

- These notes will combine what is in the lectures, possibly slightly re-organised, and with a few extra words.
- Anything new will be in **green!**

Algebra 4 - 2021

- 1) Homological algebra
- 2) Commutative algebra (rel. to alg. geo)
- 3) Representation theory of finite groups.

Course details

- Weekly lecture uploaded to IS
- Exercise class on Tuesday
- 3 marked assignments (50%) + oral exam

Homological algebra - overview of topics covered

- Ex. sequences, ch. complexes, long ex. seq.
 - Homology of ch. comp.
 - Simpl. & sing. homology
 - Cat of ch. comp & homology a Functor
 - Chain homotopies - induce equal maps on homology
 - Cochain complexes & cohomology (singular cohomology)
- Pre-additive & additive cats
 - biproducts
 - kernels, cokernels & abelian categories
 - Additive & exact Functors
 - Freyd-Mitchell embedding \rightarrow to prove something in a gen abelian cat, suff. to prove it in a $\text{Mod } R$
- Snake lemma & long ex. seq. of homology
 - Homology groups of spheres (in ex. class)
- Homology & cohomology of alg. structures
 - proj & inj. obs
 - Proj., inj resolution
 - left & right derived functors $L_n F, R_n F$
 - Functoriality & their props: $L_0 F, L_1 F$ & les.
 - Def of Ext & Tor via der. functors
 - Ext via extensions
 - Mentioned group cohomology (very briefly)
- Proj & inj. dimension
 - Proj, inj dim. of modules
 - Rel. to Ext
 - Calc for abelian groups
 - P.d. of a ring, global dim.
 - Hilb. S & theorem (no proof)

Homological algebra

Exact sequences

- R a ring
- Mod_R cat of left R -modules.
- Consider $A \xrightarrow{f} B \xrightarrow{g} C \in \text{Mod}_R$
in which $g \circ f = 0$, the zero homomorphism.
- Then $g(fa) = 0$ all $a \in A$ so
that $\text{im}(f) \subseteq \text{Ker}(g)$.

Def) - The sequence is said to be exact at B if $\text{im}(f) = \text{Ker}(g)$.

Example

- $0 = \{0\} \rightarrow A \xrightarrow{f} B$ is exact $\Leftrightarrow \text{Ker} f = 0$
 $\Leftrightarrow f$ is injective (mono)
- $A \xrightarrow{f} B \rightarrow 0$ is exact $\Leftrightarrow \text{im} f = B$
 $\Leftrightarrow f$ is surjective (epi - we will see)
later
- A short exact sequence is a sequence
 $0 \rightarrow A \xrightarrow{f} B \xrightarrow{g} C \rightarrow 0$
which is exact in each position :
 - ex@A : f is injective

- ex @ C : g is surjective

- ex @ B : $\text{im } f = \ker g$, but since g is surjective $C \cong B / \ker g = B / \text{im } f$.

Exercise : Using this, show that each ses is of the form

$$\underline{0 \longleftarrow A \longleftarrow B \longrightarrow B/A \longrightarrow 0}$$

for A a submodule of B ,
(up to isomorphism of sequences)

Defⁿ . A chain complex A is a sequence

$$\dots \xrightarrow{d_{n+1}} A_n \xrightarrow{d_n} A_{n-1} \longrightarrow \dots$$

of R -modules (for $n \in \mathbb{Z}$)

where $d_n \circ d_{n+1} = 0 \quad \forall n$.

- The elts of $Z_n := \ker d_n \subseteq A_n$ are called n -cycles
- $B_n := \text{im } d_{n+1} \subseteq \ker d_n$ are called n -boundaries

Defⁿ) The n -th homology of A is the

R -module $H_n(A) = \frac{\ker d_n}{\operatorname{im} d_{n+1}} = \frac{Z_n}{B_n}$.

Remark) - Observe that

$$A \text{ is exact @ } A_n \iff H_n(A) = 0.$$

Thus the n -th homology measures the failure of A to be exact at A_n .

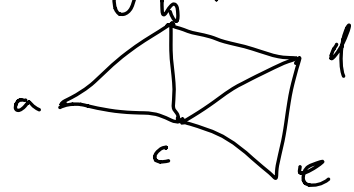
- If A is exact at all n , it is called a long exact sequence.

Examples

① Simpl. homology
 K a geom. simp. complex - eg. triang. space

- K_n set of n -simplices.

Each has $n+1$ faces.



- If set of vertices is ordered, so are these maps (omit i 'th vertex)

so $K_n \xrightarrow{d_i} K_{n-1}$ for $i=0, \dots, n$.

- Forming free R -modules, gives

$C_n \xrightarrow{d_i} C_{n-1}$ applying to basis elts
 & then $\dots C_n \xrightarrow{\sum_{i=0}^n (-1)^i d_i} C_{n-1} \dots$ is a chain
 complex C_K ,

with $H_n(C_K) =$ simplicial homology of K .

② Singular homology:

X a top. space,

$S_n(X) =$ Free R -module

on set of cts maps $\Delta_n \rightarrow X$
 where Δ_n is the stand. n -simplex.

- As before get Face maps, &

ch. comp $S_K(X) \xrightarrow{d = \sum (-1)^i d_i} S_{K-1}(X)$

where homology is
singular homology of X . (same as S)
 (simp h. ...)

- A chain map $f: A \rightarrow B$ of chain complexes consists of maps $f_n: A_n \rightarrow B_n$ such that

$$\begin{array}{ccc} A_{n+1} & \xrightarrow{d_{n+1}} & A_n \\ f_{n+1} \downarrow & \approx & \downarrow f_n \\ B_{n+1} & \xrightarrow{d_{n+1}} & B_n \end{array} \quad \forall n.$$

Notation: One often just writes $d: B_{n+1} \rightarrow B_n$ when the context is clear.

- Chain complexes and chain maps form a category $\text{Ch}(\text{Mod}_R)$.

Proposition

The n -th homology determines a functor $H_n: \text{Ch}(\text{Mod}_R) \rightarrow \text{Mod}_R$.

Proof) It sends $A \mapsto H_n(A)$.

At $f: A \rightarrow B$, then given

$x \in Z_n(A) = \ker(d: A_n \rightarrow A_{n-1})$ we have

$d f x = f d x = 0$ so $f x \in Z_n(B)$;

similarly if $x \in B_n(A)$ then $f(x) \in B_n(B)$;

hence we obtain

$$H_n(A) = \frac{\mathbb{Z}_n A}{B_n A} \longrightarrow \frac{\mathbb{Z}_n B}{B_n B} = H_n(B)$$
$$[x] \longmapsto [fx]$$

This is clearly functorial. \square

Homotopy & homology

Defⁿ) let $f, g: A \rightrightarrows B$ be chain maps. A chain htpy s from f to g (written $s: f \rightsquigarrow g$) is a sequence of homomorphisms $s_n: A_n \rightarrow B_{n+1}$

as in the picture below :

$$\begin{array}{ccccccc}
 \dots & A_{n+1} & \xrightarrow{d_{n+1}} & A_n & \xrightarrow{d_n} & A_{n-1} & \dots \\
 & \downarrow f_{n+1} - g_{n+1} & \searrow s_n & \downarrow f_n - g_n & \searrow s_{n-1} & \downarrow f_{n-1} - g_{n-1} & \\
 \dots & B_{n+1} & \xrightarrow{d_{n+1}} & B_n & \xrightarrow{d_n} & B_{n-1} & \dots
 \end{array}$$

such that

$$d_{n+1} s_n + s_{n-1} d_n = g_n - f_n \quad \forall n \in \mathbb{Z}.$$

- A chain map $f: A \rightarrow B$ is null homotopic if $f \rightsquigarrow 0$.

