# Selten's game and non-credible threat



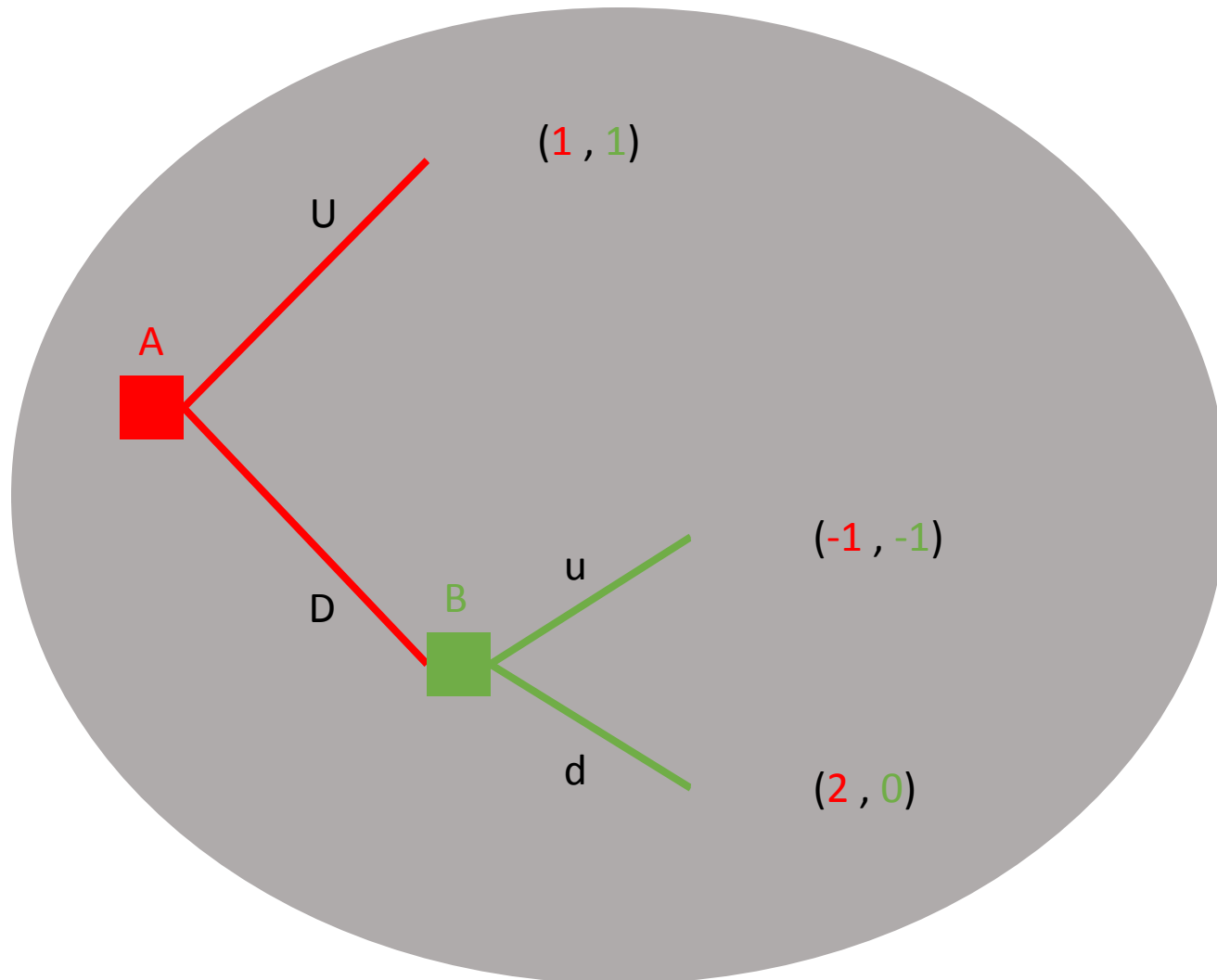
# Subgame perfection



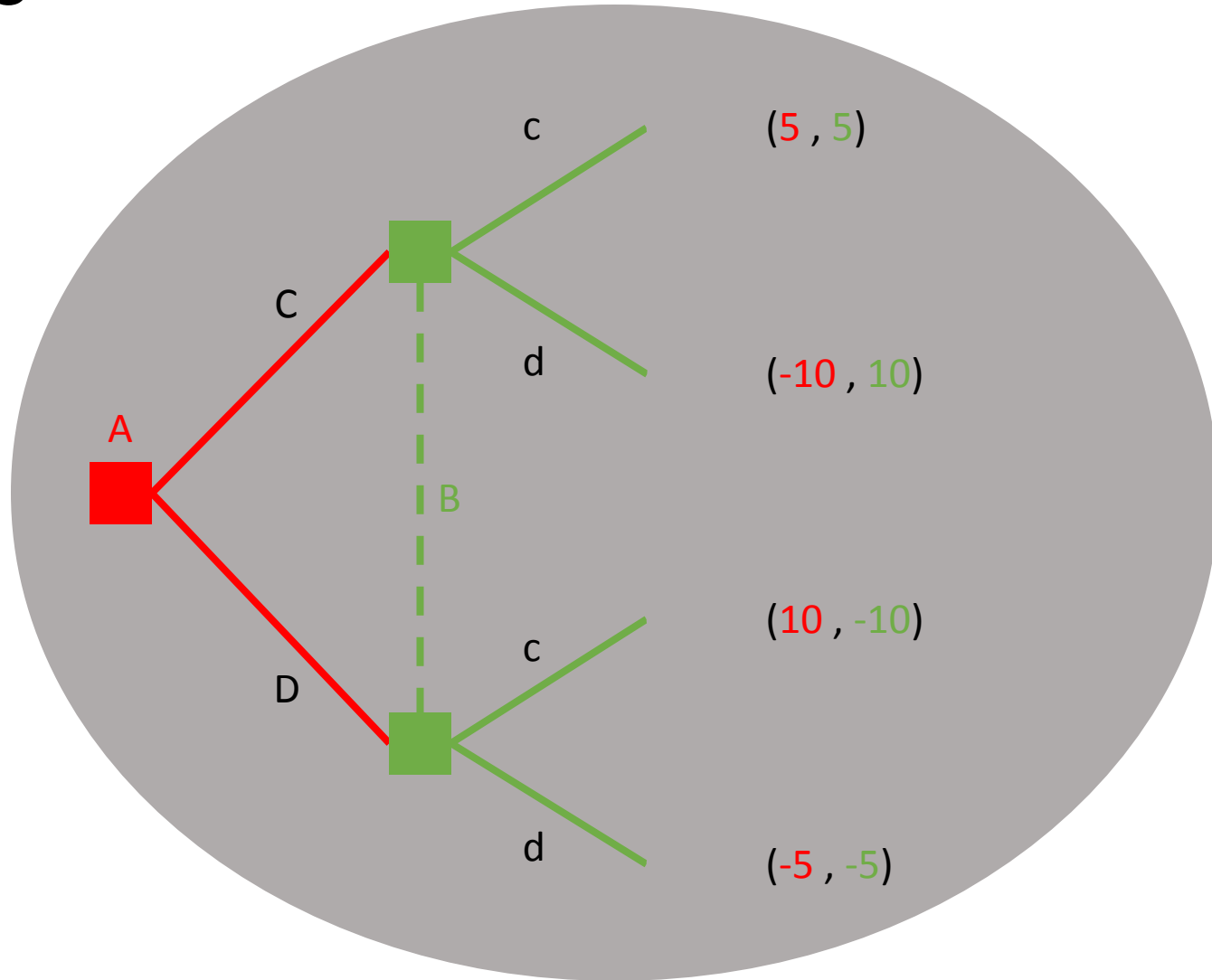
# Subgame

- Subset of the nodes of a game starting with initial node and including all its successors that preserves all information sets of the game and over which a (new) game is defined by the restriction of the original game elements to these nodes
- Subgame can be treated as a game in its own right
- To ensure each player rational, there's requirement that strategies are optimal in all proper subgames

