

# Game theory 2

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# Sum-up of the previous lecture

Opponent

		$s_1$	$s_3$
Me	$s_2$	0 , 3	6 , 4
	$s_3$	1 , 5	5 , 2

Social welfare

# Social welfare

- Situation where **sum of all payoffs** of an outcome is at its **maximum**
- Might lead to rationally unstable solutions
- Does not provide a solid analytical tool

# Game M

		B	
		Right	Left
A	Right	2 , 2	0 , 0
	Left	0 , 0	1 , 1

# Game M – Social welfare

B

		Right	Left
A	Right	4	0
	Left	0	2

# Game N

		B	
		l	r
A	L	2, 2	4, 0
	R	2, 3	8, -1

# Game N – Social welfare

B

		B	
		l	r
A	L	4	4
	R	5	7



# Prisoner's dilemma – Social welfare

B

		B	
		c	d
A	C	5, 5	0, 7
	D	7, 0	3, 3

# Prisoner's dilemma – Social welfare

B

		c	d
A	C	10	7
	D	7	6

# Prisoner's dilemma – Social welfare

		B	
		c	d
A	C	5, 50	0, 70
	D	7, 0	3, 30

A 2x2 payoff matrix for a Prisoner's dilemma. The rows represent Player A's strategies (C and D) and the columns represent Player B's strategies (c and d). The payoffs are shown as (A's payoff, B's payoff). The cell for (D, d) is highlighted in yellow.

A \ B	c	d
C	5, 50	0, 70
D	7, 0	3, 30

# Prisoner's dilemma – Social welfare

A

		B	
		c	d
C	C	55	70
	D	70	33

# Prisoner's dilemma – Social welfare

		B	
		C	d
A	C	50, 5	0, 7
	D	70, 0	30, 3

# Prisoner's dilemma – Social welfare

B

		C	D
A	C	55	7
	D	70	33

Pareto efficiency

# Game M

		B	
		Right	Left
A	Right	2 , 2	0 , 0
	Left	0 , 0	1 , 1



# Game M

		B	
		Right	Left
A	Right	<b>2</b> , 2	0, 0
	Left	0, 0	<b>1</b> , 1

# Game M – pure strategy equilibriums

B

		B	
		Right	Left
A	Right	<b>2</b> , <b>2</b>	<b>0</b> , <b>0</b>
	Left	<b>0</b> , <b>0</b>	<b>1</b> , <b>1</b>

# Pareto efficiency

- Outcome is Pareto efficient (Pareto optimal), if there is **no other outcome** which is **better or equal for all players** and **strictly better for some player**
- Conversely, outcome A is **Pareto dominated**, if there is outcome B that makes **all players as good** (weakly better) and **one player strictly better** compared to outcome A
- **Pareto dominated** outcome is **not Pareto efficient**
- Might lead to rationally unstable solutions