- It is a htpy equivalence if $\exists g: B \rightarrow A$ such that $fg \sim 1_B$ & $gf \sim 1_A$.

Lemma

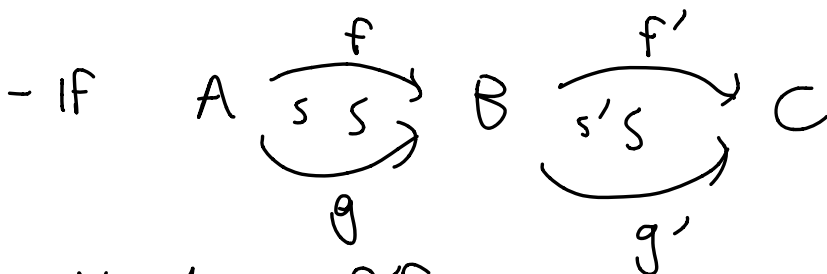
Homotopy is an e-rel, compatible with composition.

Proof - Taking $s_n = 0$ shows $f \sim f$.

- If $s: f \sim g$ & $t: g \sim h$ take $stt: f \sim h$; indeed

$$h - f = (h - g) + (g - f) = (dstsd) + (dt + td) = d(stt) + (stt)d.$$

- If $s: f \sim g$ then $-s: g \sim f$ where $(-s)_n = -s_n$.



must show $f' \sim g'g$.

- Suff. to show $f' \sim f'g \sim g'g$ by above.

- Consider $f's$ with $(f's)_n = f'_{n+1}s_n$.

$$\begin{aligned}
 - \text{Then } d(fs) + (fs)d &= f'ds + f'sd \\
 &= f'(ds + sd) \\
 &= f'(g - f) \\
 &= f'g - f'f
 \end{aligned}$$

$$\begin{aligned}
 \text{Sim. } d(s'g) + (s'g)d &= ds'g + s'dg \\
 &= (ds' + s'd)g \\
 &= (f' - g')g = f'g - g'g \quad \square
 \end{aligned}$$

lemma

If $f \sim g$, they induce the same map
 $H_n f = H_n g$ on homology.

Proof - Consider $x \in Z_n(A) = \ker(d: A_n \rightarrow A_{n-1})$
 - we must show $f(x), g(x) \in Z_n(B)$
 coincide

modulo $B_n(B) = \text{Im}(d_{n+1}: B_{n+1} \rightarrow B_n)$.

$$\begin{aligned}
 - \text{But } g(x) - f(x) &= sdx + dsx = dsx \\
 &\in \text{Im}(d_{n+1}) = B_n(B)
 \end{aligned}$$

$$\text{as } sdx = s0 = 0.$$

- Hence $H_n(f) = H_n(g)$.

Corollary

If $f: A \longrightarrow B$ is htpy equiv,
then $H_n f: H_n A \longrightarrow H_n B$ is invertible.

Proof

If $gf \sim 1_A$ & $fh \sim 1_B$ then

$$H_n(g)H_n(f) = \text{by funet.}$$

$$H_n(gf) = \text{by previous prop}$$

$$H_n(1_A) = \text{by funet.}$$

$$1_{H_n(A)}.$$

$$\text{Sim } H_n(f)H_n(g) = 1_{H_n(B)}.$$

Cohomology

- So far we have talked about chain complexes & homology
- Dually we have cochain complexes & cohomology
- A cochain complex in \mathcal{C} is a diagram
$$\dots X^{n-1} \xrightarrow{d^{n-1}} X^n \xrightarrow{d^n} X^{n+1} \dots$$
 st.
$$d^n \circ d^{n-1} = 0 \text{ all } n \in \mathbb{Z}.$$

Remark: Equally, a chain complex in the opposite abelian cat, in the sense of what follows shortly.

Everything about chain complexes in an abelian cat thus has dual version for cochain complexes

- The cohomology of a cochain complex X is defined as

$$H^n X := \ker d^n / \operatorname{im} d^{n-1}$$

Examples

① If X is a ch. complex of ab. groups
& A an abelian group

then $Ab(X^n, A) \xrightarrow{Ab(d^{n+1}, A)} Ab(X^{n+1}, A) \dots$
is a cochain complex in Ab .

② Recall if X a top space, can form a
chain complex in Ab

$S_n X$ whose homology is the
singular homology of X .

If A is an abelian group, the
cohomology of $Ab(S_n X, A)$ is
called

singular cohomology of X
of X with coefficients in A .

Lecture 2 - Abelian categories

- What is correct categorical context for homological algebra?
 - Need to talk about zero maps, add & subtract morphisms & form kernels, images & quotients & these should behave like in Mod_R
 - Resulting notion: abelian category

Defⁿ) A pre-additive category (Ab-enriched category)

\mathcal{C} is a cat in which hom-sets $\mathcal{C}(a, b)$ has structure of an abelian gp.

(ie. $a \begin{matrix} \xrightarrow{f} \\ \xrightarrow{g} \end{matrix} b \longmapsto a \begin{matrix} \xrightarrow{f+g} \\ \xrightarrow{g} \end{matrix} b, a \xrightarrow{-f} b, a \xrightarrow{0} b$)

& moreover pre & postcomposition preserves ab. group str:

$$\left(x \begin{matrix} \xrightarrow{r} a \\ \xrightarrow{v} \end{matrix} \begin{matrix} \xrightarrow{f} \\ \xrightarrow{g} \end{matrix} b \xrightarrow{s} y \quad \text{we have } \begin{matrix} (f+g)r = fr + gr \\ 0_{a,b}r = 0_{x,b} \end{matrix} \right)$$

$$\& \quad \begin{matrix} s(f+g) = sf + sg \\ s0 = 0 \end{matrix}$$

Example

- Mod_R is pre-additive.

Given $M \begin{matrix} \xrightarrow{f} \\ \xrightarrow{g} \end{matrix} N$, $f+g: M \rightarrow N$ is defined by $x \mapsto fx + gx$

which is ab. gp. hom. by commutativity of $+$.

$$\begin{aligned} (f+g)(rx) &= f(rx) + g(rx) = rf(x) + rg(x) \\ &= r(f(x) + g(x)) = \\ &= r \cdot (f+g)(x) \end{aligned}$$

$$-f: M \rightarrow N: x \mapsto -(fx)$$

$$0: M \rightarrow N: x \mapsto 0$$

- $s(f+g) = sf + sg$ holds as s a homom. & $(f+g)r = fr + gr$ is trivial.