# Subgames of Selten's game



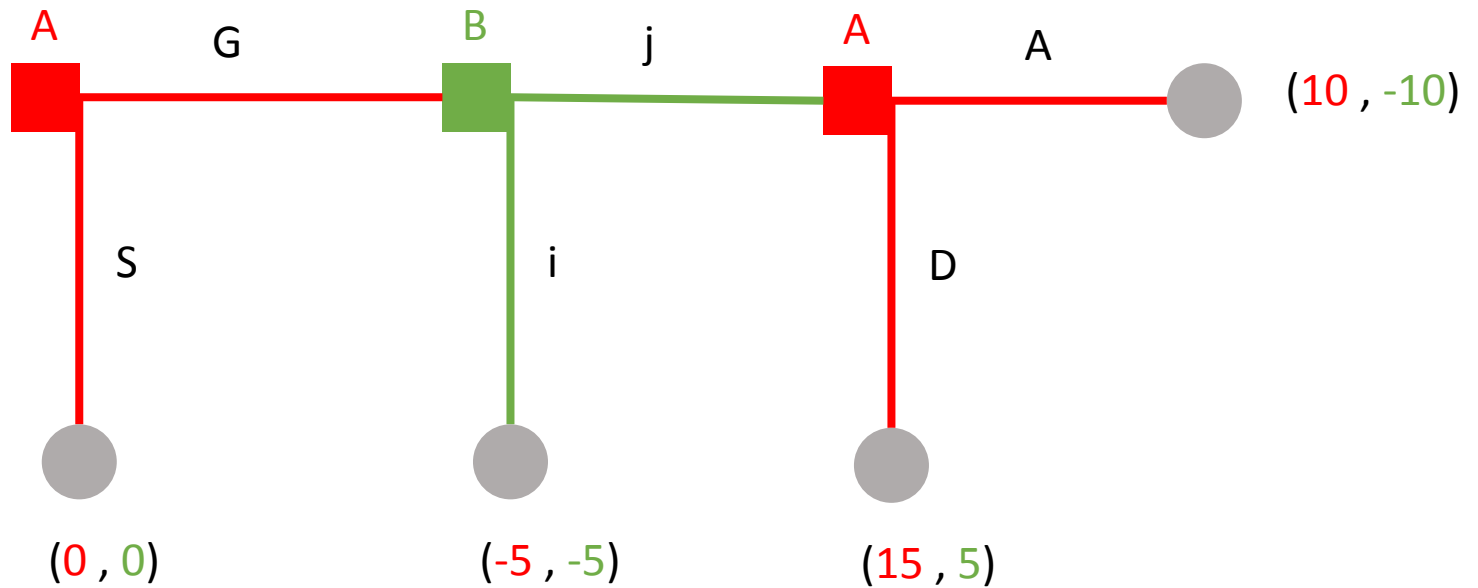
# Subgames in PD in EF



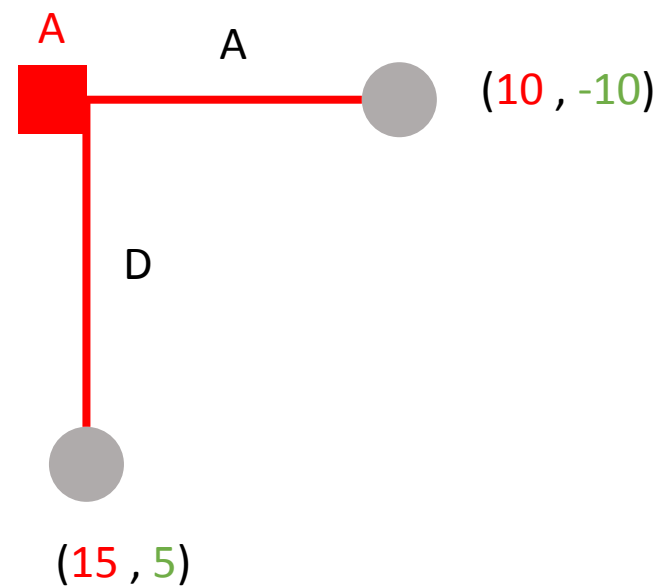
# Information sets

- Player's information/knowledge about moves, which have been played in the game previously
- In games of perfect information, each node with its actions is information set – each player in each moment of the game knows, which moves were played before
- If player is uncertain about previous moves, information set contains more than one node
- Information sets capture player's ignorance about previous moves in game