# Game M – Pareto efficiency

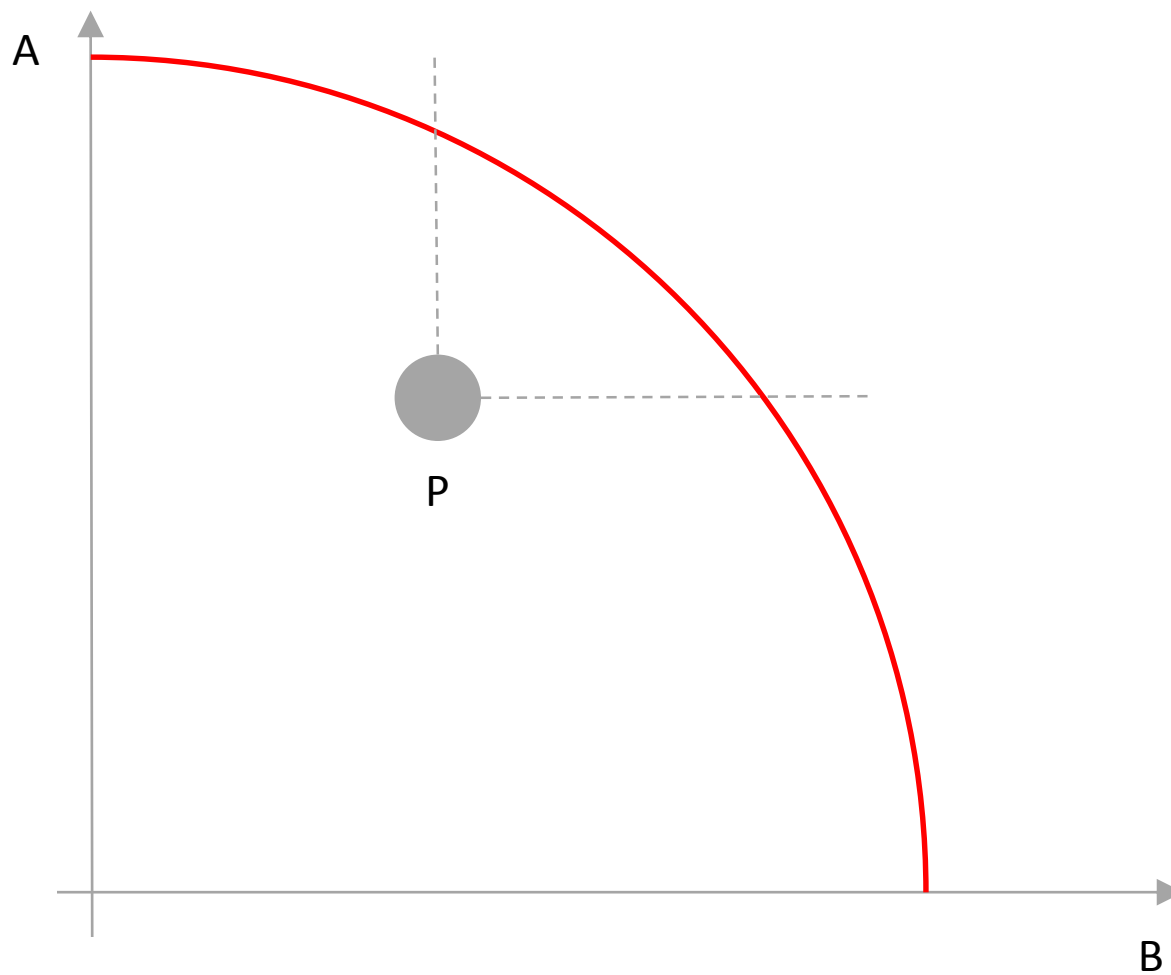
B

		Right	Left
A	Right	2, 2	0, 0
	Left	0, 0	1, 1

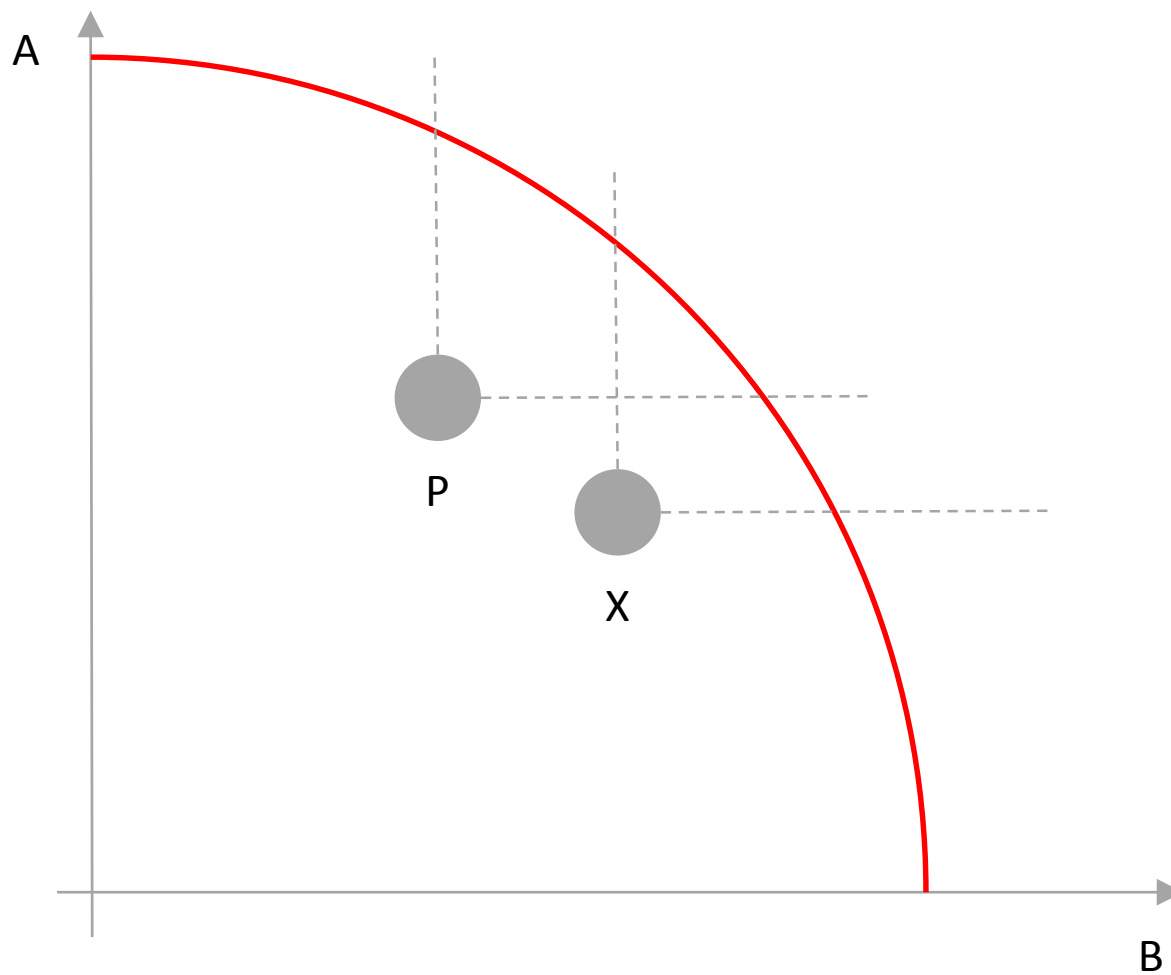
# Pareto efficiency



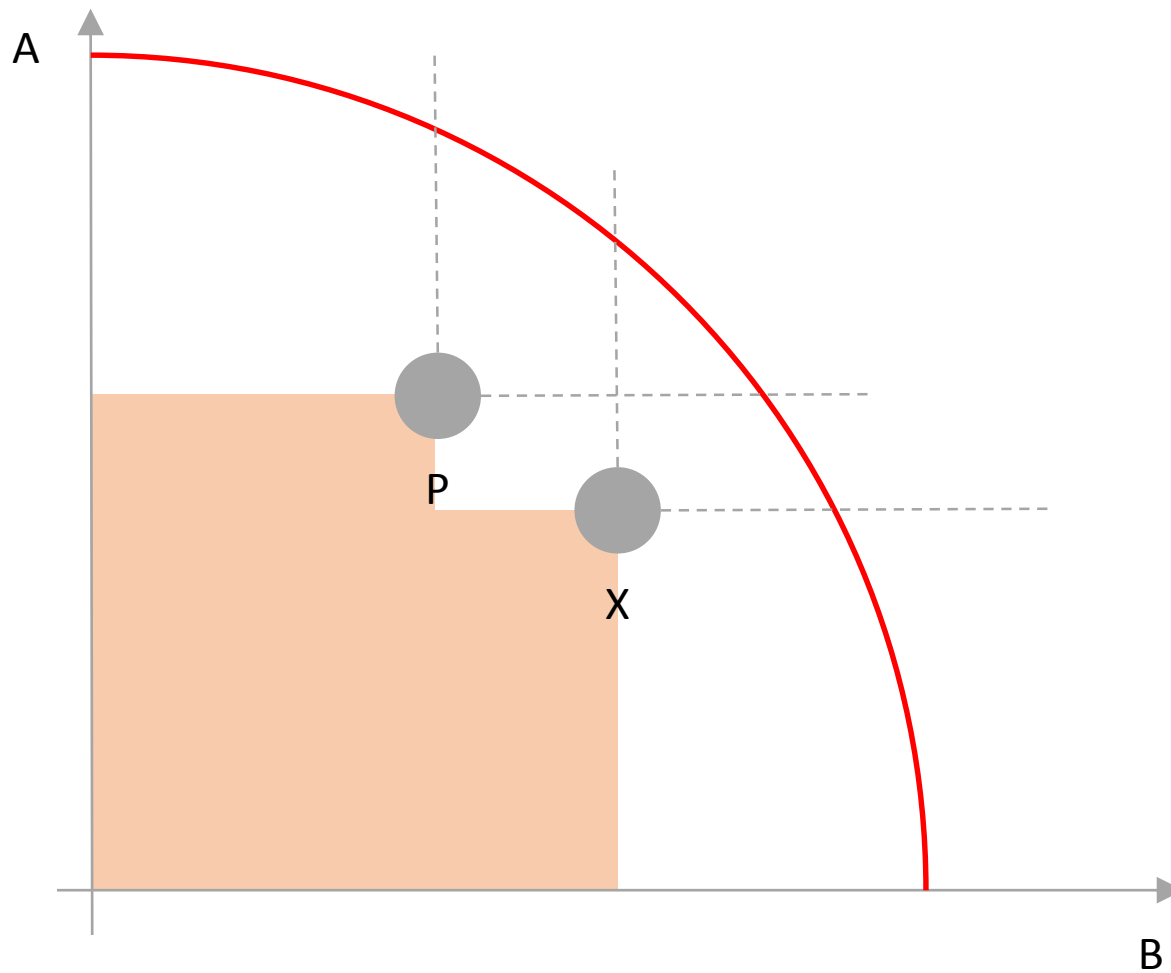
# Pareto efficiency



# Pareto efficiency

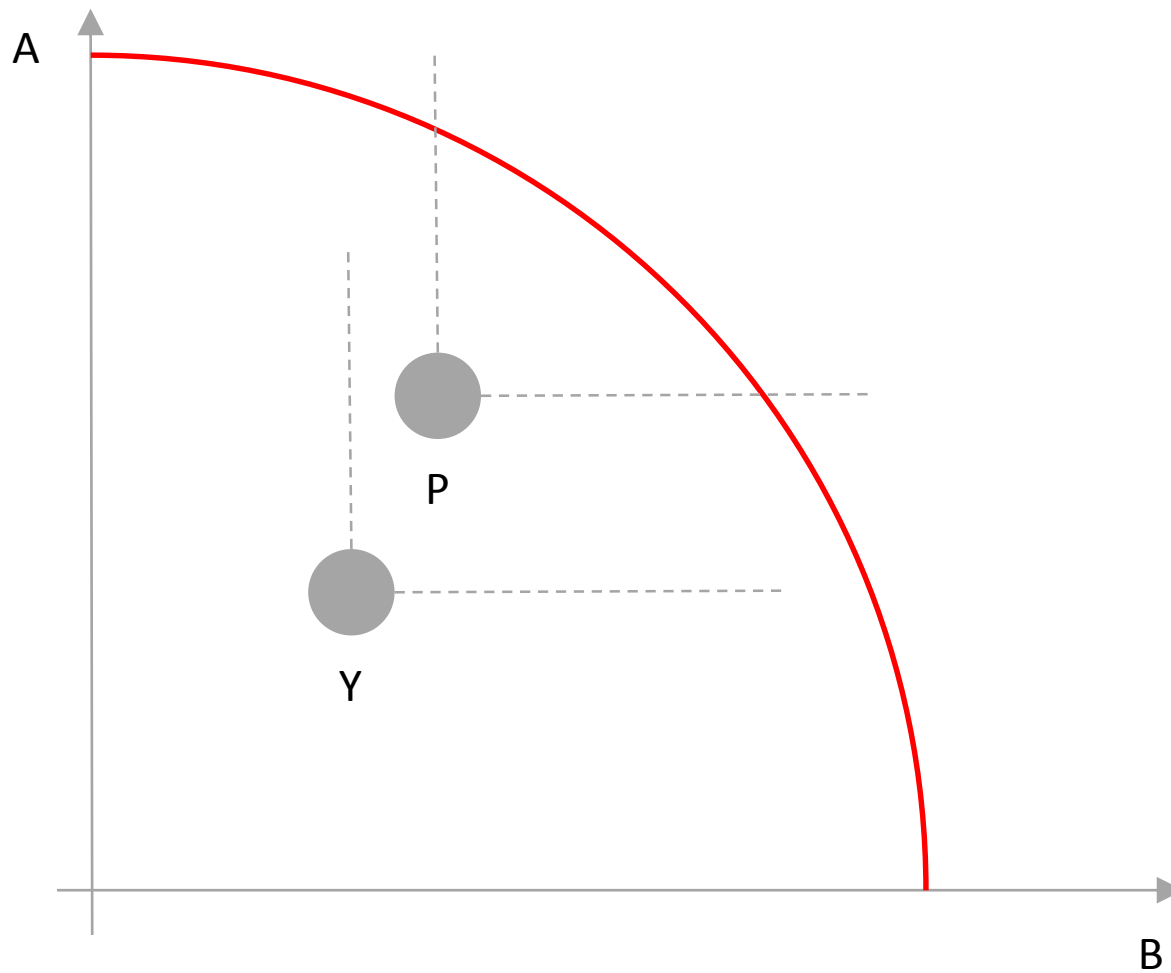


# Pareto efficiency

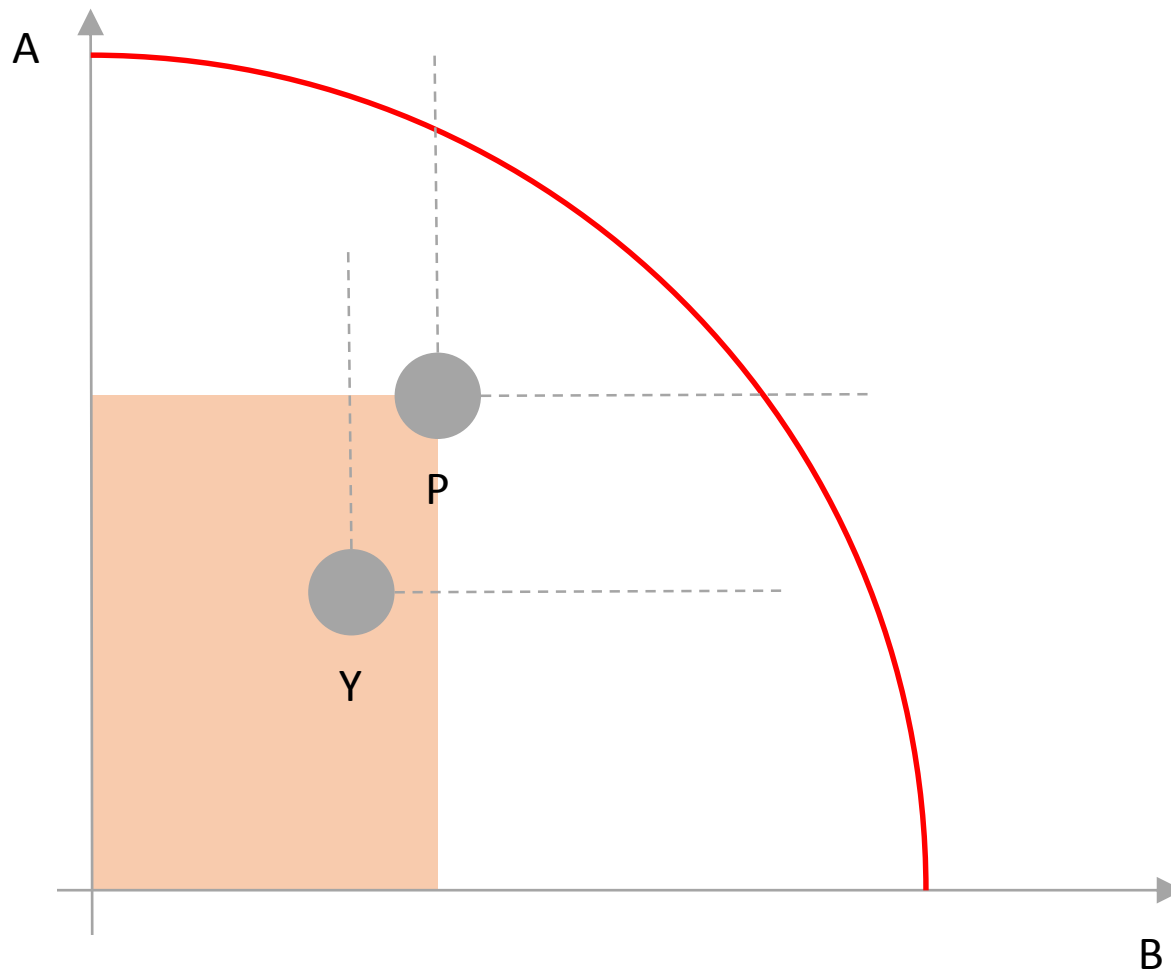




# Pareto efficiency



# Pareto efficiency



# Prisoner's dilemma – Pareto efficiency

B

		B	
		c	d
A	C	5, 5	0, 7
	D	7, 0	3, 3

# Prisoner's dilemma – Pareto efficiency

B

		B	
		c	d
A	C	5, 5	0, 7
	D	7, 0	3, 3

# Prisoner's dilemma – Pareto efficiency

B

		B	
		c	d
A	C	5, 5	0, 7
	D	7, 0	3, 3

# Game N

		B	
		l	r
A	L	2, 2	4, 0
	R	2, 3	8, -1

# Game N – Pareto efficiency

B

		B	
		l	r
A	L	2, 2	4, 0
	R	2, 3	8, -1

# Game N – Pareto efficiency

B

		B	
		l	r
A	L	2, 2	4, 0
	R	2, 3	8, -1



Pareto optimality solid tool for  
comparing equilibriums

Mixed-strategy  
Nash equilibrium

# Matching pennies

- Two players
- Players choose heads or tails
- If players match heads/tails, I (Player 1) win both coins
- If players don't match heads/tails, opponent (Player 2) wins both coins

# Matching pennies

My pair