Remark

\mathcal{C} is pre-additive $\Leftrightarrow \mathcal{C}^{\mathcal{P}}$ is .

Propⁿ

Let \mathcal{C} be pre-additive.

- ① Then \mathcal{C} has term. ob. $\Leftrightarrow \mathcal{C}$ has init. ob.
- ② \mathcal{C} has binary prods. $\Leftrightarrow \mathcal{C}$ has binary coproducts.

Proof

① - Let t be terminal.

- Then $t \xrightarrow{0} t = 1_t$.
- Now given x , we have $0: t \rightarrow x$ & must show it is unique, so consider $f: t \rightarrow x$.
Then $f = f \circ 1_t = f \circ 0 = 0$, so t initial.
- Converse is dual.

② Let $\begin{array}{ccc} & & a \\ & p_1 \nearrow & \\ c & & \\ & p_2 \searrow & \\ & & b \end{array}$ be a product diagram.

We have

$$\begin{array}{ccc} & \text{id} \nearrow & a \\ a & & \\ & 0 \searrow & b \end{array}$$

indices! $a \xrightarrow{i_1 = \langle \text{id}, 0 \rangle} c$

Similarly

$$\begin{array}{ccc} & 0 \nearrow & a \\ b & & \\ & \text{id} \searrow & b \end{array} \quad \text{ind. ! } b \xrightarrow{i_2 = \langle 0, \text{id} \rangle} c$$

& we have

$$\begin{array}{ccc} a & \xrightarrow{i_1} & c \\ & \nearrow i_2 & \\ b & & \end{array}$$

, which we will show is coprod. diagram.

Key point: we have diagram

$$a \begin{array}{c} \xrightarrow{i_1} \\ \xleftarrow{p_1} \end{array} c \begin{array}{c} \xleftarrow{i_2} \\ \xrightarrow{p_2} \end{array} b \quad \text{satisfying}$$

- $p_1 i_1 = 1_A, p_2 i_1 = 0$
- $p_1 i_2 = 0, p_2 i_2 = 1_B$
- $i_1 p_1 + i_2 p_2 = 1_C$

Such a diag. is called a biproduct diagram.

To see last equation,

$$p_1(i_1 p_1 + i_2 p_2) = p_1 i_1 p_1 + p_1 i_2 p_2 = 1 p_1 + 0 p_2 = p_1(1)$$

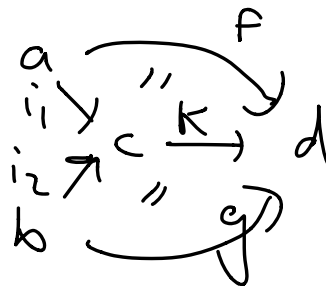
& sim for p_2 , as required.

- Any biprod. diag. is both product & coproduct - let's show coproduct.

- Indeed given



must find k st



But if we have such a k then

$$k = k \cdot 1_c = k(i_1 p_1 + i_2 p_2) = k i_1 p_1 + k i_2 p_2 = f p_1 + g p_2$$

so must have $k = f p_1 + g p_2$.

let's check that $f p_1 + g p_2$ has required

props:

$$(f p_1 + g p_2) i_1 = f p_1 i_1 + g p_2 i_1 = f + 0 = f$$

$$\dots \dots \dots i_2 \qquad \qquad \qquad = g$$

as required.

Converse is dual \square

Defⁿ) A preadd. cat is additive if it has finite products (equivalently, by above, finite coproducts).

Notation: In additive cat, write 0 for the terminal = initial object.

Example Mod_R is additive.

Prop 3 generalises fact from Alg. 3 that in Mod_R binary prods & coproducts coincide.

Kernels & quotients

consider $a \xrightarrow{f} b \in \mathcal{C}$ preadditive.

- The kernel of f $\ker f$ comes with

$\ker f \xrightarrow{i} a \xrightarrow{f} b$ & $f \circ i = 0$ & is universal with this property:
(ie. given $g: c \rightarrow a$ & $f \circ g = 0 \exists ! \bar{g}: c \rightarrow \ker f$ st $i \circ \bar{g} = g$.)

- The cokernel of f is dual: we have

$a \xrightarrow{f} b \xrightarrow{p} \text{coker } f$ s.t. $p \circ f = 0$ & is universal in dual sense.

Remark: $\ker f$ & $\text{coker } f$ are equalisers & coequalisers of

$$a \begin{array}{c} \xrightarrow{f} \\ \xrightarrow{0} \end{array} b$$

This implies $\ker f \xrightarrow{i} a$ is mono

& $b \xrightarrow{f} \text{coker} f$ is epi.

- Sometimes I'll write $\ker f = \ker F$
 $\text{coker} f = \text{coker} F$.
- If $i: a \hookrightarrow b$ is mono,
write $\text{coker}(i) := b/a$.

Example

- In $\text{Mod } R$, given $f: A \rightarrow B$,
 $\ker f = \{a \in A: fa = 0\}$
 $\text{coker} f = B/\text{im} f$.

Lemma

In an additive cat.,

- ① $A \xrightarrow{f} B$ is mono $\Leftrightarrow \ker f = 0$
- ② $A \xrightarrow{f} B$ is epi $\Leftrightarrow \text{coker} f = 0$

Proof

- By duality, it suffices to prove ①.

- Let $f: A \rightarrow B$ mono & consider

$$0 \xrightarrow{0} A \xrightarrow{f} B$$

- Given $g: C \rightarrow A$ s.t. $fg = 0$ then
 $fg = 0 = f \circ 0$ so $g = 0$; then
 $C \xrightarrow{g=0} A$ as 0 terminal.

$$!0 \vee 0 \xrightarrow{!} 0$$

Hence $0 = \ker f$.

- Conversely, let $0 \xrightarrow{0} A \xrightarrow{f} B$ be kernel
& consider $C \xrightarrow{g} A$ s.t. $fg = fh$.

Equip. $f(g-h) = 0$ so $\exists! C \xrightarrow{t} 0$ s.t.

$g \circ h = 0 \circ t = 0$ so $g = h$ & f is mono. \square

• In an additive cat with kernels & cokernels, we can factor each morphism in two ways:

(1)
$$k_f \xrightarrow{i} a \xrightarrow{f} b$$

$$p \searrow \text{c}k_f \xrightarrow{\exists! t} b$$
 where t is induced by $f \circ i = 0$.

(2)
$$a \xrightarrow{f} b \xrightarrow{q} \text{c}f$$

$$\exists! l \searrow k_{c_f} \xrightarrow{m}$$

where l is induced by $q \circ f = 0$

As $0 = q \circ f = q \circ t \circ p$ & p is epi, we have $q \circ t = 0$ so we get

! map $k_{c_f} \xrightarrow{\alpha} k_f$ such that

$$a \xrightarrow{f} b$$

$$p \searrow \text{c}k_f \xrightarrow{\alpha^{-1}} k_f \xrightarrow{m} b$$

Defⁿ An add. cat w' kernels & cokernels is abelian if

$\alpha: \text{c}k_f \rightarrow k_f$ is invertible.