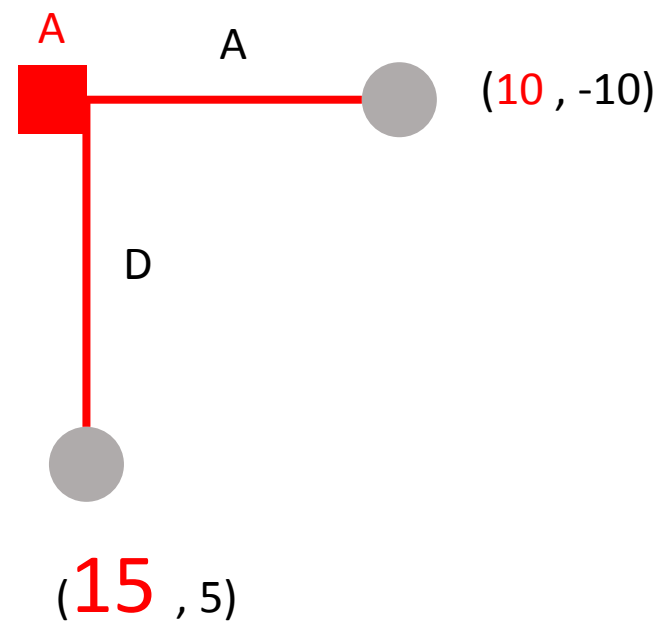
# Game A – Backwards induction



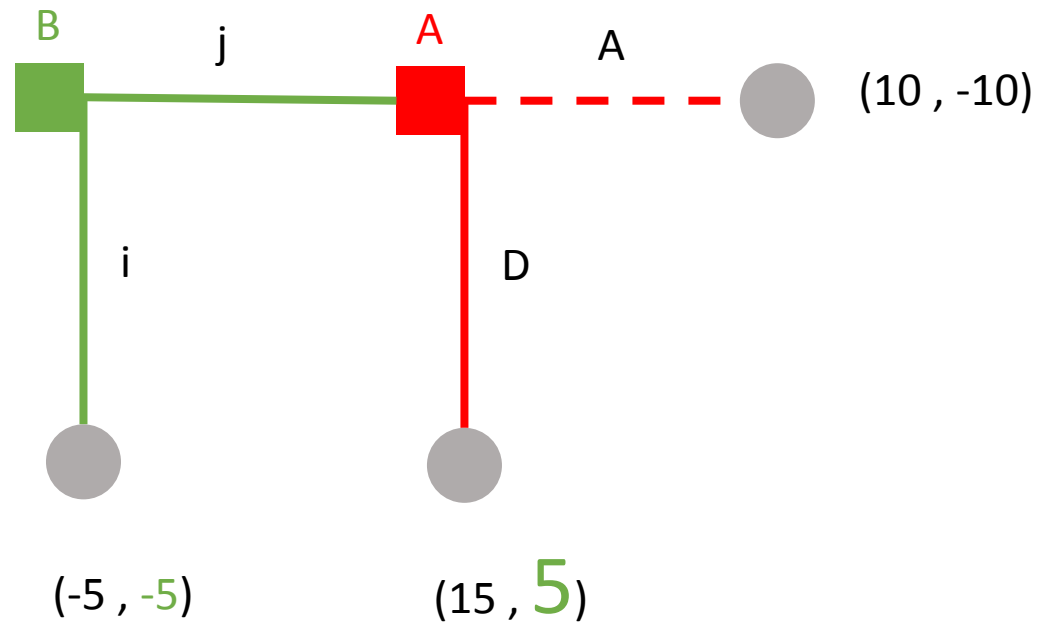
# Game A – Backwards induction



# Game A – Backwards induction

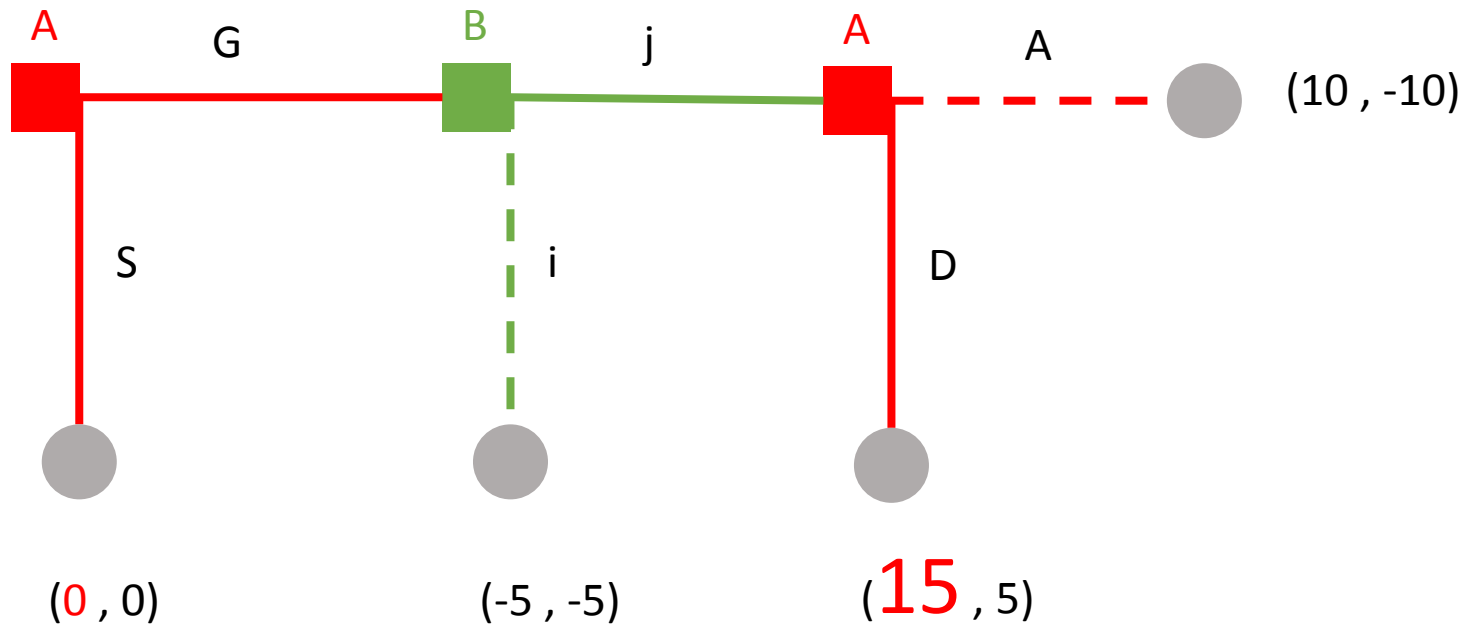


# Game A – Backwards induction

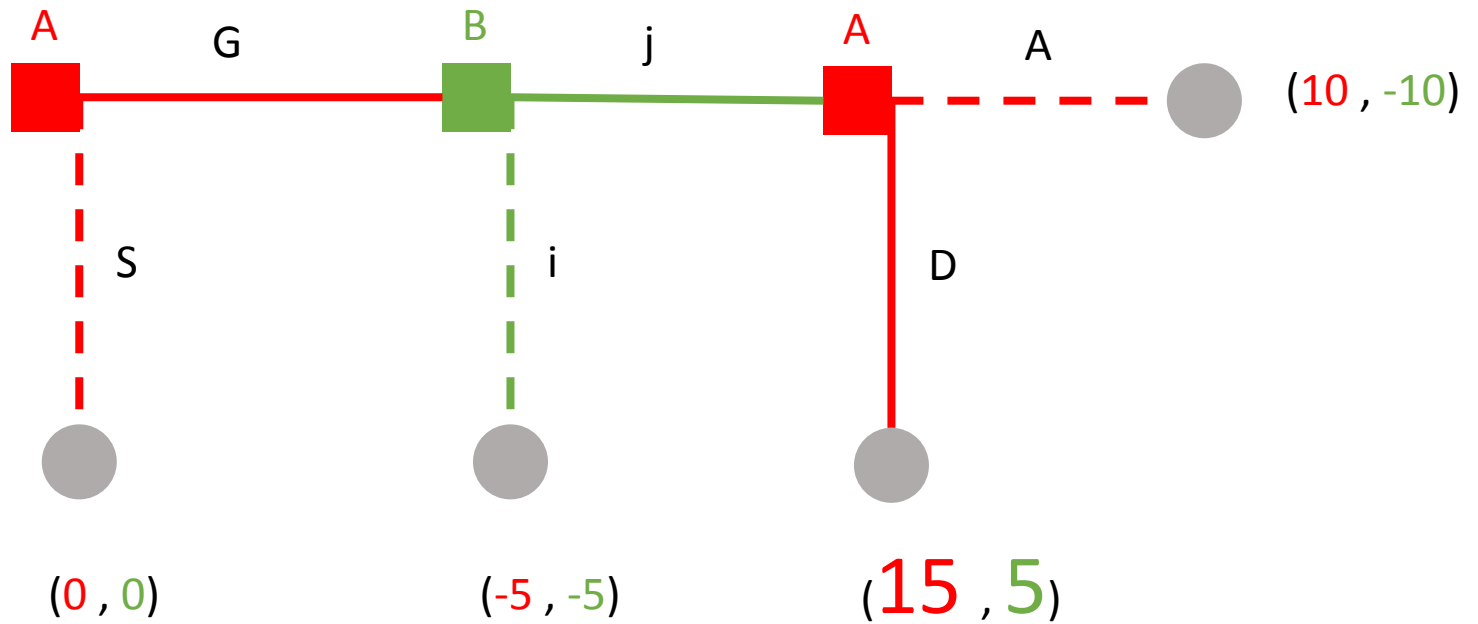




# Game A – Backwards induction



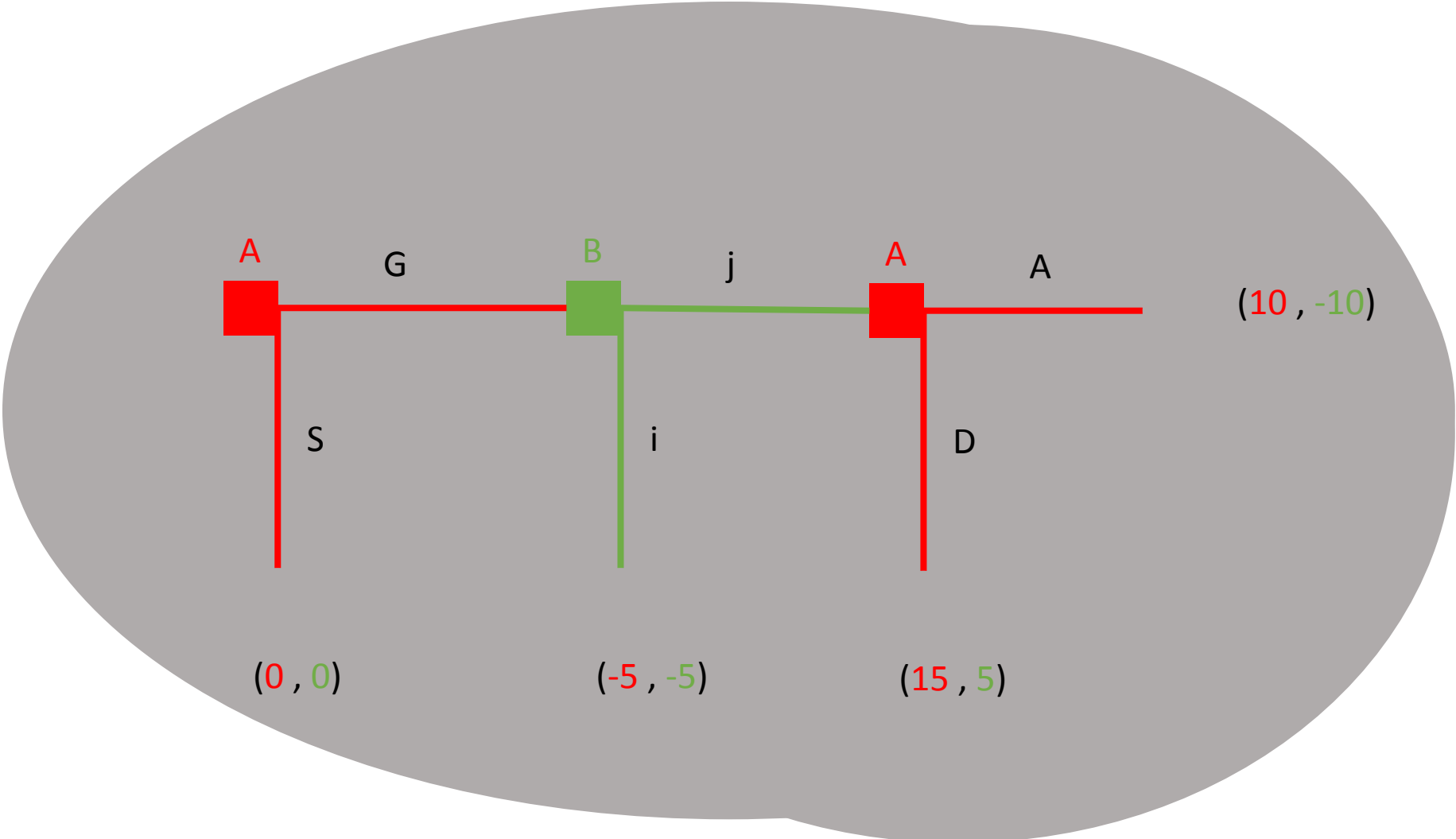
# Game A – Backwards induction solution



# Subgame perfection

- Set of strategies is subgame perfect, if for every proper subgame, the restriction of those strategies to the subgame forms a Nash Equilibrium
