

# Game theory 1

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Can suicide terrorism  
be rational?

# Many criticisms against Rational choice

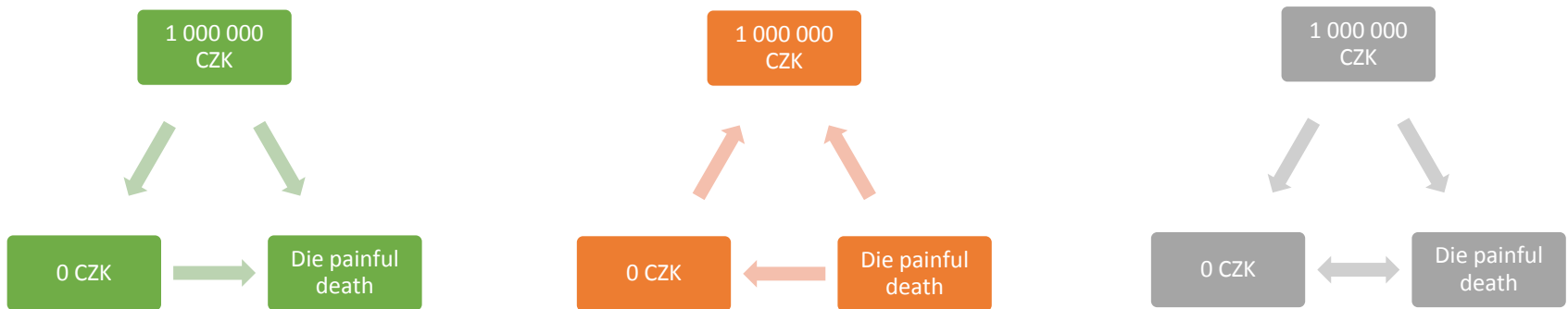
- Common criticism of rational choice – people behave irrationally
- Many times incorrect
- Rationality  $\neq$  Sensibility
- Ordering preferences
- I can mostly prefer taking over the world and least painful death, but equally prefer most painful death and least taking over the world

# Rationality

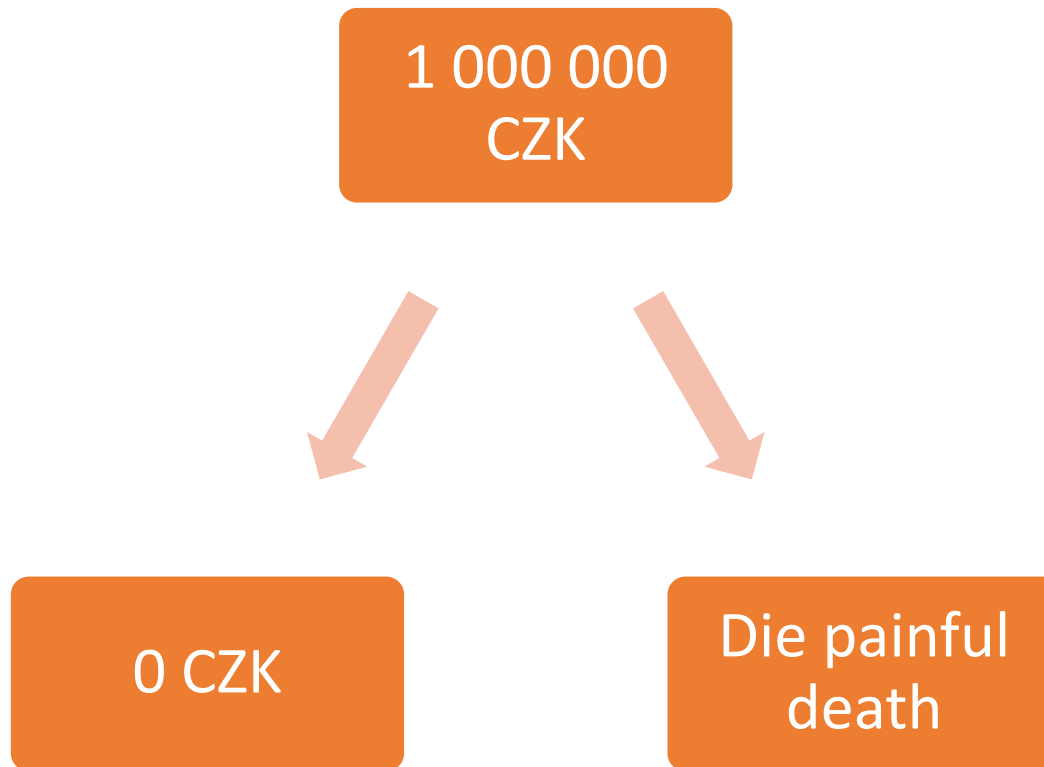
- Defined by two key premises
  - Completeness
  - Transitivity
- Indifferent to normative assessment of preferences and choices

