# Game theory 1

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# Can rational choice explain suicide terrorism?

### Many criticisms against Rational choice

- Common criticism of rational choice people behave irrationally
- Many times incorrect
- Rationality ≠ Sensibility
- Ordering preferences
- I can mostly prefer taking over the world and least painful death, but equally prefer most painful death and least taking over the world

### Rationality

- Defined by two key premises
  - Completeness
  - Transitivity
- Indifferent to normative assessment of preferences and choices

### Completeness

- Preference ordering complete if and only if for any two outcomes X and Y individual:
  - A) Prefers X to Y strong preference relation
  - B) Prefers Y to X strong preference relation
  - C) Is indifferent weak preference relation



Incomplete preferences



### Transitivity

• For any three outcomes X, Y and Z, if X is preferred to Y and Y is preferred to Z, X must be preferred to Z











### Intransitive preferences

- Prefer X to Y, Y to Z and Z to X
- Doesn't make sense



### Other notions about preferences

- Preferences over outcomes are stable and do not change in the time of making decision – are fixed
- Preferences are ordinal they order actions but the difference between the two values has no meaning unless they state utility
- Compare two situations
  - u(C<sub>1</sub>) = 1, u(C<sub>2</sub>) = 2, u(C<sub>3</sub>) = 0
  - u(C<sub>1</sub>) = 1, u(C<sub>2</sub>) = 200, u(C<sub>3</sub>) = -50
- Both situations have same preference ordering
  - C<sub>2</sub> p C<sub>1</sub> p C<sub>3</sub>

### Other notions about rationality

- Rational choice theory is not attempting to explain cognitive processes happening in individuals
- Rationality tells nothing about preferences over outcomes
- Rational actors may differ in choices in same situation
- Rational actors can err

# Types of games

# Types of games

- Games of perfect information
- Games of imperfect information
- Cooperative games
- Non-cooperative games
- Constant-sum game
- Positive-sum game

# Games of perfect/imperfect information

#### Perfect information games

- All players know other players' strategies available to them
- All players know payoffs over actions
- All players know other players know

### Imperfect information games

 Some information about other players' actions is not know to the player

### Cooperative/non-cooperative games

#### **Cooperative games**

- Actors are allowed to make enforceable contracts
- Players do not need to cooperate, but cooperation is enfoceable by an outside party

#### Non-cooperative games

- Actors unable to make enforceable contracts outside of those specifically modeled in the game
- Players might cooperate, but any cooperation must be selfenforcing

### Constant-sum/Positive-sum games

#### **Constant sum games**

- Sum of all players' payoffs is the same for any outcome
- Gain for one participant is always at the expense of another
- Special case of zero-sum game where all outcomes involve a sum of all player's payoffs of 0

#### Positive-sum games

- Combined payoffs of all players are not the same in every outcome of the game
- Positive-sum game implies that players may have interests in common, to achieve an outcome that maximizes total payoffs.

# Introducing a game

### What makes a game the game

- Players
- Actions
- Strategies
- Outcomes
- Payoffs of player

### Game of grades

- Each pair can choose 2 actions:  $\alpha$  or  $\beta$
- If both choose  $\alpha$ , both will receive C
- If both choose  $\beta$ , both will receive B
- If one chooses  $\alpha$  and other  $\beta$ , one will receive A and other D

### Game of grades – my grades



### Game of grades – my opponent's grades



# Game of grades – normal form

Me

	α	β
α	<mark>C</mark> ,C	<b>A</b> , D
β	<b>D</b> , A	В,В

# Games in normal form

### Normal form representation of a game

- Called also "strategic form" or "matrix form"
- Visualized as a matrix
- Represents a game as if agents were acting simultaneously

### Utilities (Payoffs)

- Grades are not utilites
- Utilities for game:
  - EU(A) = 3
  - EU(B) = 2
  - EU(C) = 1
  - EU(D) = 0
- Preference over outcomes: A > B > C > D -> APBPCPD

### Game of grades with payoffs

Me

	α	β
α	1,1	<mark>3</mark> ,0
β	<mark>0</mark> ,3	<mark>2</mark> ,2

### Solution concepts

- Nash Equilibrium
  - Dominant Strategy Equilibrium
  - Pure Strategy Equilibrium
  - Mixed Strategy Equilibrium
- Subgame Perfect Equilibrium
- Bayesian Equilibrium
- Weak Perfect Bayesian Equilibrium