		My pair	
		Heads	Tails
Me	Heads	1, -1	-1, 1
	Tails	-1, 1	1, -1

# Matching pennies – Pareto efficiency

My pair

		My pair	
		Heads	Tails
Me	Heads	1, -1	-1, 1
	Tails	-1, 1	1, -1

# Matching pennies – mixed strategy

My pair

		My pair	
		Heads (0.5)	Tails (0.5)
Me	Heads (0.5)	1, -1	-1, 1
	Tails (0.5)	-1, 1	1, -1

Calculation  
of mixed-strategy NE

# Game Y

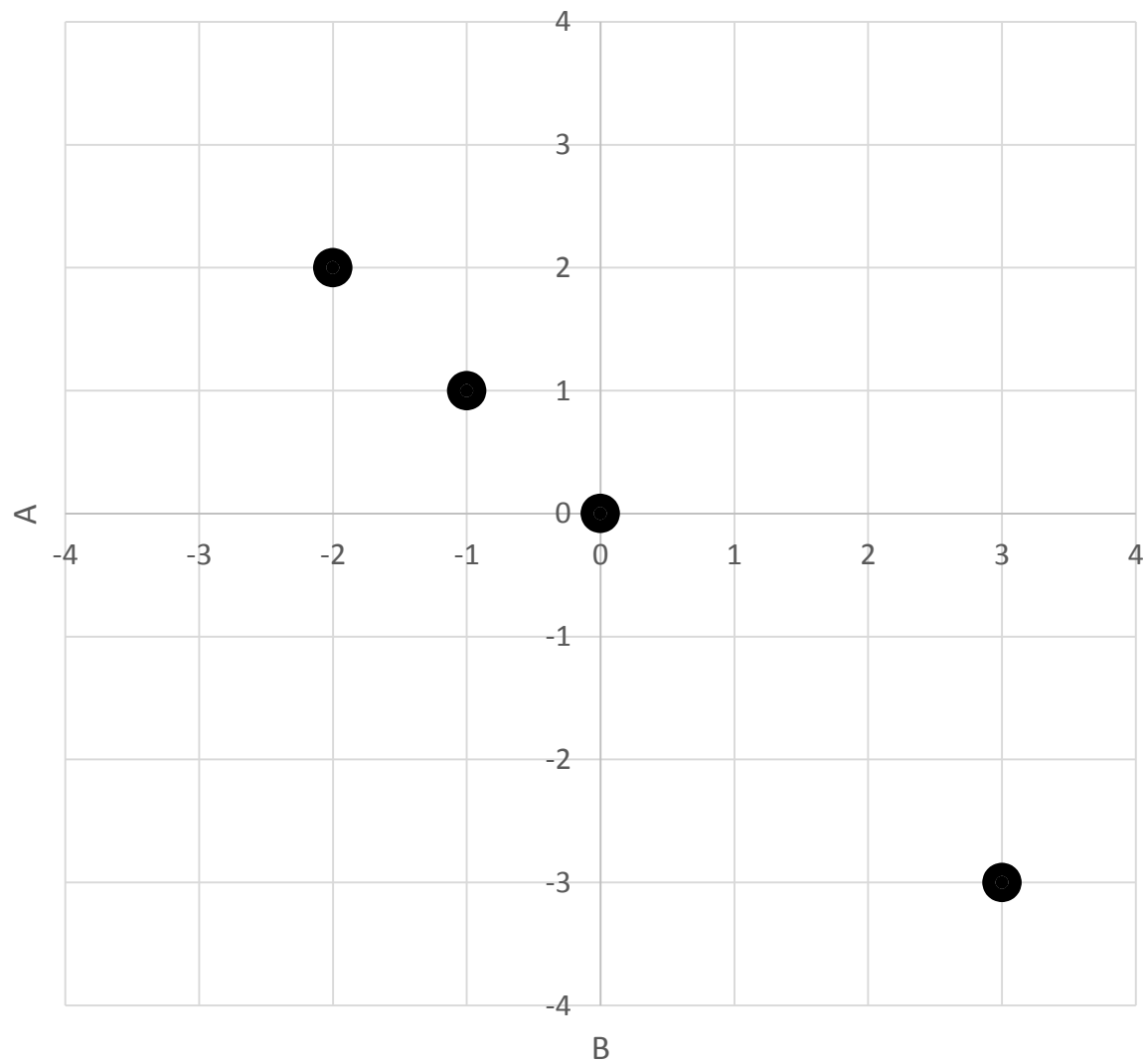
		B	
		L	R
A	U	3, -3	-2, 2
	D	-1, 1	0, 0



# Game Y – Pareto efficiency

		B	
		L	R
A	U	3, -3	-2, 2
	D	-1, 1	0, 0

# Game Y – Pareto efficiency?



# Game Y

		B	
		L (q)	R (1 - q)
A	U (p)	3 , -3	-2 , 2
	D (1 - p)	-1 , 1	0 , 0

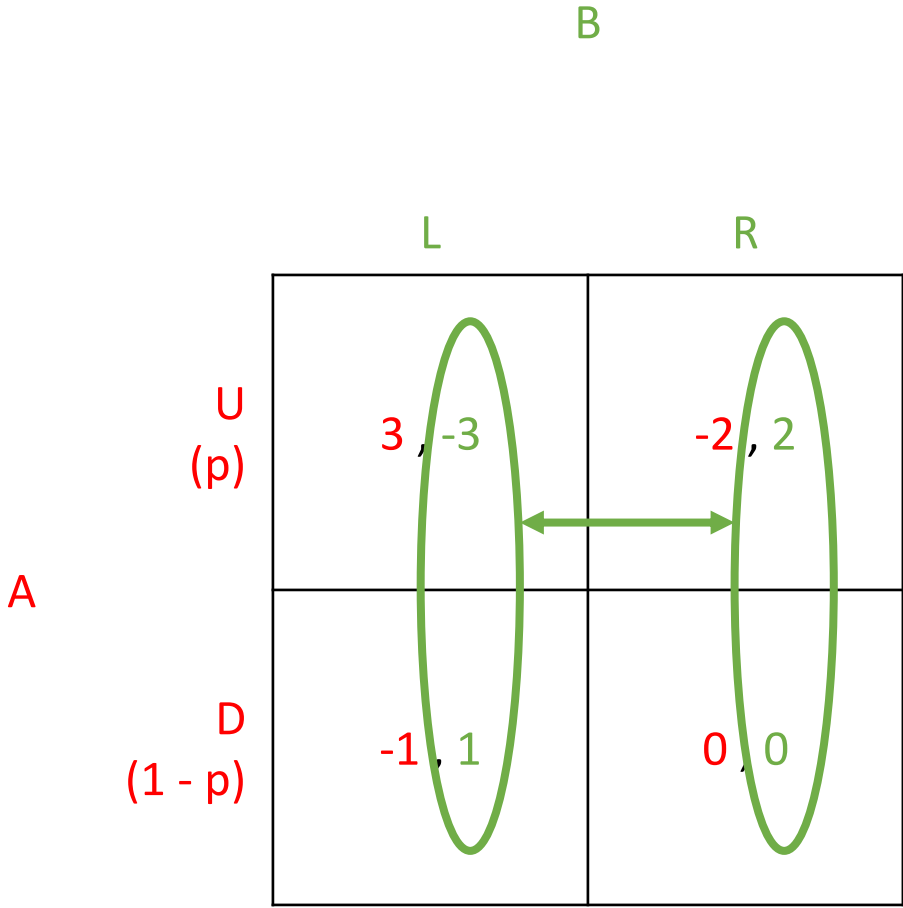
# Game Y – Player A

- Player **A** plans to mix **Up** and **Down** strategy at a certain ratio **p**
- Player **B** might play **Left** or **Right**
- Player **A** must find such a **probability** of playing **U** and **D** that makes Player **B** **indifferent** to selecting **L** or **R**
- Player **B** **has to gain same utility** from B's choice **Left** and **Right**
  - $EU_L = EU_R$
- Expected utility of Player B choosing Left:
  - $EU_L = f(p)$
- Expected utility of Player B choosing Right:
  - $EU_R = f(p)$