# Completeness

- Preference ordering complete if and only if for any two outcomes X and Y individual:
  - A) Prefers X to Y – strong preference relation
  - B) Prefers Y to X – strong preference relation
  - C) Is indifferent – weak preference relation

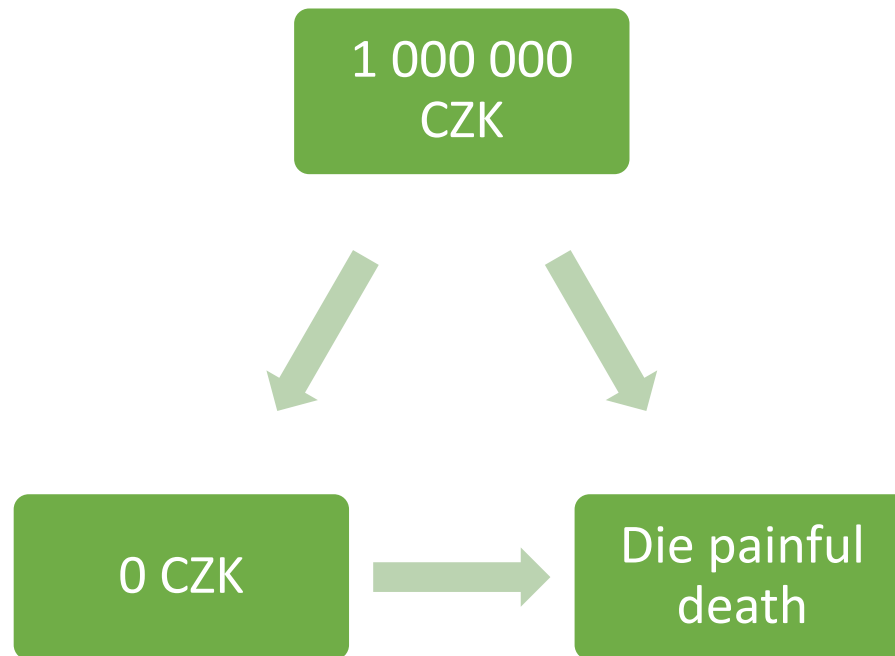