Me





Me





### My opponent



Me

### Prisoner's dilemma

- Both players are tempted to defect, since cooperate is strictly dominated by defect
- The outcome of the game is that both players betray the other one and end up choosing  $\boldsymbol{\alpha}$
- Both will end up with outcome that is less preferred than the optimal outcome  $\beta$ ,  $\beta$  by seeking maximal gain from own action
- β, β is Pareto Efficient outcome brings best outcomes for all players – no one could be better-off without making someone worse-off

# Dominance

### Dominant Strategy Equilibrium

• Strategy might be dominant

Two types of dominance

- Strict (strong) dominance
- Weak dominance

### Strict dominance

- Player i
- Payoff u<sub>i</sub>
- Dominant strategy s<sub>i</sub>
- Dominated strategy s<sub>i</sub>'
- Strategy of all other players s<sub>-i</sub>
- Player i's strategy si' is strictly dominated by player i's strategy si if and only if
- $u_i(s_i, s_{-i}) > u_i(s_i', s_{-i})$  for **all**  $s_{-i}$
- utility of playing s<sub>i</sub> against others' strategies s<sub>-i</sub> is greater than utility of playing s<sub>i</sub>' against others's strategies s<sub>-i</sub> for all others' strategies s<sub>-i</sub>

### Game of grades – strict dominance



Me
#### Weak dominance

- Player i
- Payoff u<sub>i</sub>
- Dominant strategy s<sub>i</sub>
- Dominated strategy s<sub>i</sub>'
- Strategy of all other players s<sub>-i</sub>
- Player i's strategy si' is weakly dominated by player i's strategy si if
- $u_i(s_i, s_{-i}) \ge u_i(s_i', s_{-i})$  for **all**  $s_{-i}$  and
- $u_i(s_i, s_{-i}) > u_i(s_i', s_{-i})$  for **some**  $s_{-i}$
- utility of playing s<sub>i</sub> against others' strategies s<sub>-i</sub> is greater or equal to utility of playing s<sub>i</sub>' against others's strategies s<sub>-i</sub> for all others' strategies s<sub>-i</sub> and greater for some others' strategies s<sub>-i</sub>

#### Game of grades – weak dominance



Me

Never play dominated strategies

- Dominated strategy brings lesser payoffs than dominant strategy
- Dominated strategy brings lesser payoffs no matter what strategy is selected by other player
- Can't control minds of others to force them not to play dominant strategy
- Event if **could** control minds of others and be sure they'll play dominated strategy, than **rational to play dominant strategy anyway**

## Choosing numbers

- Choose integer between 1 100 incl.
- All numbers will be averaged
- Winner is the one who will be closest to the 2/3 of the group's average

#### Choosing numbers

- Average = 100
- 2/3 of average = ~ 66.66
- X > 67 is strictly dominated strategy
  - Even if everyone else selected 100
  - One selected 67
  - I selected 68
  - Outcome 68 is dominated by 67
- What is the rational choice for this game?

# If all players were strictly rational, result is 1

### I know you know

- I know
  - Numbers above 67 are never rational
- You know that I know
  - You'll never select number above 67, therefore numbers above 46 are never rational either
- I know You know that I know
  - I know that You'll never select above 46, hence I should never select number higher than 30
- You know that I know that You know that I know
  - You know that I won't select above 30, therefore I should never select number above 20

Get into opponent's shoes

#### Real life results

- 2012 Game theory online course
- 10 000 + players
- Mean 34
- Mode 50
- Median 33
- Winner 23
- Spikes: 50, 33, 20, 1

Iterated deletion of dominated strategies

## Iterated deletion of dominated strategies

- Can delete dominated strategies as if they were not present in the game
- Game becomes simpler than the original one
- Can find equilibriums quickly games are dominance-solvable

## Game of grades





#### My pair



#### My pair

Me



This game is dominance-solvable

#### Opponent



 $S_1 vs S_2$ 

		S <sub>1</sub>	s <sub>2</sub>	S <sub>3</sub>
Me	S <sub>1</sub>	<mark>0</mark> ,1	- <mark>2</mark> , 3	4,-1
	S <sub>2</sub>	<mark>0</mark> ,3	<b>3</b> ,1	<mark>6</mark> ,4
	S <sub>3</sub>	1,5	4 , 2	5,2