• In an abelian category, the two ways of factoring f coincide (up to ! iso).

we write

$$\begin{array}{ccc}
 & f & \\
 a & \xrightarrow{\quad} & b \\
 \text{epi} \sim \underset{p}{\parallel} \overset{e}{\searrow} & \text{im} f & \xrightarrow{t} \sim \text{mono as} \\
 & \parallel & = \text{ma} \\
 & \text{ck} f \cong \text{k} c f &
 \end{array}$$

Example

In $\text{Mod } R$, given $f: A \rightarrow B$ we have
 $A \rightarrow \text{ck} f = A \rightarrow A/\text{ker} f$
 $\& B \rightarrow \text{c} f = B \rightarrow B/\text{im} f$

so $\text{k} c f \hookrightarrow A = \text{im} f$ & the induced map $\text{ck} f \xrightarrow{\alpha} \text{k} c f$ is the hom.

$$\begin{array}{ccc}
 A/\text{ker} f & \longrightarrow & \text{im} f \\
 [a] & \longmapsto & f a
 \end{array}$$

which is an iso by first iso theorem.
 In particular $\text{Mod } R$ is abelian.

Another characterisation of abelian cats is as follows.

Propⁿ \mathcal{C} is abelian \Leftrightarrow

- ① Each mono is kernel of its cokernel
- ② Each epi is cokernel of its kernel.

Remark: ② shows that epis are

coequaliser maps i.e. regular epi, so in an ab. cat

epi \equiv regular epi.

- In particular, this is true in $\text{Mod } R$.

- Remember in Rng , $\mathbb{Z} \rightarrow \mathbb{Q}$ is epi but not reg. epi.

So Rng is not abelian.

Proof of prop.

- We will only show \Rightarrow & leave \Leftarrow for interested.

- Let $A \xrightarrow{f} B$ be mono, so $\ker f = 0$.

- Then consider

$$\begin{array}{ccccccc} \ker f = 0 & \longrightarrow & A & \xrightarrow{f} & B & \longrightarrow & \text{coker } f \\ & & \downarrow 1 & \searrow f & \nearrow m & & \\ & & \text{coker } f = A & \xrightarrow[\alpha]{\cong} & \ker \text{coker } f & & \end{array}$$

where $A \xrightarrow{1} A$ is $\text{coker } f$ as $\ker f = 0$.

Therefore f is iso to $\ker \text{coker } f \rightarrow B$, as required.

- To show each epi is cokernel of kernel is dual.

□

Examples

① $\text{Mod } R$

② IF A is abelian so is A^{op} !

So can apply duality to abelian cats.

③ IF A is abelian, we can consider chain complexes in A :

$$\dots A_{n+1} \xrightarrow{d_{n+1}} A_n \xrightarrow{d_n} A_{n-1} \dots$$

$\underbrace{\hspace{10em}}_{\parallel}$

We have $\text{im}(d_{n+1}) \overset{\circ}{\hookrightarrow} \text{Ker}(d_n)$
 $\searrow \quad \swarrow$
 A_n

& so can form

$$h_n(A) := \text{Ker}(d_n) / \text{im}(d_{n+1})$$

- Thus can speak of homology of chain complexes in abelian cat.

- We can also consider exact sequences, chain maps, chain homotopies in A &

the results of last week hold in this setting.

Moreover we can form the category $\text{Ch}(A)$ of chain complexes in A & it is

again abelian:

indeed $(f+g)_n = f_n + g_n$ for chain maps $(-f)_n = -f_n$ & $(\ker f)_n = \ker(f_n)$, $(\text{coker } f)_n = \text{coker}(f_n)$ are constructed componentwise.

As before,

$H_n : \text{Ch}'(A) \longrightarrow A$
is a functor.

④ If \mathcal{C} is a small cat, then $[\mathcal{C}, A]$ the functor cat is abelian with componentwise structure.

⑤ X a top. space, the poset

of open sets $\mathcal{O}(X)$ is a cat,
& cat of presheaves valued
in $\text{Mod}_R [\mathcal{O}(X)^{\text{op}}, \text{Mod}_R]$
is abelian, as is its
full subcat of sheaves.
important in alg. geometry.

Lecture 3

Additive & exact Functors

Defⁿ) Let A, B be abelian cats. A functor $F: A \rightarrow B$ is additive if each function $F_{x,y}: A(x,y) \rightarrow B(Fx, Fy)$ is a homom. of abelian groups (i.e. $F(f+g) = Ff + Fg$ & $F_{x,y} \circ 0_{A(x,y)} = 0_{B(Fx, Fy)}$).

- From last week,

- term/in. ob are char. by diags $a \xrightarrow{0=id} a$ which we call zero object diag & a zero ob., denoted by 0 .
- bin. products/coproducts are characterised by diagrams of form

$$a \begin{array}{c} \xrightarrow{i_1} \\ \xleftarrow{p_1} \end{array} c \begin{array}{c} \xleftarrow{i_2} \\ \xrightarrow{p_2} \end{array} b \quad \text{sat.}$$

- $p_1 i_1 = id$, $p_2 i_2 = id$, $i_1 p_1 + i_2 p_2 = id_c$
- $p_1 i_2 = 0$, $p_2 i_1 = 0$,

which are called biproduct diagrams, and often denote biprod. by $a \oplus b$.

Proposition

Any additive functor F preserves finite prods & finite coproducts - i.e. zero ob. & biproducts.

Proof

F preserves zero ob. diag / biproduct diagrams \square

Defⁿ) An additive functor $F: A \rightarrow B$

bet. abelian cats is

- left exact if it pres. kernels (lex)
- right exact if it pres. cokernels (rex)
- exact if it preserves both. (ex)

Remark : Lex functors are

those preserving finite limits
(fin. prods + equalisers \equiv kernels).

Rex functors preserve finite colimits
& ex functors preserve both.

Lemma

- ① F is lex \Leftrightarrow it preserves exactness of sequences $0 \rightarrow A \xrightarrow{f} B \xrightarrow{g} C$
- ② F is rex \Leftrightarrow it preserves exactness of sequences $A \xrightarrow{f} B \xrightarrow{g} C \rightarrow 0$
- ③ F is ex \Leftrightarrow it preserves exactness everywhere
 \Leftrightarrow it pres ses $0 \rightarrow A \xrightarrow{f} B \xrightarrow{g} C \rightarrow 0$.

Proof

First we prove ①. observe $0 \rightarrow A \xrightarrow{f} B \xrightarrow{g} C$
is exact $\Leftrightarrow \ker f = 0$ (ie F is mono)

& $\ker g = f$ ($A \rightarrow \text{im } f$ is an iso)

Hence if F pres. kernels, it pres ex. of such sequences.