# Game Y

		B	
		L	R
A	U (p)	3, -3	-2, 2
	D (1 - p)	-1, 1	0, 0

# Game Y



# Game Y - Player A's strategy

- $EU_L$ :
  - Some % of time ( $p$ ) gets B utility  $-3$
  - Rest of the time ( $1 - p$ ) gets B utility  $1$

- $EU_L = (p) * (-3) + (1 - p) * (1)$

- $EU_L = -3p + 1 - p$

- $EU_L = 1 - 4p$

		B	
		L ( $q$ )	R ( $1 - q$ )
A	U ( $p$ )	3, -3	-2, 2
	D ( $1 - p$ )	-1, 1	0, 0

# Game Y - Player A's strategy

- $EU_R$ :
  - Some % of time ( $p$ ) gets B utility 2
  - Rest of the time ( $1 - p$ ) gets B utility 0

- $EU_R = (p) * (2) + (1 - p) * (0)$

- $EU_R = 2p + 0 - 0p$

- $EU_R = 2p$

		B	
		L ( $q$ )	R ( $1 - q$ )
A	U ( $p$ )	3, -3	-2, 2
	D ( $1 - p$ )	-1, 1	0, 0



# Player A's strategy – making B indifferent

## Comparison of $EU_L$ with $EU_R$

- $EU_L = 1 - 4p$

- $EU_R = 2p$

- $EU_L = EU_R$

- $1 - 4p = 2p$        $+4p$

- $1 = 6p$        $/6$

- $p = \mathbf{1/6}$

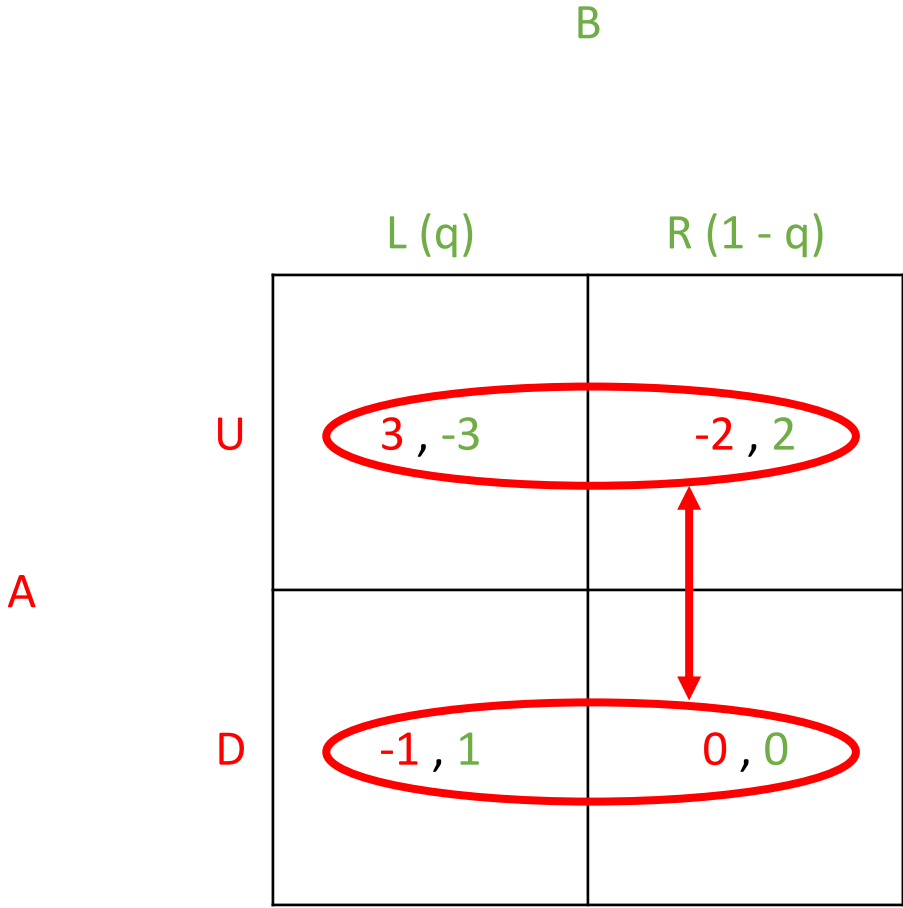
- $\mathbf{1 - p = 1 - 1/6 = 5/6}$

- We've found the ideal mixed strategy for Player A

- If Player A plays Up 1/6 of time and Down 5/6 of time, Player B is indifferent to choosing Left or Right

- We need to do the same for player B

# Game Y



# Game Y - Player B's strategy

- $EU_U$ :
  - Some % of time ( $q$ ) gets A utility 3
  - Rest of the time ( $1 - q$ ) gets A utility -2

- $EU_U = (q) * (3) + (1 - q) * (-2)$

- $EU_U = 3q - 2 + 2q$

- $EU_U = 5q - 2$

B

		B	
		L ( $q$ )	R ( $1 - q$ )
A	U ( $p$ )	3, -3	-2, 2
	D ( $1 - p$ )	-1, 1	0, 0

# Game Y - Player B's strategy

- $EU_D$ :
  - Some % of time ( $q$ ) gets A utility **-1**
  - Rest of the time ( $1 - q$ ) gets A utility **0**

- $EU_D = (q)*(-1) + (1 - q)*(0)$

- $EU_D = -1q + 0 - 0q$

- $EU_D = -q$

B

		B	
		L ( $q$ )	R ( $1 - q$ )
A	U ( $p$ )	3, -3	-2, 2
	D ( $1 - p$ )	-1, 1	0, 0

# Player B's strategy – making A indifferent

## Comparison of $EU_U$ with $EU_D$

- $EU_U = 5q - 2$

- $EU_D = -q$

- $EU_U = EU_D$

- $5q - 2 = -q$        $-5q$

- $-2 = -6q$        $/-6$

- $q = 1/3$

- $1 - q = 1 - 1/3 = 2/3$

- We've found the ideal mixed strategy for Player B

- If Player B plays Left 1/3 of time and Down 2/3 of time, Player A is indifferent to choosing Up or Down

Mixed strategy NE

(  $\frac{1}{6}$  U ,  $\frac{1}{3}$  L )

# Game Y - MSNE

B

		B	
		L (1/3)	R (2/3)
A	U (1/6)	3, -3	-2, 2
	D (5/6)	-1, 1	0, 0

# Battle of sexes

- Want to go out together but have no means of communication
  - Have 2 choices – ballet or car show
  - Player A prefers car show (C)
  - Player B prefers ballet (B)
  - Both prefer being together than being alone (A)
- 
- Preferences for player A:  $C > B > A$
  - Preferences for player B:  $B > C > A$



# Battle of sexes

		B	
		b	c
A	B	1, 2	0, 0
	C	0, 0	2, 1

# Battle of sexes – PS Nash equilibriums

B

		b	c
A	B	1, 2	0, 0
	C	0, 0	2, 1

# Equilibriums

- 2 pure-strategies equilibriums
- How would they coordinate?
- Apart from pure strategies equilibriums there is one mixed strategy equilibrium for this game
- (  $\frac{1}{3} B$ ,  $\frac{2}{3} b$  )

# Battle of sexes – mixed strategy equilibrium

B

		B	
		b $\frac{2}{3}$	c $\frac{1}{3}$
A	B $\frac{1}{3}$	1, 2	0, 0
	C $\frac{2}{3}$	0, 0	2, 1

Calculation of MS NE payoffs

# Battle of sexes – mixed-strategy NE payoffs

B

		B	
		b $\frac{2}{3}$	c $\frac{1}{3}$
A	B $\frac{1}{3}$	$1, 2$ $\frac{1}{3} * \frac{2}{3}$	$0, 0$ $\frac{1}{3} * \frac{1}{3}$
	C $\frac{2}{3}$	$0, 0$ $\frac{2}{3} * \frac{2}{3}$	$2, 1$ $\frac{2}{3} * \frac{1}{3}$

# Battle of sexes – mixed-strategy NE payoffs

B

		B	
		b $\frac{2}{3}$	c $\frac{1}{3}$
A	B $\frac{1}{3}$	1, 2 $\frac{2}{9}$	0, 0 $\frac{1}{9}$
	C $\frac{2}{3}$	0, 0 $\frac{4}{9}$	2, 1 $\frac{2}{9}$

# BoS – Payoffs for player A

- We simply multiply payoffs for player A and probabilities for each outcome and then sum them together

- Player A's payoffs:

- $u(B, b) = 1 * 2/9 = 2/9$

- $u(B, c) = 0 * 1/9 = 0$

- $u(C, b) = 0 * 4/9 = 0$

- $u(C, c) = 2 * 2/9 = 4/9$

- $EU(A) = 2/9 + 0 + 0 + 4/9$

- $EU(A) = 6/9$

- $EU(A) = 2/3$

		B	
		b 2/3	c 1/3
A	B 1/3	1, 2 2/9	0, 0 1/9
	C 2/3	0, 0 4/9	2, 1 2/9



# BoS – Payoffs for player B

- We simply multiply payoffs of player B and probabilities for each outcome and then sum them together

- Player A's payoffs:

- $u(B, b) = 2 * 2/9 = 4/9$
- $u(B, c) = 0 * 1/9 = 0$
- $u(C, b) = 0 * 4/9 = 0$
- $u(C, c) = 1 * 2/9 = 2/9$

- $EU(B) = 4/9 + 0 + 0 + 2/9$
- $EU(B) = 6/9$
- $EU(B) = 2/3$

B

		b 2/3	c 1/3
B 1/3	$1, 2$ 2/9	$0, 0$ 1/9	
C 2/3	$0, 0$ 4/9	$2, 1$ 2/9	

A

# Battle of sexes NE

- Pure strategies NE

- ( **B** , **b** )
  - $EU(A) = 1$
  - $EU(B) = 2$
- ( **C** , **c** )
  - $EU(A) = 2$
  - $EU(B) = 1$

- Mixed strategies NE

- (  $\frac{1}{3}$  **B** ,  $\frac{2}{3}$  **b** )
  - $EU(A) = \frac{2}{3}$
  - $EU(B) = \frac{2}{3}$