 $S_1 vs S_3$ 

		S <sub>1</sub>	s <sub>2</sub>	S <sub>3</sub>
	S <sub>1</sub>	<mark>0</mark> ,1	- <mark>2</mark> , 3	4,-1
Me	S <sub>2</sub>	0,3	3,1	6,4
	S <sub>3</sub>	<b>1</b> ,5	4,2	<mark>5</mark> , 2

 $S_2 vs S_3$ 

		S <sub>1</sub>	S <sub>2</sub>	S <sub>3</sub>
Me	S <sub>1</sub>	0,1	-2 , 3	4,-1
	S <sub>2</sub>	<mark>0</mark> ,3	<mark>3</mark> ,1	<mark>6</mark> ,4
	S <sub>3</sub>	<b>1</b> ,5	<b>4</b> , 2	5,2

 $\mathbf{S}_1~\mathbf{VS}~\mathbf{S}_3$ 

		S <sub>1</sub>	s <sub>2</sub>	S <sub>3</sub>
	S <sub>1</sub>	o, <b>1</b>	-2,3	4 , -1
Me	S <sub>2</sub>	0,3	3,1	6 <b>, 4</b>
	S <sub>3</sub>	1,5	4 , 2	5,2

 $\mathbf{S}_1 \; \mathbf{VS} \; \mathbf{S}_2$ 

		S <sub>1</sub>	s <sub>2</sub>	S <sub>3</sub>
Me	S <sub>1</sub>	0,1	-2 , 3	4,-1
	S <sub>2</sub>	o , <b>3</b>	3,1	6,4
	S <sub>3</sub>	1,5	4,2	5,2

 $S_2 VS S_3$ 

		S <sub>1</sub>	s <sub>2</sub>	S <sub>3</sub>
	<b>S</b> <sub>1</sub>	0,1	-2,3	4,-1
Me	S <sub>2</sub>	0,3	3,1	6 <b>, 4</b>
	<b>S</b> <sub>3</sub>	1,5	4 , <mark>2</mark>	5 , <b>2</b>





#### Opponent

	S <sub>1</sub>	S <sub>2</sub>	S <sub>3</sub>
S <sub>2</sub>	<b>0</b> ,3	<b>3</b> ,1	<mark>6</mark> , 4
S <sub>3</sub>	<mark>1</mark> , 5	4,2	5,2

Me

#### $s_1 vs s_3$ after deletion





## $s_1 vs s_2$ after deletion





#### $s_2$ vs $s_3$ after deletion







Me

#### Opponent

## $S_1$ S<sub>3</sub> S<sub>2</sub> <mark>0</mark>,3 <mark>6</mark>,4 S<sub>3</sub> 5,2 1,5

Opponent

Me



#### Opponent

## $S_1$ S<sub>3</sub> S<sub>2</sub> 0,3 6,**4** S<sub>3</sub> 5,<mark>2</mark> 1,5

Opponent

Me

## Sometimes not solvable, but simplified

## Limits of iterated deletion of dominated strategies

- Strictly dominated strategies may be deleted in a random order
- Deleting weakly dominated strategies in some order might delete equilibriums
- This solution concept is not always applicable sometimes game simply don't have dominance

#### How to solve the game without dominance?





#### How to solve the game without dominance?




# Nash Equilibrium

#### Nash Blonde Game

- 2 or more lusty males
- Several interested females
- At least one more female than male
- Just one female blonde
- Every male prefers blonde to brunette and brunette to no companion

#### Nash Blonde Game – normal form



M2

## Nash Equilibrium

- Set of strategies, one for each player, such that no player has incentive to unilaterally change her action
- Players are in equilibrium if a change in strategies by any one of them would lead player to earn less (considering strategies of others') than if she remained with her current strategy
- Mutual best response to others' choices

		L	С	R
	т	1,1	<mark>0</mark> ,0	<mark>0</mark> ,0
В	М	<mark>0</mark> ,2	<mark>1</mark> ,1	<mark>2</mark> ,-1
	В	<mark>0</mark> ,0	<mark>1</mark> ,2	<mark>2</mark> ,1







# Games might have more NE

### Pure strategy equilibrium

- Two equilibriums in this game
- ( T , L )
  - u(A) = 1
  - u(B) = 1
- ( <mark>C</mark> , B )
  - u(A) = 1
  - u(B) = 2
- These are pure strategy equilibriums