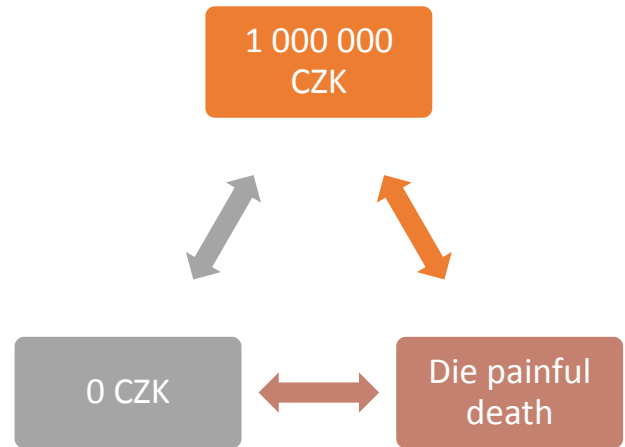
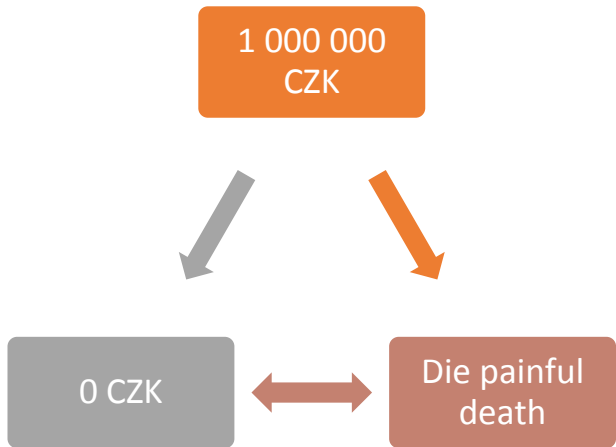
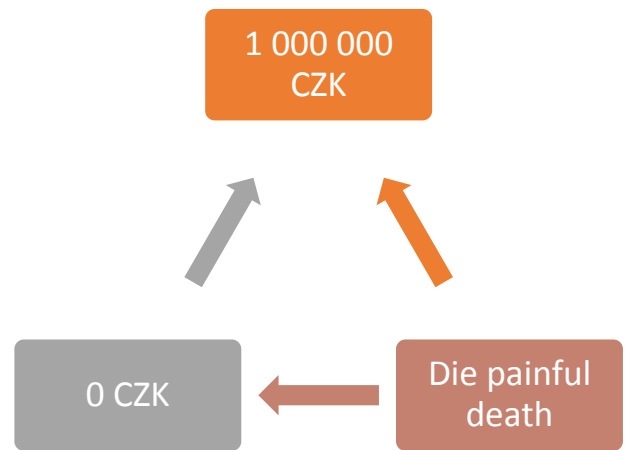
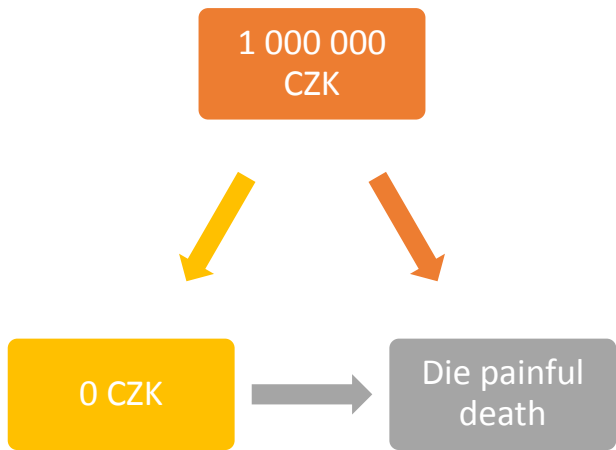
# Incomplete preferences



# Transitivity

- For any three outcomes X, Y and Z, if X is preferred to Y and Y is preferred to Z, X must be preferred to Z