- SPNE must thus form NE in every subgame of the game, from the game itself to the last subgame
- Restricts equilibriums in the game which are not a credible threat throughout the game

# Game A - Subgames



# Strategic form of Subgame 1

			B
			j
A		10, -10	
	A		
D		15, 5	

# NE in strategic form of Subgame 1

			B
			j
A		10, -10	
A			
D		15, 5	

# Strategic form of Subgame 2

		B	
		j	i
A	A	10 , -10	-5 , -5
	D	15 , 5	-5 , -5

# NE in strategic form of Subgame 2

B

		J	i
A		10, -10	-5, -5
D		15, 5	-5, -5

A



# Strategic form of Subgame 3

		B	
		J	i
A	GA	10 , -10	-5 , -5
	GD	15 , 5	-5 , -5
	SA	0 , 0	0 , 0
	SD	0 , 0	0 , 0

# NE in strategic form of Subgame 3

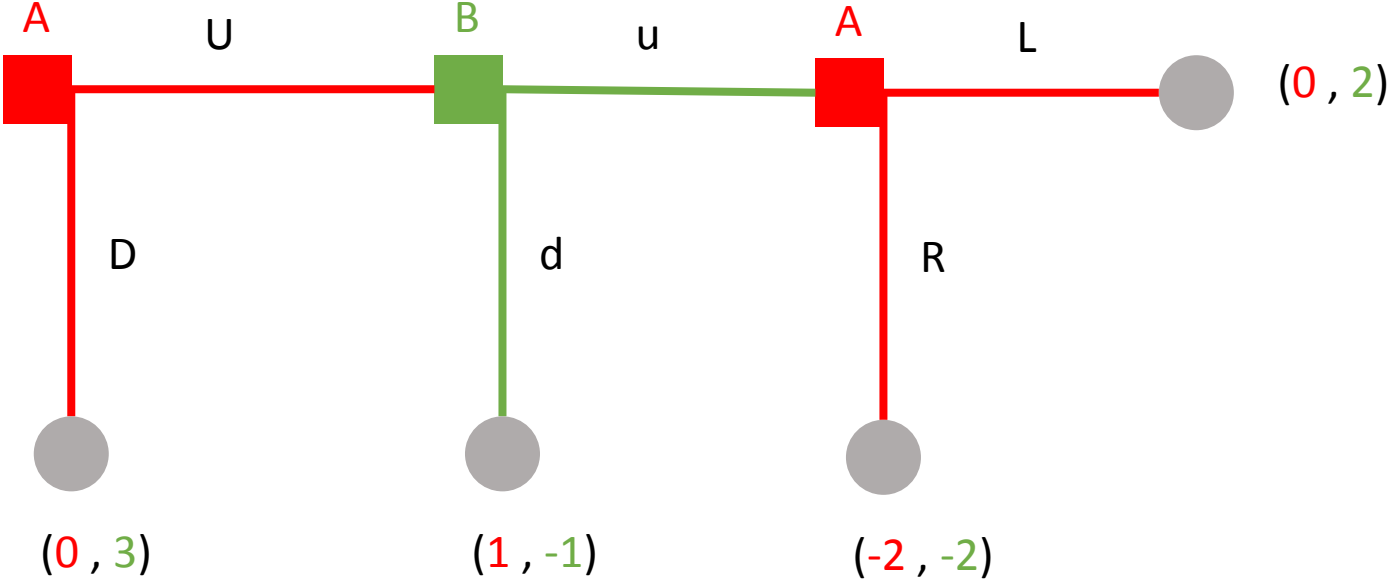
B

		J	i
A	GA	10, -10	-5, -5
	GD	15, 5	-5, -5
	SA	0, 0	0, 0
	SD	0, 0	0, 0

# NE in the Game A

- Strategic form NE
  - NE 1:  $\langle G, D; j \rangle$
  - NE 2:  $\langle S, A; i \rangle$
  - NE 3:  $\langle S, D; i \rangle$
- Backwards induction NE
  - NE 1:  $\langle G, D; j \rangle$
- Subgame perfect NE
  - NE 1:  $\langle G, D; j \rangle$