Conversely, consider the exact sequence

$0 \rightarrow \ker f \xrightarrow{i} A \xrightarrow{f} B$ which F preserves
 so $0 \rightarrow F\ker f \xrightarrow{Fi} FA \xrightarrow{FF} FB$ is exact
 $\Rightarrow F\ker f = \ker FF$. \square

② is dual to ①.

For ③, if F is ex. it preserves kernels & cokernels & so images. Therefore it preserves exactness everywhere & so in particular ses.

Suppose F pres ses. We will show F pres. monos & epis. if f is mono then $0 \rightarrow A \xrightarrow{f} B \rightarrow \operatorname{coker} f \rightarrow 0$ is ses which F preserves, so by ex. @ Fa , FF is mono.

Dually F preserves epis.

Now show F preserves kernels (cokernels dual.)

At general $f: A \rightarrow B$ consider ses on

$$\begin{array}{c}
 \xrightarrow{f} \\
 \nearrow i \\
 \xrightarrow{f} B
 \end{array}$$

hor.
row.

$$0 \rightarrow \ker f \xrightarrow{q} A \xrightarrow{f} \operatorname{im} f \rightarrow 0$$

Then Fq, Fi mono, Fp epi & hor. sequence exact after applying F , so

$$\begin{aligned}
 \ker(FF) &= \ker(Fi) \quad \text{as } Fi \text{ mono} \\
 \ker(FF) &= F(\ker f) \quad \text{as } F \text{ pres ses} \\
 \operatorname{im}(Fq) &= F(\operatorname{im} q) \quad \text{as } Fq \text{ mono} \\
 \operatorname{im}(Fq) &= F(\ker f) \quad \text{so } F \text{ pres. kernels. } \square
 \end{aligned}$$

Examples

① If \mathcal{C} abelian cat & $A \in \mathcal{C}$ the functor

$$\mathcal{C}(A, -) : \mathcal{C} \longrightarrow \text{Ab}$$

$$\begin{array}{ccc} X & & \mathcal{C}(A, X) \\ f \downarrow & & f_* \downarrow \\ Y & & \mathcal{C}(A, Y) \end{array} \quad \begin{array}{c} g \\ \downarrow \\ fg \end{array}$$

is additive & preserves all limits - therefore it is lex.

② If $\mathcal{C} = \text{Mod}_R$, we have

$$\text{Mod}_R(A, -) : \text{Mod}_R \longrightarrow \text{Ab} \quad \& \quad \text{this}$$

has a left adjoint $A \otimes_R - : \text{Ab} \rightarrow \text{Mod}_R$.

- Here $A \otimes_R B$ classifies functions

$$K : A \times B \longrightarrow \textcircled{C} \sim R\text{-module}$$

& $K(a, -) : B \rightarrow \textcircled{C}$ is hom. of ab. groups

& $K(-, b) : A \rightarrow \textcircled{C}$ is hom. of R -modules,

& $A \otimes_R B$ is constructed as a quotient sim.

to the tensor prod. of R -modules in Alg. 3.

In particular, as a left adjoint it preserves colimits & so is right exact.

More gen, if right adjoint is lex then left adjoint

is rex & coreflex.

③ The forgetful functor

$U: \text{Mod}_R \longrightarrow \text{Ab}$ is left exact,
since $U \cong \text{Mod}_R(R, -)$ which is
lex by ① (or directly).

④ If \mathcal{C} is abelian, so is $\text{Ch}(\mathcal{C})$

& then

$$\begin{array}{ccc} \text{Ch}(\mathcal{C}) & \xrightarrow{(-)_n} & \mathcal{C} \\ X & \xrightarrow{\quad} & X_n \end{array} \text{ is}$$

exact since kernels & cokernels
are componentwise in $\text{Ch}(\mathcal{C})$.

Theorem (Freyd - Mitchell)

If \mathcal{C} is a small abelian cat.,
 \exists a ring R & an exact fully faithful embedding

$$F: \mathcal{C} \longrightarrow \text{Mod } R$$

- We will not prove it.

Consequence :- When proving things about diagrams in an abelian cat it suffices to supp. we are working in a cat. of R -modules, since the theorem lets us view our category as a full subcat. of $\text{Mod } R$ such that the inclusion pres all structure - kernels, cokernels etc.

- For instance, suffices to prove snake lemma in $\text{Mod } R$.

The snake lemma & long exact sequence of homology

- The main reason that homology can be easily computed is the long exact sequence of homology, which we turn to now.

The snake lemma & long exact sequence of homology

Snake lemma

Given a diagram

$$\begin{array}{ccccccc}
 A & \xrightarrow{F} & B & \xrightarrow{g} & C & \longrightarrow & 0 \\
 \downarrow \alpha & & \downarrow \beta & & \downarrow \gamma & & \\
 0 & \longrightarrow & A' & \xrightarrow{f'} & B' & \xrightarrow{g'} & C'
 \end{array}$$

with exact rows in an abelian cat.,

there is an exact sequence

$$\ker(\alpha) \xrightarrow{f} \ker(\beta) \xrightarrow{g} \ker(\gamma) \xrightarrow{\delta} \operatorname{coker}(\alpha) \xrightarrow{f'} \operatorname{coker}(\beta) \xrightarrow{g'} \operatorname{coker}(\gamma)$$

Proof - Maps except for δ are ind. by u.p.'s.

- Exactness @ $\ker(\beta)$ & $\operatorname{coker}(\beta)$ are dual. So prove ex @ $\ker(\beta)$. Enough to prove it in $\operatorname{Mod} R$ by FM-Prop which assume for rest of proof.
- let $b \in \ker(\beta)$ & suppose $gb=0$. By ex. $\exists a$ st $fa=gb$. Then $f'\alpha a = \beta fa = \beta b = 0$ & as f' is mono (since bottom row exact) therefore $\alpha a = 0$ so $a \in \ker(\alpha)$ w' $fa = b$. Hence ex @ $\ker(\beta)$.
- Now we construct $\delta: \ker \gamma \rightarrow \operatorname{coker} \alpha$. Consider $x \in \ker \gamma$. As g epi $\exists x' \in B$ st $gx' = x$. Consider $\beta x'$. Then $g'\beta x' = \gamma gx' = \gamma x = 0$ so $\beta x' \in \ker(g') = \operatorname{im}(f')$ so $\exists x'' \in A'$ as f' is mono s.t. $f'(x'') = \beta x'$. Define $\delta(x) = x'' + \operatorname{im}(\alpha) \in \operatorname{coker}(\alpha)$.
- To show δ well defined, let $gy' = x$.
- Then $g(y' - x') = 0$ so by ex. of top row, $\exists a \in A$ st $fa = y' - x'$, so $\alpha a \in A'$. Then $f'\alpha a = \beta fa = \beta y' - \beta x' = f'y'' - f'x''$ so as f' mono $\alpha a = y'' - x''$. Hence $y'' + \operatorname{im} \alpha = x'' + \operatorname{im} \alpha$.