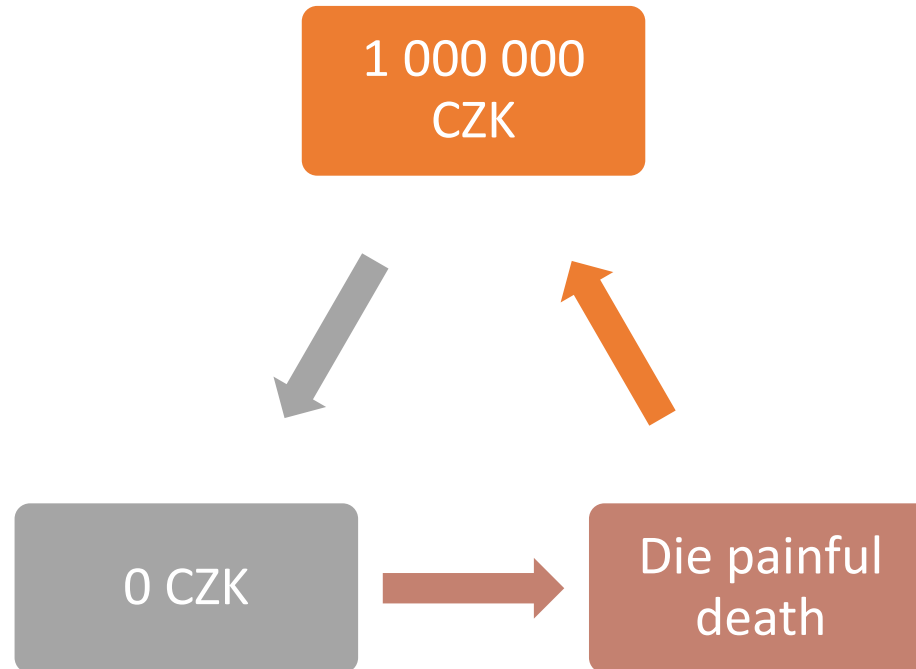






# Intransitive preferences

- Prefer X to Y, Y to Z and Z to X
- Doesn't make sense



# Other notions about preferences

- Preferences over outcomes are stable and do not change in the time of making decision – are fixed
- Preferences are ordinal – they order actions but the difference between the two values has no meaning unless they state utility
- Compare two situations
  - $u(C_1) = 1, u(C_2) = 2, u(C_3) = 0$
  - $u(C_1) = 1, u(C_2) = 200, u(C_3) = -50$
- Both situations have same preference ordering
  - $C_2 \succ C_1 \succ C_3$

# Other notions about rationality

- People do not calculate their actions – the definition of rationality is narrower than common-sense one
- Rationality tells nothing about preferences over outcomes
- Rational actors may differ in choices in same situation
- Rational actors can err

# Types of games

# Types of games

- Games of perfect information
- Games of imperfect information
  
- Cooperative games
- Non-cooperative games
  
- Constant-sum game
- Positive-sum game

# Games of perfect/imperfect information

## Perfect information games

- All players know other players' strategies available to them
- All players know payoffs over actions
- All players know other players know

## Imperfect information games

- Some information about other players' actions is not known to the player

# Cooperative/non-cooperative games

## Cooperative games

- Actors are allowed to make enforceable contracts
- Players do not need to cooperate, but cooperation is enforceable by an outside party

## Non-cooperative games

- Actors unable to make enforceable contracts outside of those specifically modeled in the game
- Players might cooperate, but any cooperation must be self-enforcing

# Constant-sum/Positive-sum games

## Constant sum games

- Sum of all players' payoffs is the same for any outcome
- Gain for one participant is always at the expense of another
- Special case of zero-sum game where all outcomes involve a sum of all player's payoffs of 0

## Positive-sum games

- Combined payoffs of all players are not the same in every outcome of the game
- Positive-sum game implies that players may have interests in common, to achieve an outcome that maximizes total payoffs.



Introducing a game

# What makes a game the game

- Players
- Actions
- Strategies
- Outcomes
- Payoffs of player

# Game of grades

- Each pair can choose 2 actions:  $\alpha$  or  $\beta$
- If both choose  $\alpha$ , both will receive C
- If both choose  $\beta$ , both will receive B
- If one chooses  $\alpha$  and other  $\beta$ , one will receive A and other D – applies to both players equally

# Game of grades – my grades

My pair

	$\alpha$	$\beta$
$\alpha$	C	A
$\beta$	D	B

Me

# Game of grades – my pair's grades

My pair

	$\alpha$	$\beta$
$\alpha$	C	D
$\beta$	A	B

Me

# Game of grades – normal form

My pair

		My pair	
		$\alpha$	$\beta$
Me	$\alpha$	C, C	A, D
	$\beta$	D, A	B, B

# Games in normal form

# Normal form representation of a game

- Called also “strategic form” or “matrix form”
- Visualized as a matrix
- Represents a game as if agents were acting simultaneously



# Utilities (Payoffs)

- Grades are not utilities
- Utilities for game:
  - $EU(A) = 3$
  - $EU(B) = 2$
  - $EU(C) = 1$
  - $EU(D) = 0$
- Preference over outcomes:  $A > B > C > D \rightarrow$   
APBPCPD

# Game of grades with payoffs

My pair

		My pair	
		$\alpha$	$\beta$
Me	$\alpha$	1, 1	3, 0
	$\beta$	0, 3	2, 2

# Solution concepts

- Nash Equilibrium
  - Dominant Strategy Equilibrium
  - Pure Strategy Equilibrium
  - Mixed Strategy Equilibrium
- Subgame Perfect Equilibrium
- Bayesian Equilibrium
- Weak Perfect Bayesian Equilibrium