B

		b	c
B	1, 2	0, 0	
C	0, 0	2, 1	

A

# FSS entrance game

- Two students meet at the main faculty entrance
- Both simultaneously decide whether to walk or stop
- If both walk, they **collide** and both **get a bruise** (payoff -5)
- If one stops and other walks
  - Student who stopped gets **good karma** for letting the other pass with **payoff 1**, but at the same time gets **delayed**, which is **completely offsetting** the value of the good karma
  - Student who walked gets to **pass quickly** and thus gets **payoff 1**
- If **both stop**, each would get **good karma** for letting the other pass, but both will get **delayed**

# FSS entrance game

		B	
		W	S
A	W	-5, -5	1, 0
	S	0, 1	0, 0

# FSS entrance game NE

B

		B	
		W 1/6	S 5/6
A	W 1/6	-5, -5	1, 0
	S 5/6	0, 1	0, 0

# Stag hunt

		B	
		s	r
A	S	5, 5	0, 3
	R	3, 0	3, 3

# Stag hunt NE

- Pure strategies NE

- $(S, s)$

- $EU(A) = 5$

- $EU(B) = 5$

- $(R, r)$

- $EU(A) = 3$

- $EU(B) = 3$

- Mixed strategies NE

- $(\frac{3}{5} S, \frac{3}{5} s)$

- $EU(A) = 3$

- $EU(B) = 3$

B

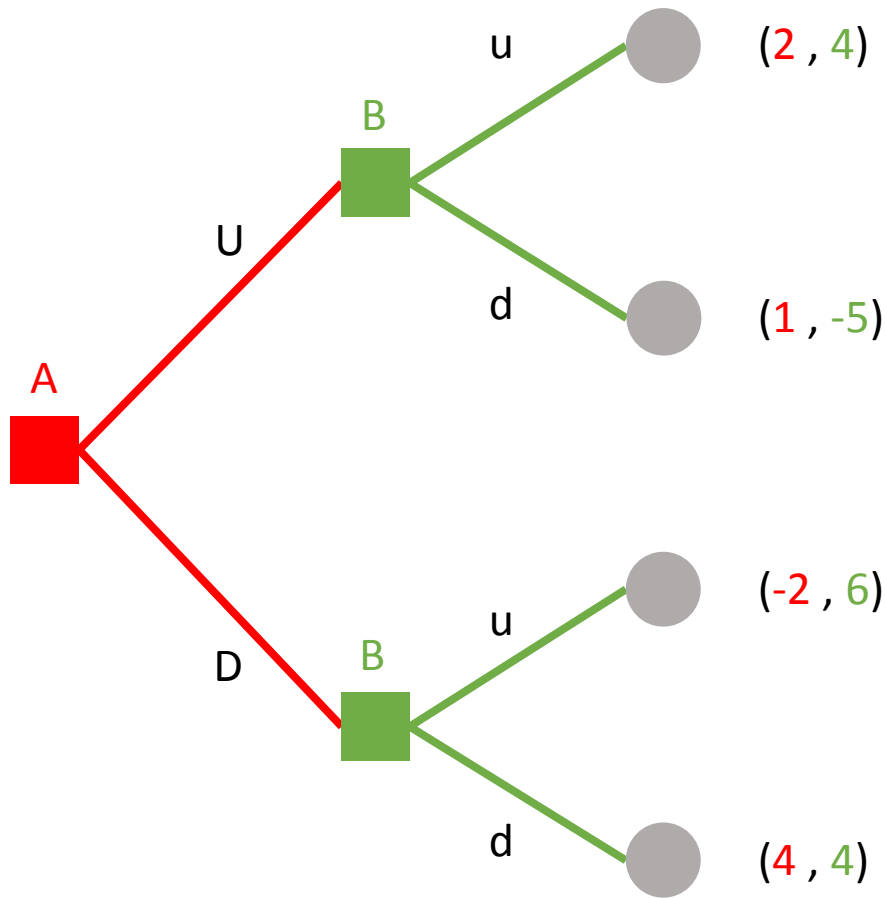
		s	R
A	S	5, 5	0, 3
	R	3, 0	3, 3

Extensive form games



# Extensive form games

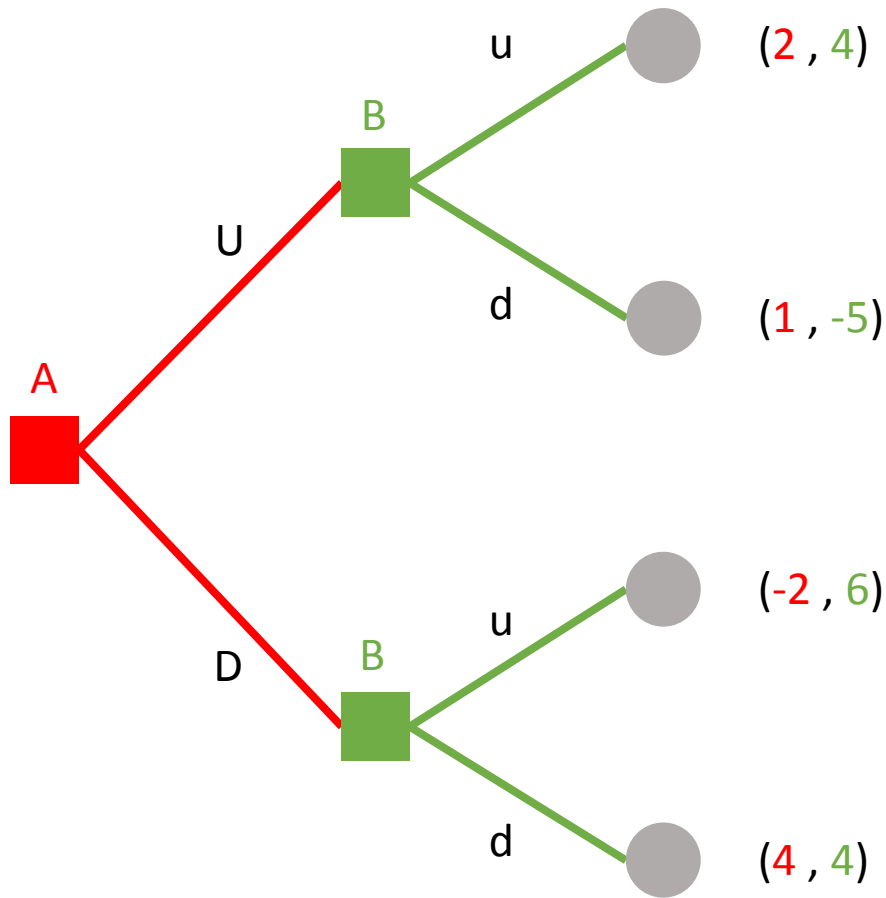
- Visualized as a **game (decision) tree**
- Players move **sequentially**
- Captures **time** in game
- Captures knowledge of agents – sometimes agents do not have information where they are located in game



# Basic terminology

- Each square is called **node**
- Each line represents an **action** an **owner of the node** has at his disposal
- Nodes might either **trigger other action** or **end**
- Every circle is an end of the tree – it's called **terminal node**
- Every circle must yield **payoffs** for all the actors
- Each moment actor has **information about all previous moves** called **information set**