# Game B



# Strategic form of Subgame 1

			B
			u
	L	0, 2	
A	R	-2, -2	

# NE in strategic form of Subgame 1

			B
		u	
A	L	0, 2	
	R	-2, -2	

# Strategic form of Subgame 2

		B	
		u	d
A	L	0, 2	1, -1
	R	-2, -2	1, -1

# NE in strategic form of Subgame 2

B

		B	
		u	d
A	L	0, 2	1, -1
	R	-2, -2	1, -1



# Strategic form of Subgame 3

		B	
		u	d
A	UL	0, 3	0, 3
	UR	0, 3	0, 3
	DL	0, 2	1, -1
	DR	-2, -2	1, -1

# NE in strategic form of Subgame 3

B

		u	d
A	UL	0, 3	0, 3
	UR	0, 3	0, 3
	DL	0, 2	1, -1
	DR	-2, -2	1, -1

# NE in the game

- Strategic form NE

- NE 1:  $\langle U, L; u \rangle$
- NE 2:  $\langle U, R; u \rangle$
- NE 3:  $\langle D, L; u \rangle$
- NE 4:  $\langle D, R; d \rangle$

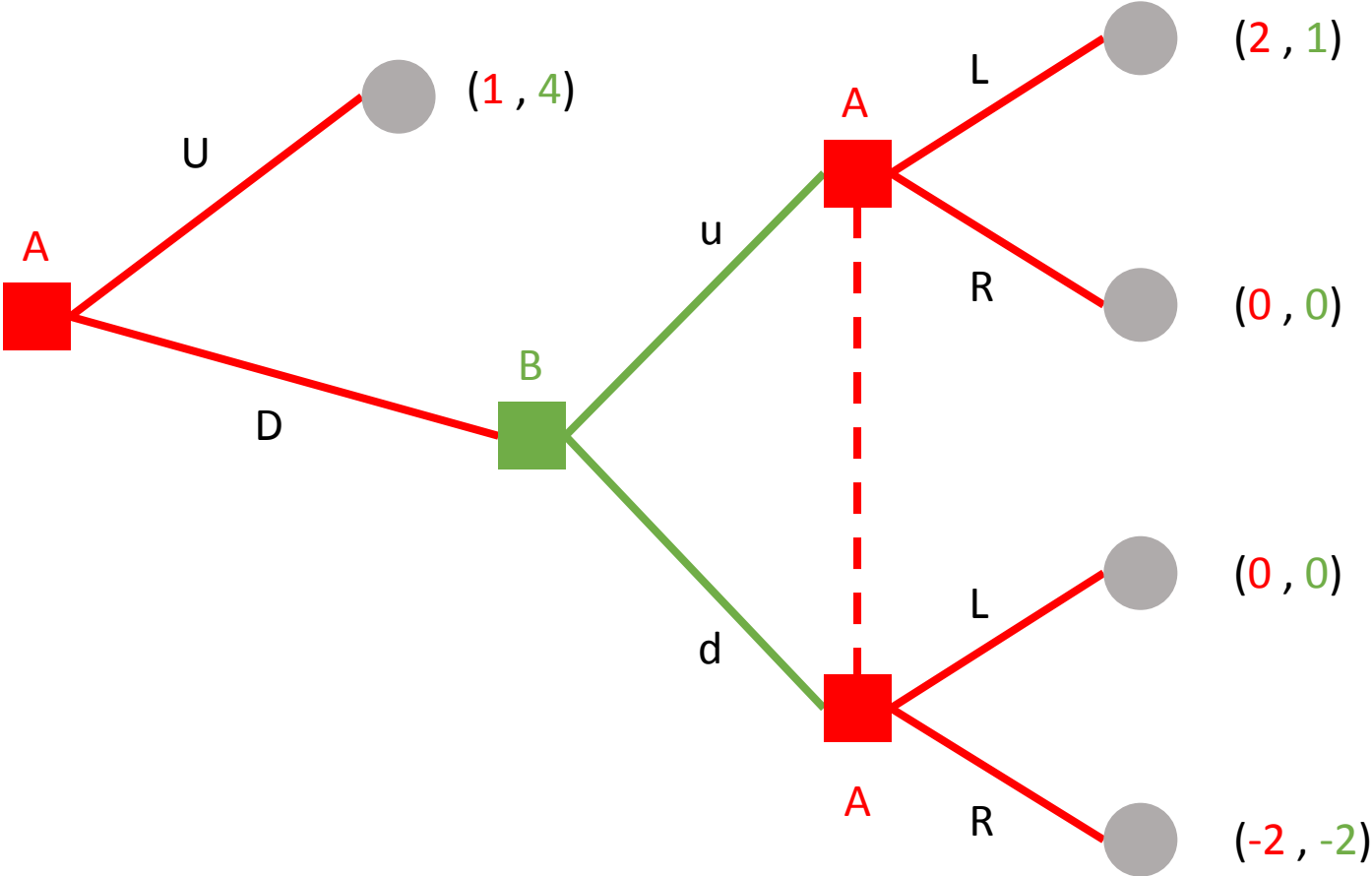
- Backwards induction NE

- NE 2:  $\langle U, R; u \rangle$
- NE 3:  $\langle D, L; u \rangle$

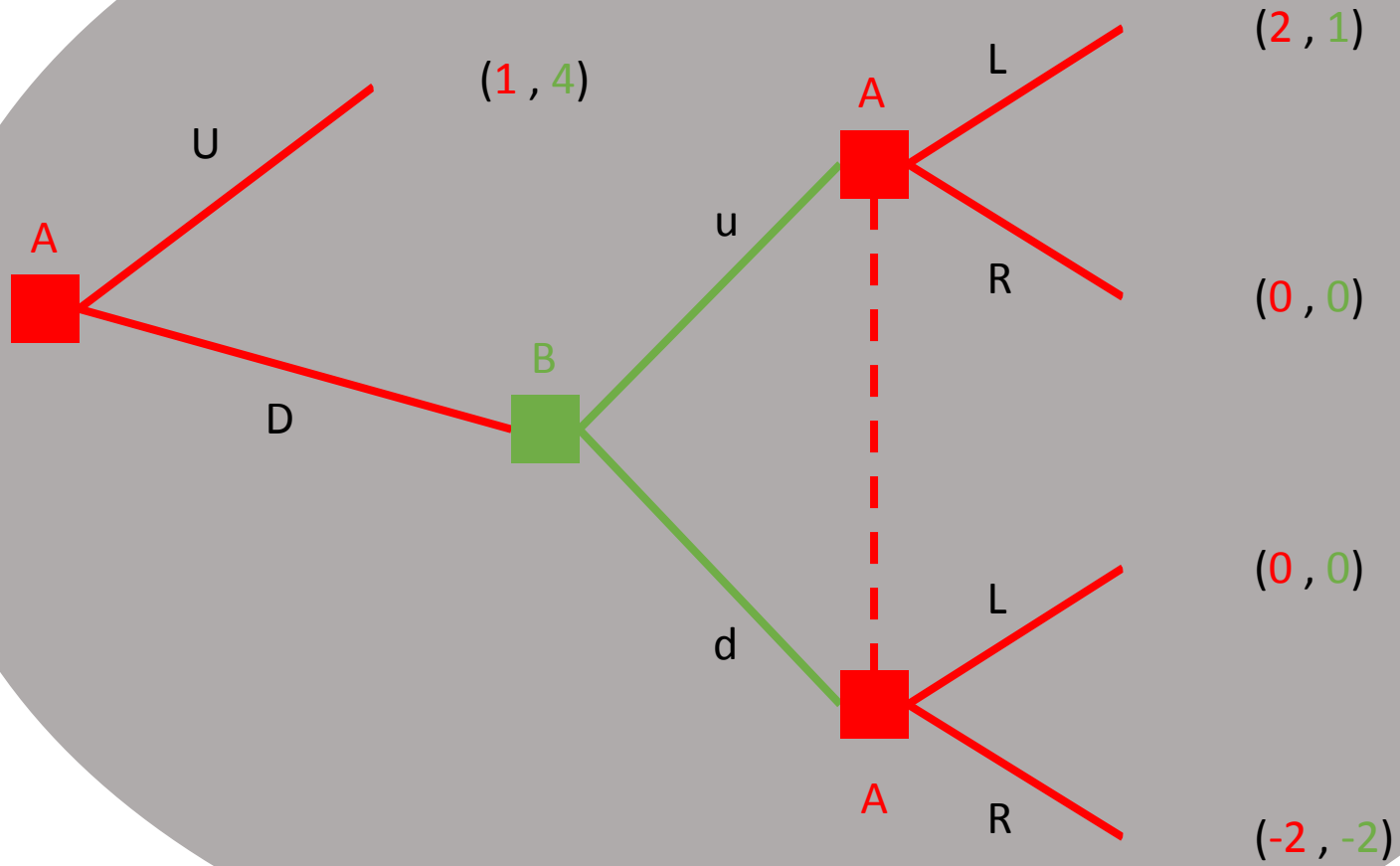
- Subgame perfect NE

- NE 3:  $\langle D, L; u \rangle$

# Game X



# Game X - Subgames



# Strategic form of Subgame of Game X

		B	
		u	d
A	L	2, 1	0, 0
	R	0, 0	-2, -2

# NE in strategic form of Subgame of Game X

B

		B	
		u	d
A	L	2, 1	0, 0
	R	0, 0	-2, -2

# Strategic form of the whole Game X

		B	
		u	d
A	UL	1 , 4	1 , 4
	UR	1 , 4	1 , 4
	DL	2 , 1	0 , 0
	DR	0 , 0	-2 , -2



# NE in strategic form of the whole Game X

B

		B	
		u	d
A	UL	1, 4	1, 4
	UR	1, 4	1, 4
	DL	2, 1	0, 0
	DR	0, 0	-2, -2

# NE in the Game X

- Strategic form NE
  - NE 1:  $\langle U, L; d \rangle$
  - NE 2:  $\langle U, R; d \rangle$
  - NE 3:  $\langle D, L; u \rangle$
- Backwards induction NE
  - Can't apply
- Subgame perfect NE
  - NE 3:  $\langle D, L; u \rangle$

Thank you for your  
attention

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