My pair

	$\alpha$	$\beta$
$\alpha$	1, 1	3, 0
$\beta$	0, 3	2, 2

Me

My pair

	$\alpha$	$\beta$
$\alpha$	1, 1	3, 0
$\beta$	0, 3	2, 2

Me

My pair

		$\alpha$	$\beta$
Me	$\alpha$	1, 1	3, 0
	$\beta$	0, 3	2, 2

My pair

	$\alpha$	$\beta$
$\alpha$	C, C	A, D
$\beta$	D, A	B, B

Me

# Prisoner's dilemma

- Both players are tempted to defect, since cooperate is strictly dominated by defect
- The outcome of the game is that both players betray the other one and end up choosing  $\alpha$
- Both will end up with outcome that is less preferred than the optimal outcome  $\beta, \beta$  by seeking maximal gain from own action
- $\beta, \beta$  is Pareto Efficient outcome – brings best outcomes for all players – no one could be better-off without making someone worse-off



**Dominance**

# Dominant Strategy Equilibrium

- Strategy might be dominant

Two types of dominance

- Strict (strong) dominance
- Weak dominance

# Strict dominance

- Player  $i$
- Payoff  $u_i$
- Dominant strategy  $s_i$
- Dominated strategy  $s_i'$
- Strategy of all other players  $s_{-i}$
  
- Player  $i$ 's strategy  $s_i'$  is strictly dominated by player  $i$ 's strategy  $s_i$  if and only if
- $u_i(s_i, s_{-i}) > u_i(s_i', s_{-i})$  for **all**  $s_{-i}$
  
- utility of playing  $s_i$  against others' strategies  $s_{-i}$  is **greater** than utility of playing  $s_i'$  against others's strategies  $s_{-i}$  for all others' strategies  $s_{-i}$

# Game of grades – strict dominance

My pair

		My pair	
		$\alpha$	$\beta$
Me	$\alpha$	1, 1	3, 0
	$\beta$	0, 3	2, 2

# Weak dominance

- Player  $i$
- Payoff  $u_i$
- Dominant strategy  $s_i$
- Dominated strategy  $s_i'$
- Strategy of all other players  $s_{-i}$
  
- Player  $i$ 's strategy  $s_i'$  is weakly dominated by player  $i$ 's strategy  $s_i$  if
- $u_i(s_i, s_{-i}) \geq u_i(s_i', s_{-i})$  for **all**  $s_{-i}$  and
- $u_i(s_i, s_{-i}) > u_i(s_i', s_{-i})$  for **some**  $s_{-i}$
  
- utility of playing  $s_i$  against others' strategies  $s_{-i}$  is **greater or equal to** utility of playing  $s_i'$  against others' strategies  $s_{-i}$  for all others' strategies  $s_{-i}$  and **greater for some** others' strategies  $s_{-i}$

# Game of grades – weak dominance

My pair

		My pair	
		$\alpha$	$\beta$
Me	$\alpha$	1, 1	3, 0
	$\beta$	0, 3	3, 2

Never play dominated  
strategies

- Dominated strategy **brings lesser payoffs** than dominant strategy
- Dominated strategy brings lesser payoffs **no matter what strategy is selected by other player**
- Can't control minds of others to force them not to play dominant strategy
- Even if **could** control minds of others and be sure they'll play dominated strategy, than **rational to play dominant strategy anyway**



# Choosing numbers

- Choose integer between 1 – 100 incl.
- All numbers will be averaged
- Winner is the one who will be closest to the  $\frac{2}{3}$  of the group's average

# Choosing numbers

- Average = 100
- $2/3$  of average =  $\sim 66.66$
- $X > 67$  is strictly dominated strategy
  - Even if everyone else selected 100
  - One selected 67
  - I selected 68
  - Outcome – 68 is dominated by 67
- What is the rational choice for this game?