Backwards induction

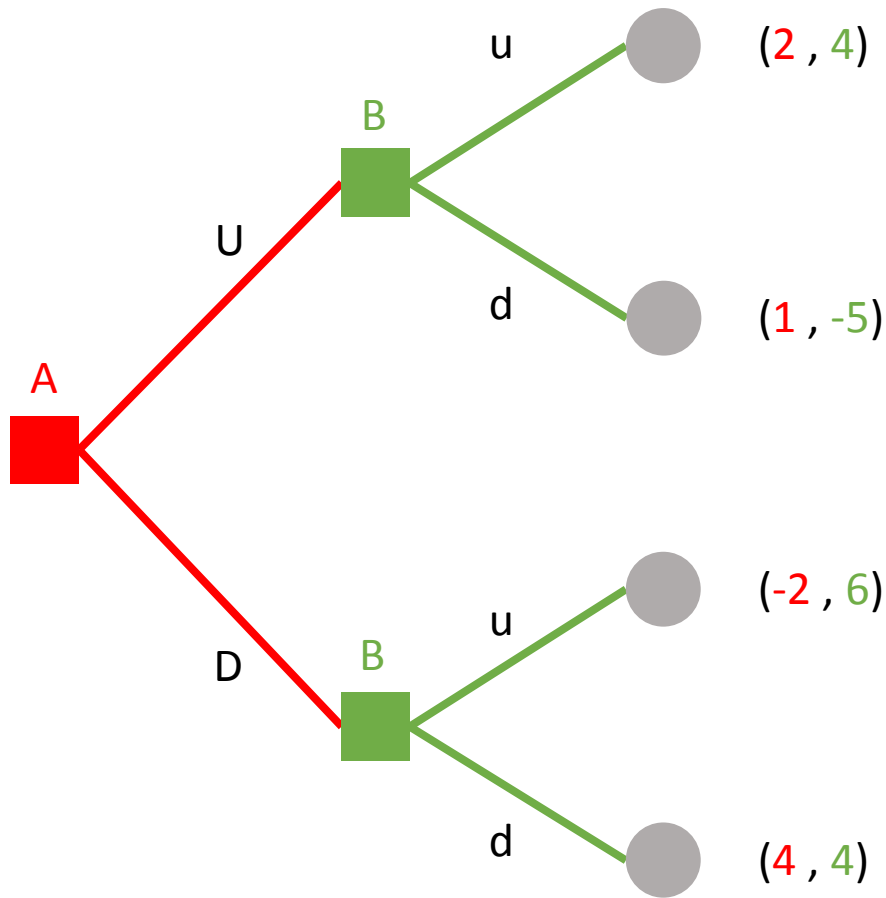


# Backwards induction

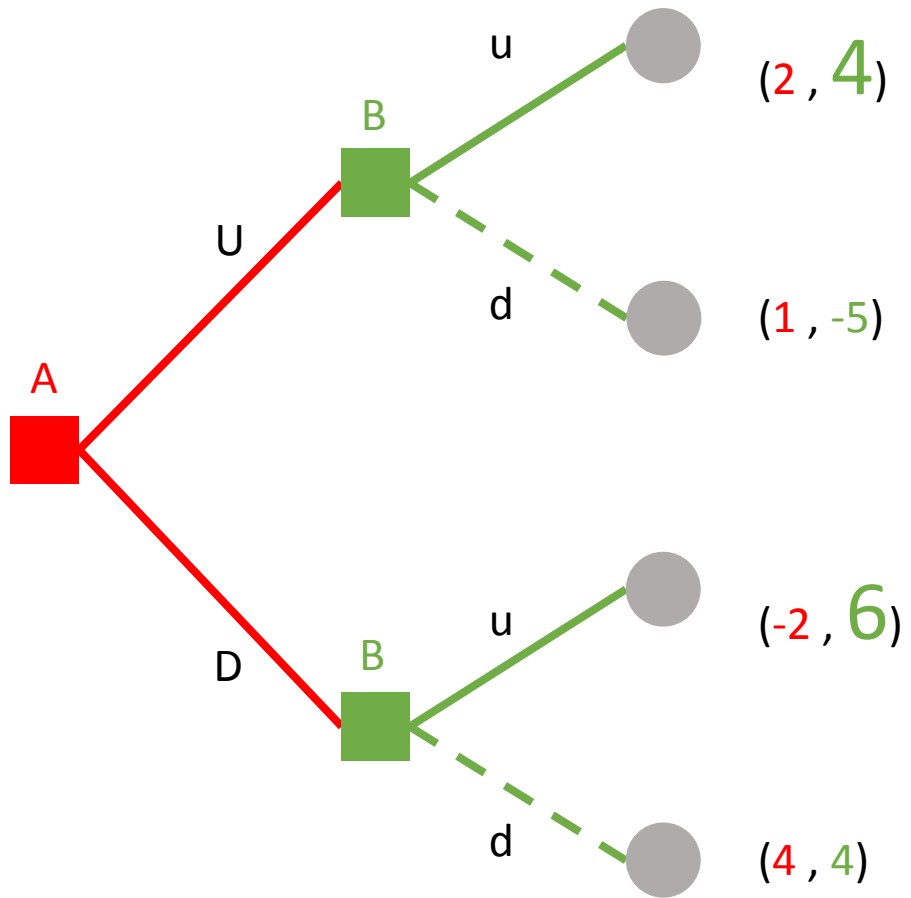
- Moves player make at nodes reached in an equilibrium are called **behavior on the equilibrium path**
- Moves player make at nodes that are not reached in an equilibrium are behavior **off the equilibrium path**

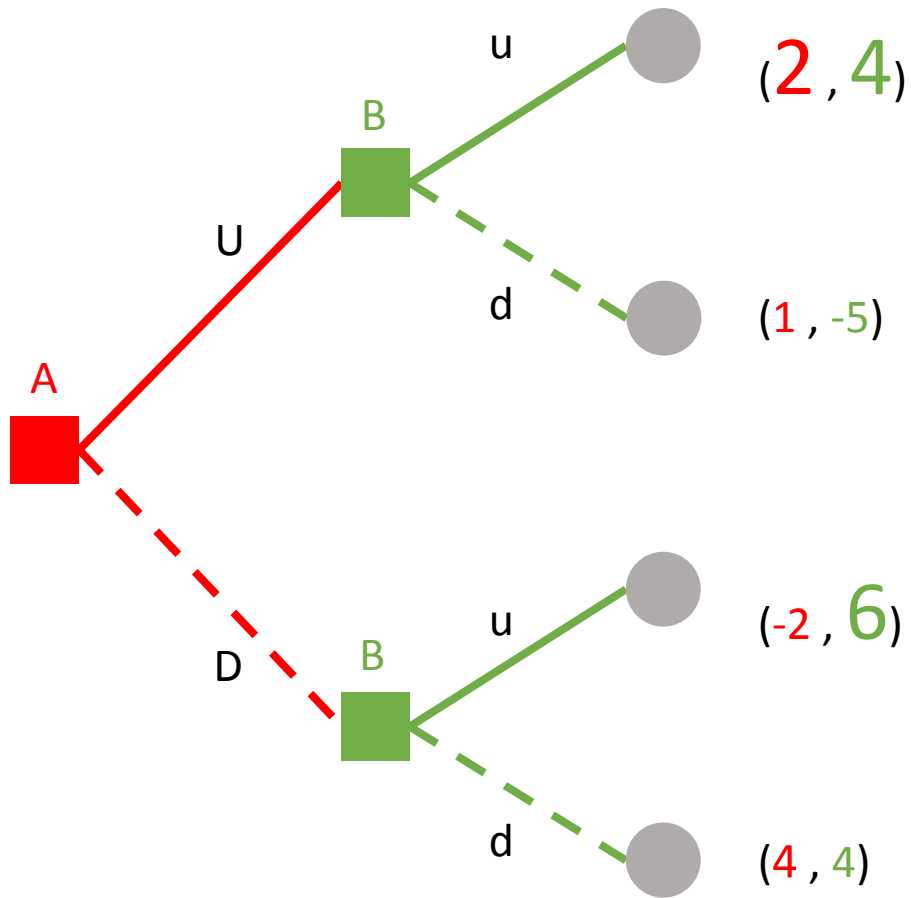
# Backwards induction

- Begin with decisions that lead only to terminal nodes
- Compare payoffs for decisions in each node leading to terminal node
- Find **best reply to alternatives of player** playing at the **current node**
- Work through nodes **backwards** and solve the outcomes of all nodes comparing payoffs for respective players









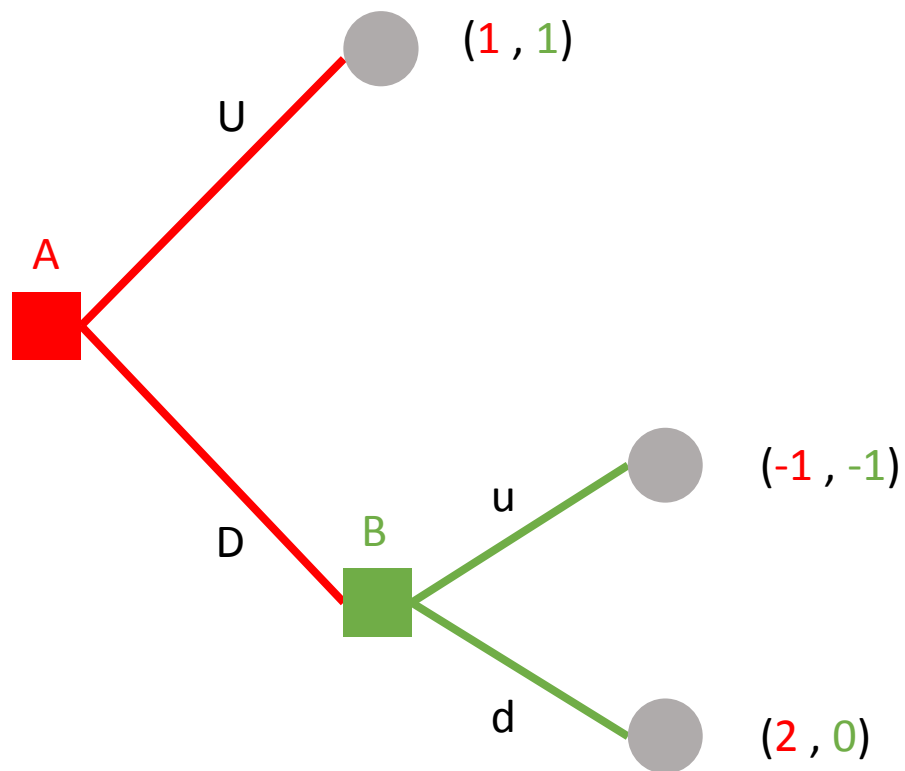
# Equilibrium of sequential game

- One Nash equilibrium in pure strategies on the equilibrium path
- ( U ; u , u )
- There may be **more NE in pure strategies**
- B decides – if A goes U than u yields better payoff in the upper node, if A goes D than u yields better payoff in the lower node
- A knows that B will choose u in both nodes, therefore compares payoffs in u for going U or D – U yields better payoff

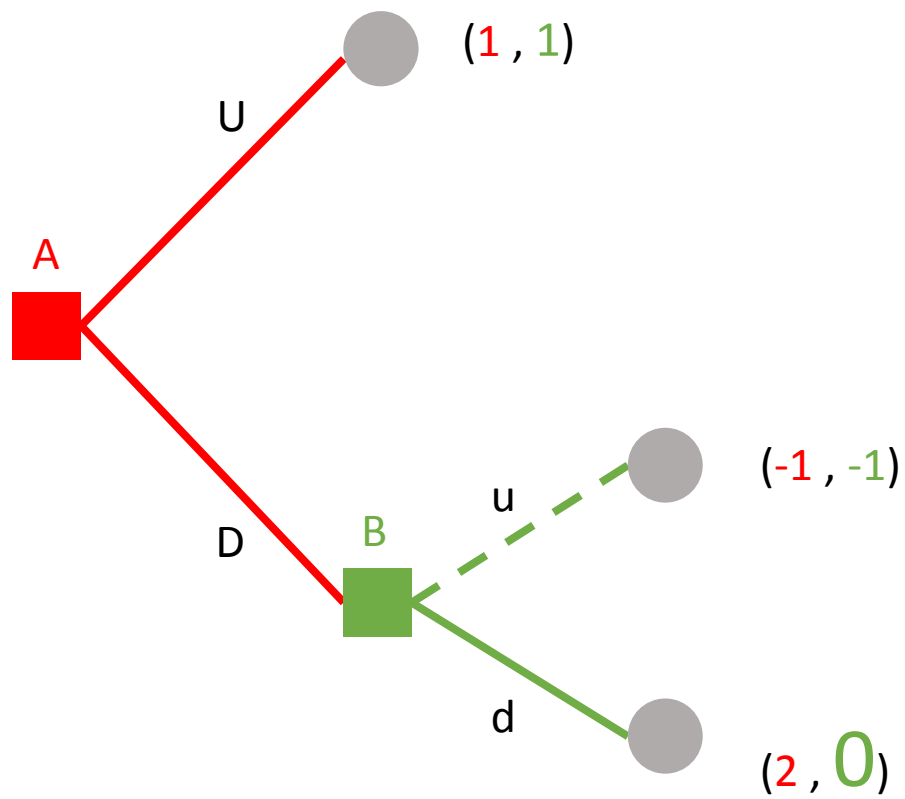
# Backwards induction

- Works in games of **perfect information**
- All actors are aware of all previous actions and can also **anticipate**, what actors will do based on their expected utilities over outcomes at subsequent nodes – actors have a **perfect recall**
- However, backwards induction assesses only rationality on the equilibrium path. **NE found off the equilibrium path will not be found** through backwards induction