Therefore \mathcal{E} is well defined.
Easy to see it is a homomorphism
& exactness @ $\ker(\gamma)$, where (α)
are left an exercise. \square

Theorem

let $0 \rightarrow A \xrightarrow{f} B \xrightarrow{g} C \rightarrow 0$ be a ses of chain complexes in an ab. cat. Then we obtain a long exact sequence of homology

$$\dots H_{n+1}(A) \xrightarrow{H_{n+1}(f)} H_{n+1}(B) \xrightarrow{H_{n+1}(g)} H_{n+1}(C) \xrightarrow{\delta} H_n(A) \xrightarrow{H_n(f)} H_n(B) \xrightarrow{H_n(g)} H_n(C) \dots$$

Proof

For a chain complex A

$$\dots A_{n+1} \xrightarrow{d_{n+1}} A_n \xrightarrow{d_n} A_{n-1} \xrightarrow{d_{n-1}} \dots$$

we have an ind. homomorphism

$$\begin{array}{ccc} A_n / \text{im}(d_{n+1}) & \xrightarrow{d} & \text{ker}(d_{n-1}) = Z_{n-1}(A) \\ x + \text{im}(d_{n+1}) & \longmapsto & d_n x \end{array}$$

- whose kernel contains elements $x + \text{im}(d_{n+1})$ with $d_n(x) = 0$ & so is precisely $\text{ker}(d_n) / \text{im}(d_{n+1}) = H_n(A)$.
- whose cokernel is $\text{ker}(d_{n-1}) / \text{im}(d_n) = H_{n-1}(A)$

Then we obtain a comm. diag.

$$\begin{array}{ccccccc} A_n / \text{im}(d_{n+1}) & \xrightarrow{f_n} & B_n / \text{im}(d_{n+1}) & \xrightarrow{g_n} & C_n / \text{im}(d_{n+1}) & \rightarrow & 0 \\ d \downarrow & & d \downarrow & & d \downarrow & & \\ 0 & \rightarrow & Z_{n-1}(A) & \xrightarrow{f_{n-1}} & Z_{n-1}(B) & \xrightarrow{g_{n-1}} & Z_{n-1}(C) \end{array}$$

& we will show these rows are

exact, & then by snake lemma we obtain les of homology. It remains to check rows are exact.

- Ex @ $C_n / \text{im}(d_n)$, $Z_{n-1}(A)$ is easy as f_{n-1} is mono & g_n is epi.

- At $x \in Z_{n-1}(B)$, suppose $gx = 0$. Then $\exists y \in A_{n-1}$ st $fa = x$. Must show $d_{n-1}a = 0$, but as F is inj, $Fd_{n-1}a = d_{n-1}fa = d_{n-1}x = 0$ implies $d_{n-1}a = 0$ so $a \in Z_{n-1}(A)$, as required.

- At $B_n / \text{im}(d_{n+1})$ let $g(x + \text{im}(d_{n+1})) = 0$. Then $gx \in \text{im}(d_{n+1})$ so $gx = dc$ some $c \in C_{n+1}$.

Then $\exists y \in B_{n+1}$ st $gy = c$. Now consider $x - dy \in B_n$. Then

$$g(x - dy) = gx - gdy = gx - dgy = gx - dc = 0, \text{ so } \exists a \in A_n \text{ st } fa = x - dy.$$

$$\begin{aligned} \text{Then } f(a + \text{im}(d_{n+1})) &= x - dy + \text{im}(d_{n+1}) \\ &= x + \text{im}(d_{n+1}), \end{aligned}$$

as required. \square

Stop here for today.

- See exercises for application to homology groups of spheres.

Homology of algebraic structures

- Homology and cohomology provide useful invariants of topological spaces.
- Can we define them for algebraic structures?

For groups, we certainly can. Indeed, a group G gives rise to a topological space BG , which is path connected, has fundamental group $\pi_1(BG) = G$, and all higher homotopy groups equal $\pi_n(BG) = 0$.

(These properties in fact characterise BG up to homotopy equivalence.)

- Indeed, when $G = \mathbb{Z}$, $BG = \bigcirc$

& when G is the free group on n elts,

$BG =$  a wedge of n circles.

- Then we can define the homology of G as the homology of BG , and the cohomology of G w' coefficients in $A \in \mathcal{A}b$ as the cohomology of BG w' coefficients in A .

- Our goal now is to describe an approach which does not use topological spaces, but is purely algebraic -

the most general approach uses the language of derived functors.

Lecture 4

Today: projectives, proj. resolutions & derived functors.
 Notation: In ab. cat, \twoheadrightarrow for epi, \rightarrow or \hookrightarrow for mono.

Projectives (see Alg 3 - we will do it quickly here)
 Defⁿ) An obj. $A \in \mathcal{C}$ in an abelian cat is projective

if given any
 diagram

$$\begin{array}{ccc} A & \xrightarrow{\alpha} & A \\ \downarrow & \exists \alpha' & \downarrow \\ B & \xrightarrow{f} & C \end{array} \quad \text{st} \quad \begin{array}{ccc} A & \xrightarrow{\alpha'} & A \\ \downarrow & & \downarrow \\ B & \xrightarrow{f} & C \end{array}$$

Remark: This says $\mathcal{C}(A, B) \xrightarrow{f_*} \mathcal{C}(A, C) \in \text{Ab}$
 $\begin{array}{ccc} \mathcal{C}(A, B) & \xrightarrow{f_*} & \mathcal{C}(A, C) \\ \downarrow g & \longmapsto & \downarrow Fg \\ \text{Ab} & & \text{Ab} \end{array}$
 is surjective / epi, so

A is proj. $\Leftrightarrow \mathcal{C}(A, -): \mathcal{C} \rightarrow \text{Ab}$ preserves epis.

Defⁿ) \mathcal{C} has enough projectives if for each $A \in \mathcal{C}$
 $\exists X$ projective & epi $X \twoheadrightarrow A$.

Propⁿ) Mod_R has enough projectives. The projectives are the retracts of free modules.

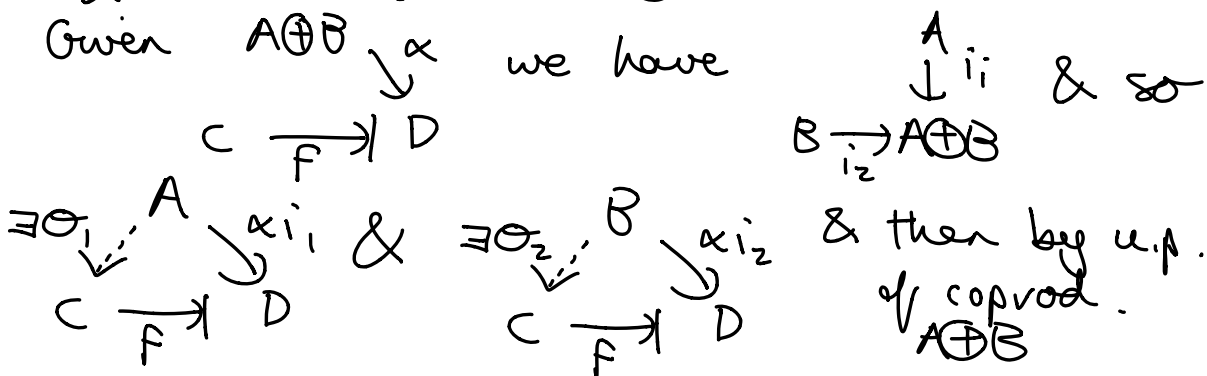
Proof See Alg. 3. Given A we take counit
 $\text{map } \mathbb{F} \circledast FUA \xrightarrow{\epsilon_A} A$ relative to
 $\text{adj } \text{Mod}_R \xrightleftharpoons[u]{\mathbb{F}} \text{Set}$, which takes formal sums to
 actual sums. \square

Proposition (Properties of projectives)

- ① Direct sums / biproducts and retracts of projectives are projective.
- ② $A \in \mathcal{C}$ is proj. $\Leftrightarrow \mathcal{C}(A, -): \mathcal{C} \rightarrow \text{Ab}$ is exact.