If all players were  
strictly rational, result  
is 1

# I know you know

- I know
  - Numbers above 67 are never rational
- You know that I know
  - You'll never select number above 67, therefore numbers above 46 are never rational either
- I know You know that I know
  - I know that You'll never select above 46, hence I should never select number higher than 30
- You know that I know that You know that I know
  - You know that I won't select above 30, therefore I should never select number above 20

Get into opponent's  
shoes

# Real life results

- 2012 Game theory online course
- 10 000 + players
- Mean 34
- Mode 50
- Median 33
- Winner 23
- Spikes: 50, 33, 20, 1

# Iterated deletion of dominated strategies

# Iterated deletion of dominated strategies

- Can delete dominated strategies as if they were not present in the game
- Game becomes simpler than the original one
- Can find equilibriums quickly – games are dominance-solvable



# Game of grades

My pair

		$\alpha$	$\beta$
Me	$\alpha$	1, 1	<del>3</del> , 0
	$\beta$	0, 3	<del>2</del> , 2

My pair

$\alpha$

$\alpha$

1, 1

Me

$\beta$

0, 3

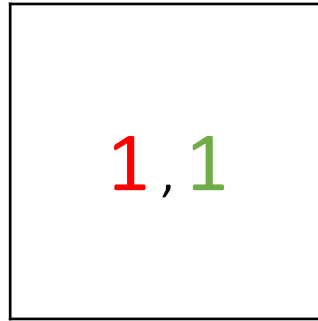
My pair

$\alpha$

Me

$\alpha$

1, 1



This game is  
dominance-solvable

# Opponent

		Opponent		
		$s_1$	$s_2$	$s_3$
Me	$S_1$	0, 1	-2, 3	4, -1
	$S_2$	0, 3	3, 1	6, 4
	$S_3$	1, 5	4, 2	5, 2

# $S_1$ VS $S_2$

Opponent

		$S_1$	$S_2$	$S_3$
Me	$S_1$	$0, 1$	$-2, 3$	$4, -1$
	$S_2$	$0, 3$	$3, 1$	$6, 4$
	$S_3$	$1, 5$	$4, 2$	$5, 2$

# $S_1$ VS $S_3$

Opponent

		$S_1$	$S_2$	$S_3$
Me	$S_1$	0, 1	-2, 3	4, -1
	$S_2$	0, 3	3, 1	6, 4
	$S_3$	1, 5	4, 2	5, 2

# $S_2$ VS $S_3$

Opponent

		$S_1$	$S_2$	$S_3$
Me	$S_1$	0, 1	-2, 3	4, -1
	$S_2$	0, 3	3, 1	<b>6</b> , 4
	$S_3$	<b>1</b> , 5	<b>4</b> , 2	<b>5</b> , 2



# $S_1$ VS $S_3$

Opponent

		$S_1$	$S_2$	$S_3$
Me	$S_1$	0, <b>1</b>	-2, 3	4, -1
	$S_2$	0, 3	3, 1	6, <b>4</b>
	$S_3$	1, <b>5</b>	4, 2	5, 2

# $S_1$ VS $S_2$

Opponent

		$S_1$	$S_2$	$S_3$
Me	$S_1$	0, 1	-2, 3	4, -1
	$S_2$	0, 3	3, 1	6, 4
	$S_3$	1, 5	4, 2	5, 2

# $S_2$ VS $S_3$

Opponent

		$S_1$	$S_2$	$S_3$
Me	$S_1$	0, 1	-2, <b>3</b>	4, -1
	$S_2$	0, 3	3, <b>1</b>	6, <b>4</b>
	$S_3$	1, 5	4, <b>2</b>	5, <b>2</b>

# Opponent

		$s_1$	$s_2$	$s_3$
Me	$s_1$	<del>0, 1</del>	<del>-2, 3</del>	<del>4, -1</del>
	$s_2$	0, 3	3, 1	6, 4
	$s_3$	1, 5	4, 2	5, 2

Opponent

		Opponent		
		$s_1$	$s_2$	$s_3$
Me	$s_2$	0, 3	3, 1	6, 4
	$s_3$	1, 5	4, 2	5, 2

# $s_1$ vs $s_3$ after deletion

Opponent

		$s_1$	$s_2$	$s_3$
Me	$s_2$	0, 3	3, 1	6, 4
	$s_3$	1, 5	4, 2	5, 2

# $s_1$ vs $s_2$ after deletion

Opponent

		$s_1$	$s_2$	$s_3$
Me	$s_2$	0, 3	3, 1	6, 4
	$s_3$	1, 5	4, 2	5, 2

# $s_2$ vs $s_3$ after deletion

Opponent

		$s_1$	$s_2$	$s_3$
Me	$s_2$	0, 3	3, 1	6, 4
	$s_3$	1, 5	4, 2	5, 2



Opponent

		$s_1$	$s_2$	$s_3$
Me	$s_2$	0, 3	<del>3, 1</del>	6, 4
	$s_3$	1, 5	<del>4, 2</del>	5, 2