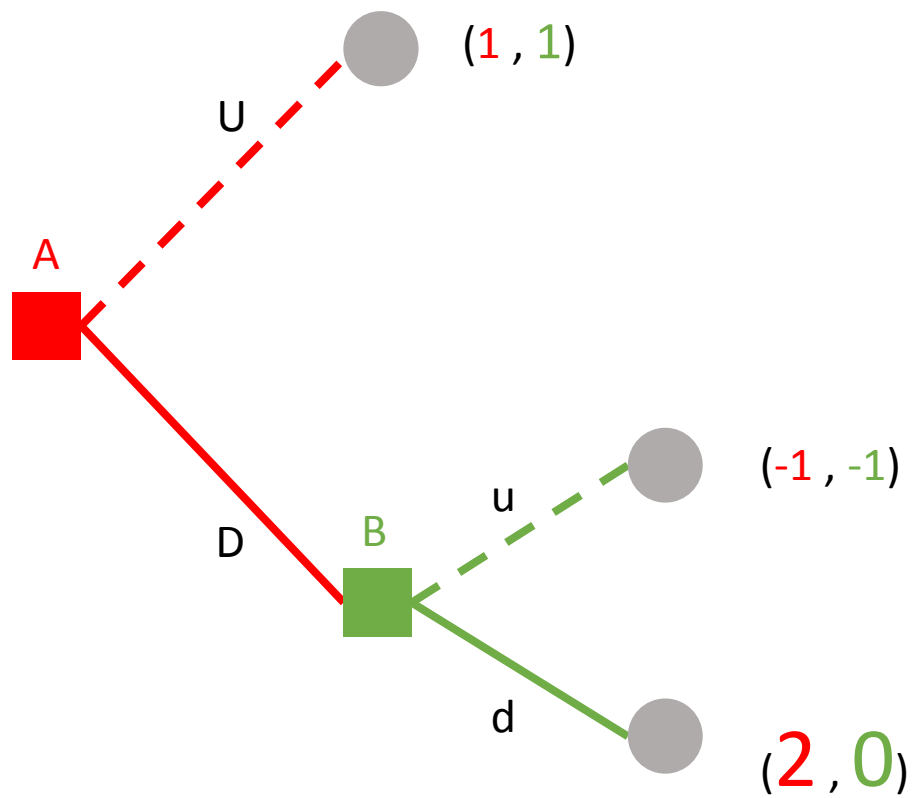
# Selten's game



# Selten's game



# Selten's game

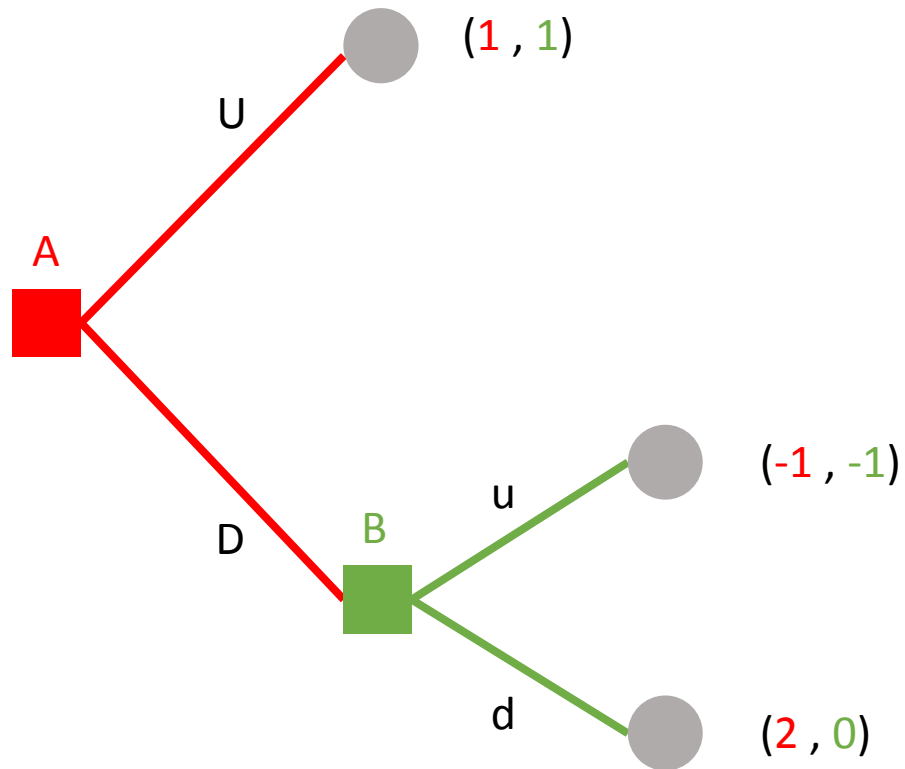


# Rewrite Selten's game into matrix

- Player A has 2 moves
  - U, D
  - Represented as 2 rows in a matrix
- Player B has also 2 moves
  - u, d
  - Represented as 2 columns in a matrix
- We have 2x2 game in normal form