Proof

① is straightforward. Consider A, B proj & direct sum $A \oplus B$.



$\exists! \theta: A \oplus B \rightarrow C$ s.t. $\theta i_1 = \theta_1, \theta i_2 = \theta_2$.

Then $A \oplus B \xrightarrow{\alpha} D$ commutes using u.p. of coproduct,

Hence $A \oplus B$ projective.

- For retracts - easy: did in proof of "Proj. Mods = retracts of frees" in Alg. 3.

② If $\mathcal{C}(A, -): \mathcal{C} \rightarrow \text{Ab}$ is exact, let $B \xrightarrow{f} C$ be epi, i.e. $B \xrightarrow{f} C \rightarrow 0$ is exact.

Then let $F = \mathcal{C}(A, -)$. Then as exact functors pres. exactness,

$FB \xrightarrow{Ff} FC \rightarrow FO = 0$ is exact, so Ff is epi.

Conversely, let F pres epis. We must show F preserves ses.

Each ses is of form $0 \rightarrow \ker f \hookrightarrow A \xrightarrow{f} B \rightarrow 0$ & as F is lex it pres. kernels so above is sent to $0 \hookrightarrow \ker FF \hookrightarrow FA \xrightarrow{FF} FB \rightarrow 0$ as F preserves epis

Projective resolutions

Notation: Ch. complex X st $X_n = 0$ for $n < 0$
can be id. w' a positive chain complex

$\dots \rightarrow X_2 \xrightarrow{d} X_1 \xrightarrow{d} X_0$. write
 $\text{Ch}(\mathcal{C})_{\geq 0}$ for cat of positive chain complexes.

Def) let $A \in \mathcal{C}$ abelian cat. A proj. resolution
of A is a chain complex $C \in \text{Ch}(\mathcal{C})_{\geq 0}$
with a map $C_0 \xrightarrow{\epsilon} A$
such that:

- ① $\dots \rightarrow C_2 \xrightarrow{d} C_1 \xrightarrow{d} C_0 \xrightarrow{\epsilon} A \rightarrow 0$ is exact
- ② each C_i is projective.

Can give more conceptual definition, as
below.

Defⁿ) - A chain map $f: A \rightarrow B \in \text{Ch}(\mathcal{C})$
quasi-isomorphism if $H_n f: H_n A \rightarrow H_n B$
is an iso $\forall n \in \mathbb{Z}$.

• For $A \in \mathcal{C}$, let $A[0]$ be chain complex

$$\dots \rightarrow 0 \rightarrow 0 \rightarrow 0 \rightarrow A$$

|
degree 0

Proposition

Projective resolution C of A

\equiv
quasi-iso $C \rightarrow A[0]$ st.
each C_i is projective.

~~Proof~~ - A chain map $C \rightarrow A[0]$ is a diag.

$$\begin{array}{ccc} \vdots & & \vdots \\ C_2 & \rightarrow & 0 \\ d \downarrow & & \downarrow \\ C_1 & \rightarrow & 0 \\ d \downarrow & \varepsilon \downarrow & \\ C_0 & \rightarrow & A \end{array}$$

is specified by a single morphism $\varepsilon: C_0 \rightarrow A$ st. $\varepsilon d = 0$.

- To say it is a quasi-iso is to say

a) $H_n(C) \xrightarrow{\cong} H_n(A[0]) = 0$ for $n \geq 1$ &

$$H_0(C) \xrightarrow{H_0(\varepsilon)} H_0(A[0])$$

$$\begin{array}{ccc} \parallel & & \parallel \\ C_0 / \text{im } d & \xrightarrow{\bar{\varepsilon}} & A \\ x + \text{im } d & \longmapsto & \varepsilon x \end{array}$$

induced map is invertible

Now $\bar{\varepsilon}$ is invertible \Leftrightarrow

b) $\bar{\varepsilon}$ is inj. (its kernel $\ker \bar{\varepsilon} / \text{im } d = 0$)

c) $\bar{\varepsilon}$ is surj. (equiv, ε is surj) so

(a), (b), (c) \Leftrightarrow

$$\dots C_2 \xrightarrow{d} C_1 \xrightarrow{d} C_0 \xrightarrow{\varepsilon} A \rightarrow 0 \text{ is exact}$$

where - (c) corr. to $\text{ex}@A$

- (b) - - - - $\text{ex}@C_0$

(c) $\text{ex}@$ other positions. \square

Proposition

If \mathcal{C} has enough projectives, then each object has a projective resolution.

Proof

- Let $A \in \mathcal{C}$ & consider $C_0 \xrightarrow{\varepsilon} A \rightarrow 0$ w' C_0 proj. Then $C_0 \xrightarrow{\varepsilon} A \rightarrow 0$ exact.
- Now form

$$C_1 \begin{array}{c} \nearrow p_0 \\ \xrightarrow{d} \end{array} \begin{array}{c} \text{ker } \varepsilon \\ \searrow i \end{array} \rightarrow C_0 \xrightarrow{\varepsilon} A \rightarrow 0$$

with C_1 proj, p_0 epi.

- Then $\text{ker } \varepsilon = \text{im } p_0 \subseteq \text{im } d \subseteq \text{ker } \varepsilon$
so $\text{im } d = \text{ker } \varepsilon$.
- Now we continue in this way

$$\dots C_2 \xrightarrow{d} C_1 \begin{array}{c} \nearrow p_0 \\ \xrightarrow{d} \end{array} \begin{array}{c} \text{ker } \varepsilon \\ \searrow i \end{array} \rightarrow C_0 \xrightarrow{\varepsilon} A \rightarrow 0$$

$\searrow \text{ker } p_0 \quad \nearrow$

obtaining a projective resolution. \square

Derived functors

let $F: A \rightarrow B$ be a right exact functor
 be abelian categories.

The n'th left derived functor $L_n F: A \rightarrow B$
 is defined as follows:

@ $X \in A$, let
^{stand. terminology} $X_\bullet \xrightarrow{d} X[0]$ be a proj. resolution

$X_\bullet = \dots \xrightarrow{d} X_2 \xrightarrow{d} X_1 \xrightarrow{d} X_0$, so can form

$FX_\bullet = \dots \rightarrow FX_2 \xrightarrow{F d} FX_1 \xrightarrow{F d} FX_0$.