Opponent

	$s_1$	$s_3$
$s_2$	0, 3	6, 4
$s_3$	1, 5	5, 2

Me

# Opponent

	$s_1$	$s_3$
$s_2$	0, 3	6, 4
$s_3$	1, 5	5, 2

Me

Opponent

		Opponent	
		$s_1$	$s_3$
Me	$s_2$	0, 3	6, 4
	$s_3$	1, 5	5, 2

Sometimes not  
solvable,  
but simplified

# Limits of iterated deletion of dominated strategies

- Strictly dominated strategies may be deleted in a random order
- Deleting weakly dominated strategies in some order might delete equilibriums
- This solution concept is not always applicable – sometimes game simply don't have dominance

# How to solve the game without dominance?

Opponent

		Opponent	
		$s_1$	$s_3$
Me	$s_2$	0, 3	6, 4
	$s_3$	1, 5	5, 2

# How to solve the game without dominance?

Opponent

		Opponent	
		$s_1$	$s_3$
Me	$s_2$	0 , 3	6 , 4
	$s_3$	1 , 5	5 , 2



# Nash Equilibrium

# Nash Blonde Game

- 2 or more lusty males
- Several interested females
- At least one more female than male
- Just one female blonde
- Every male prefers blonde to brunette and brunette to no companion

# Nash Blonde Game – normal form

		M2	
		Bl	Br
M1	Bl	0, 0	2, 1
	Br	1, 2	1, 1

# Nash Equilibrium

- Set of strategies, one for each player, such that **no player has incentive to unilaterally change** her action
- Players are in equilibrium if a change in strategies by any one of them would lead player to earn less (considering strategies of others') than if she remained with her current strategy
- **Mutual best response** to others' choices

		A		
		L	C	R
B	T	1, 1	0, 0	0, 0
	M	0, 2	1, 1	2, -1
	B	0, 0	1, 2	2, 1

		A		
		L	C	R
B	T	1, 1	0, 0	0, 0
	M	0, 2	1, 1	2, -1
	B	0, 0	1, 2	2, 1

		A		
		L	C	R
B	T	1, 1	0, 0	0, 0
	M	0, 2	1, 1	2, -1
	B	0, 0	1, 2	2, 1

		A		
		L	C	R
B	T	1, 1	0, 0	0, 0
	M	0, 2	1, 1	2, -1
	B	0, 0	1, 2	2, 1



Games might have  
more NE

# Pure strategy equilibrium

- Two equilibriums in this game
- ( T , L )
  - $u(A) = 1$
  - $u(B) = 1$
- ( C , B )
  - $u(A) = 1$
  - $u(B) = 2$
- These are **pure strategy equilibriums**

Other basic games

# Chicken

A

B

		s	h
S	5, 5	0, 10	
H	10, 0	-10, -10	

Detailed description: A 2x2 normal form game matrix for the 'Chicken' game. The row player is labeled 'A' and the column player is labeled 'B'. Player A's strategies are 'S' (Swerve) and 'H' (Hold). Player B's strategies are 's' (Swerve) and 'h' (Hold). The payoffs are given as (A's payoff, B's payoff). The matrix shows that if both players swerve, they both get 5. If one swerves and the other holds, the swerver gets 0 and the holder gets 10. If both hold, both get -10.

# Chicken NE

- Pure strategies NE

- ( **H** , **s** )

- $EU(A) = 10$

- $EU(B) = 0$

- ( **S** , **h** )

- $EU(A) = 0$

- $EU(B) = 10$

A

B

		s	h
S	5, 5	0, 10	
H	10, 0	-10, -10	

# Stag hunt

A

B

		B	
		s	r
A	S	5, 5	0, 3
	R	3, 0	3, 3

# Stag hunt NE

- Pure strategies NE

- ( S , s )

- $EU(A) = 5$

- $EU(B) = 5$

- ( R , r )

- $EU(A) = 3$

- $EU(B) = 3$

A

B

		s	R
S	5, 5	0, 3	
R	3, 0	3, 3	