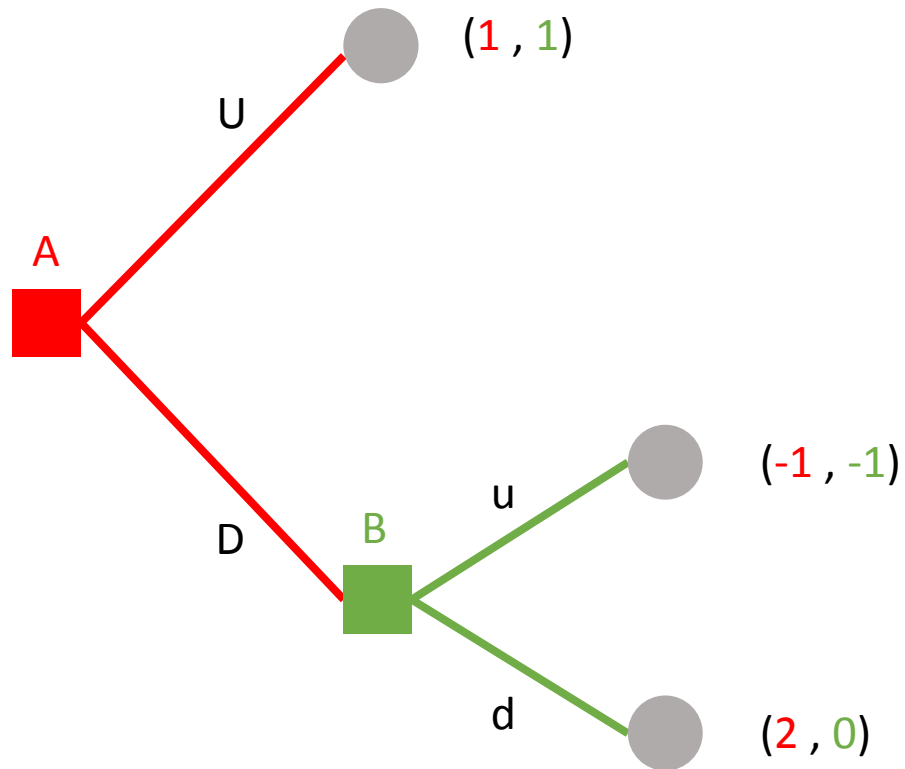
# Selten's game



# Strategic form of Selten's game

		B	
		u	d
A	U		
	D		

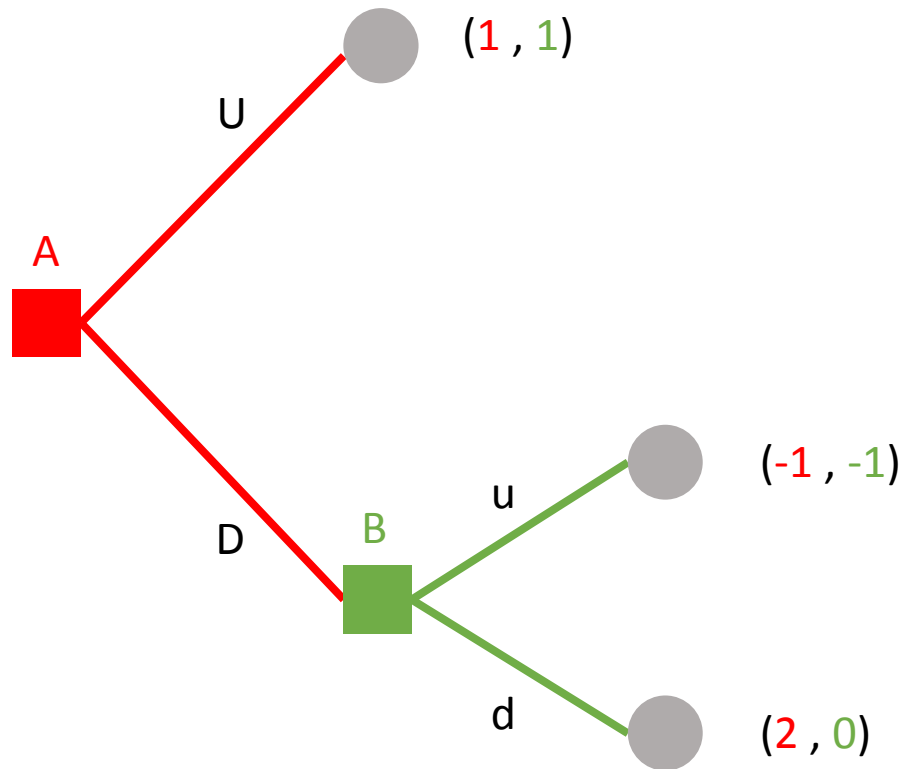
# Selten's game



# Strategic form of Selten's game

		B	
		u	d
A	U		
	D	-1, -1	2, 0

# Selten's game



# Strategic form of Selten's game

		B	
		u	d
A	U	1, 1	
	D	-1, -1	2, 0

# Strategic form of Selten's game

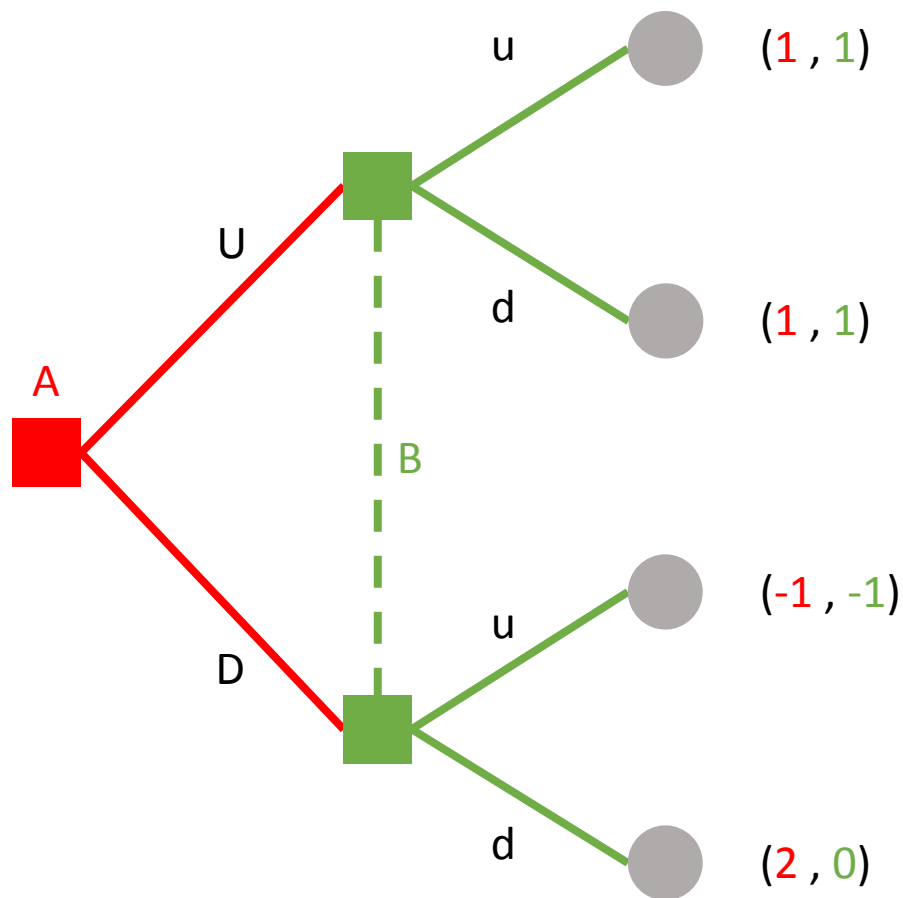
		B	
		u	d
A	U	1, 1	1, 1
	D	-1, -1	2, 0

# Rewriting extensive form into normal form

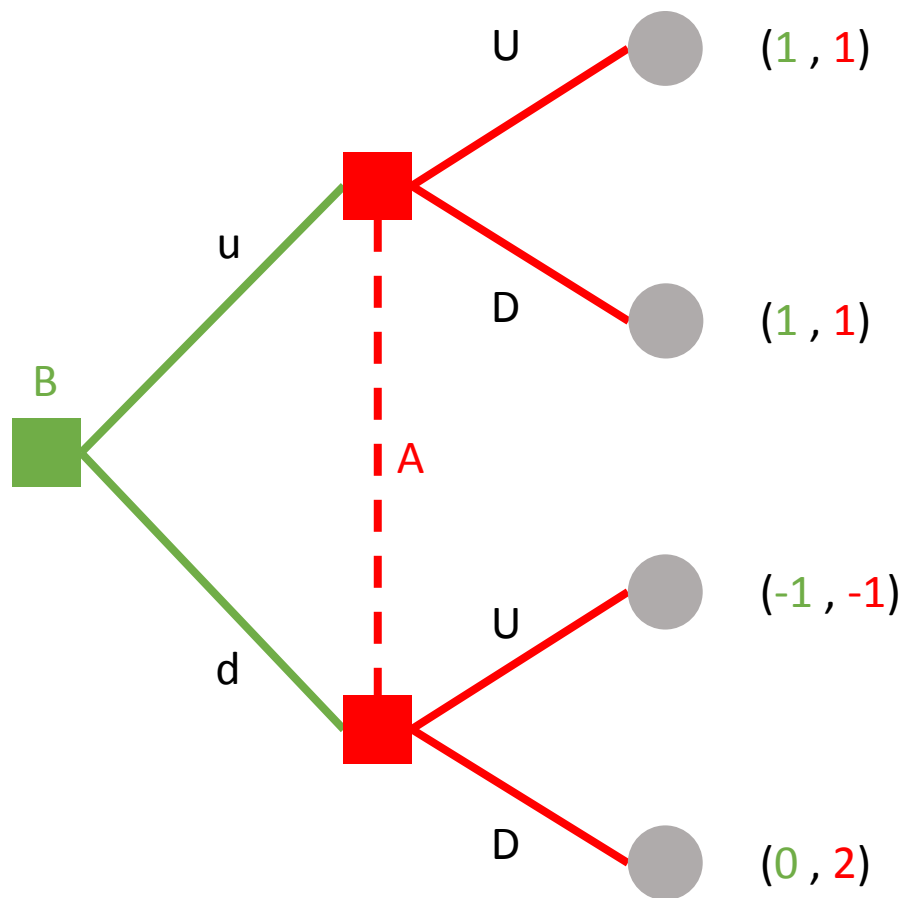
- If we rewrite extensive form game into normal form, only **one matrix will emerge** as a representation
- This **does not work the other way around**
- Since matrixes do not hold information about sequence of actions, one normal-form game might have multiple extensive-form representations that would completely alter outcomes of the game



# Selten's game from matrix



# Selten's game from matrix



# Selten's game – Nash equilibrium

		B	
		u	d
A	U	1, 1	1, 1
	D	-1, -1	2, 0

# Selten's game – Nash equilibrium

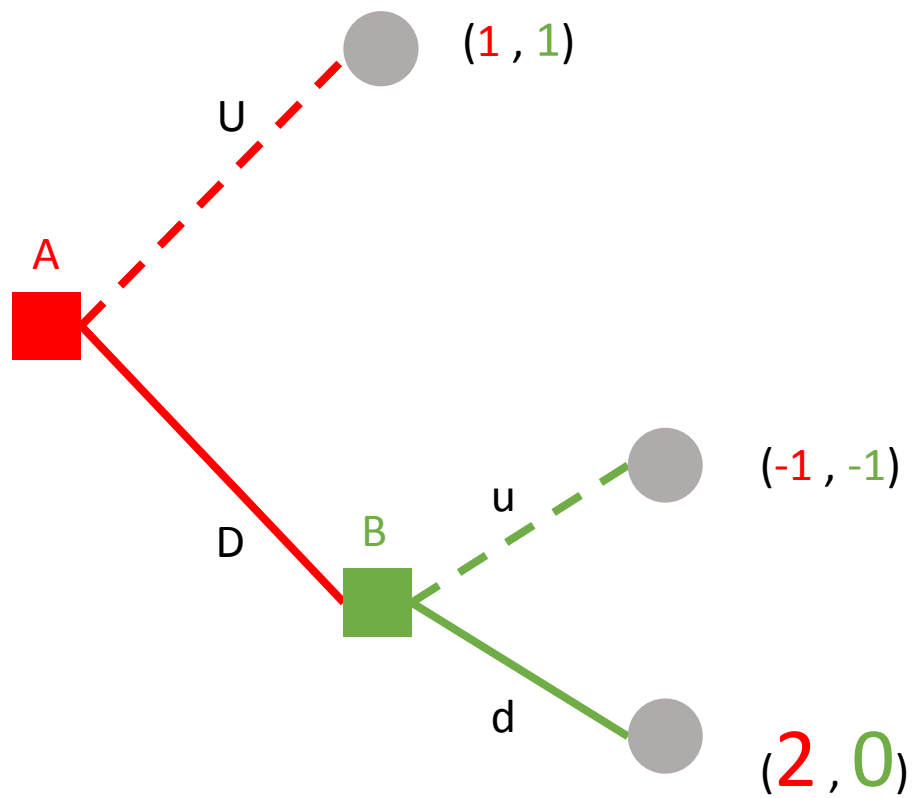
		B	
		u	d
A	U	1, 1	1, 1
	D	-1, -1	2, 0

# Selten's game – Nash equilibrium

B

		B	
		u	d
A	U	1, 1	1, 1
	D	-1, -1	2, 0

# Selten's game



# NE off the equilibrium path

- This game has another NE which is represented by action U which is not revealed in extensive form
- This equilibrium (U, u) is called a **non-credible threat**
  - B is willing to get its largest payoff, which will result from A playing U (deciding not to play the game)
  - B can change its payoff from A's decision about U and D only by threatening that it will play u if A plays D
  - But B will never play u, since it brings lower payoff than playing d

# Selten's game and non-credible threat