Then we set

$$L_n F(X) = H_n(FX_\bullet).$$

Still have to define $L_n F$ on morphisms.

Remark: Strange definition, since pr. resol.
 not functorial nor unique up to iso.
 However they are so, up to homotopy, &
 we will use this to define $L_n F$ on morphisms.

Lemma

Consider $f: A \rightarrow B \in \mathcal{C}$ & proj. resolutions

$A_\bullet \xrightarrow{d} A[0]$ & $B_\bullet \xrightarrow{d} B[0]$ of A & B .

Then \exists a chain map f_\bullet s.t. square

$$\begin{array}{ccc} A_\bullet & \xrightarrow{f_\bullet} & B_\bullet \\ d \downarrow & & \downarrow d \\ A[0] & \xrightarrow{f[0]} & B[0] \end{array}$$

* commutes, & it is unique w' this prop. up to chain homotopy.
 f in deg. 0, else 0.

Proof) - Will const. f_0 inductively.

$$\begin{array}{ccccccc}
 \dots & A_1 & \xrightarrow{d} & A_0 & \xrightarrow{d} & A & \rightarrow 0 \\
 & \textcircled{3} \downarrow f_1 & \textcircled{2} \downarrow \exists & \textcircled{1} \downarrow f_0 & \parallel & \downarrow f & \\
 \dots & B_1 & \xrightarrow{d} & B_0 & \xrightarrow{d} & B & \rightarrow 0
 \end{array}$$

Since row ex., $d: B_0 \rightarrow B$ surj. As A_0 proj., $\exists f_0$ as in $\textcircled{1}$. Next, $df_0d = Fd = 0$ so f_0d factors through $\ker d \rightarrow B_0$. As row is ex., $\text{im}(B_1) = \ker(d)$ so $B_1 \rightarrow \ker d$ surj., so as A_1 is proj., $\textcircled{2}$ factors through B_1 as a map f_1 , as in $\textcircled{3}$. Then continue in same way.

- For uniqueness, suppose we have $g_0: A_0 \rightarrow B_0$ making $*$ commute.

Then we have

$$\begin{array}{ccccccc}
 \dots & A_2 & \xrightarrow{d} & A_1 & \xrightarrow{d} & A_0 & \xrightarrow{d} A \rightarrow 0 \\
 & \textcircled{4} \downarrow h_1 & \textcircled{3} \downarrow h_0 & \textcircled{2} \downarrow f_0 - g_0 & \parallel & \downarrow f & \\
 \dots & B_2 & \xrightarrow{d} & B_1 & \xrightarrow{d} & B_0 & \xrightarrow{d} B \rightarrow 0
 \end{array}$$

Then $d(f_0 - g_0) = Fd - Fd = 0$, so $\textcircled{1}$ factors as $\textcircled{2}$ through $\ker d$ & as $B_1 \rightarrow \ker d$ is surj & A_0 proj, we obtain $\textcircled{3}$ h_0 satisfying $dh_0 = f_0 - g_0$.

Next we need a map $h_1: A_1 \rightarrow B_2$ s.t.

$dh_1 + h_0d = f_1 - g_1$ or equiv.

$dh_1 = f_1 - g_1 - h_0d := k$. Now $dk = 0$ as $df_1 - dg_1 - dh_0d = Fd - Fd - (f_0 - g_0)d = 0$, so k factors through $\ker d$ as in $\textcircled{4}$, & now as

$A_1 \text{ proj}, B_2 \rightarrow \ker d \text{ epi}$, obtain $\textcircled{5} h_1$ sat
 $dh_1 = K$, as required.
 Then continue inductively to constr h_2, \dots .

□

With this in place, we can define

$L_n F$ on morphisms:

$\textcircled{c} F: A \rightarrow B$ obtain $f_0: A_0 \rightarrow B_0$ sat. *

so $Ff_0: FA_0 \rightarrow FB_0$ inducing

$$L_n F(A) \xrightarrow{L_n F(f_0)} L_n F(B)$$

$$\parallel \quad H_n(FA_0) \xrightarrow{H_n(Ff_0)} H_n(FB_0) \quad \parallel$$

Proposition

$L_n F$ is a functor.

Proof

Firstly observe that if

$$\begin{array}{ccc} & \xrightarrow{g_0} & \\ A_0 \xrightarrow{f_0} & B_0 & \text{as in } * \text{, then by} \\ d \downarrow & \downarrow d & \text{lemma } f_0 \stackrel{h}{\simeq} g_0. \\ A[0] \xrightarrow{Ff_0} & B[0] & \end{array}$$

Then $Ff_0 \stackrel{Fh}{\simeq} Fg_0$ so

$H_n(Ff_0) = H_n(Fg_0)$ as
 homology identifies homotopic maps.
 Hence $L_n F$ is

well defined on morphisms.

Consider $A \xrightarrow{f} B \xrightarrow{g} C$. Then

$$\begin{array}{ccccc}
 A & \xrightarrow{f} & B & \xrightarrow{g} & C \\
 d \downarrow & & \downarrow d & \cong & \downarrow d \quad \text{so} \\
 A[0] & \xrightarrow{f[0]} & B[0] & \xrightarrow{g[0]} & C[0] \\
 & \searrow & \swarrow & \nearrow & \\
 & & \text{"}gf[0]\text{"} & &
 \end{array}$$

by lemma, have $g \circ f \cong (gf)$,
 so $F(g \circ f) \cong F(gf)$ so

$$\text{Hn}F(g \circ f) = \text{Hn}F(gf)$$

$$\text{"} \quad \text{L}_n F(g) \text{L}_n F(f) = \text{L}_n F(gf) \text{"}$$

Similarly, $\text{L}_n F$ pres. identities &
 so is a functor. \square

Next time, more theory of
 deriv. functors & examples -

Ext, Tor.

Lecture 5

Today :- finish left derived functors

- dual cohomology & right derived functors
- examples : Ext, Tor

Recap : For $F: A \rightarrow B$ right exact, & A

missing in video having enough projectives, defined
 $L_n F: A \rightarrow B : X \mapsto H_n(FX_*)$ *proj resolution of X*

Proposition

For $F: A \rightarrow B$ right exact functor (between abelian categories) as above, we have a natural isomorphism $L_0 F \cong F$.

~~Proof~~ At $X \in A$, we have $X \xrightarrow{d} X[0]$ & so
 $L_0 F X = H_0 F X \xrightarrow{H_0 F(d)} H_0 F X[0] = X$ gives the components of a natural transformation

$$L_0 F \rightarrow F.$$

For invertibility, as F is rex,

$$FX_1 \xrightarrow{Fd_1} FX_0 \xrightarrow{Fd_0} FX \rightarrow 0 \text{ is exact,}$$

so $\ker Fd_0 = \text{im } Fd_1$ so

$$H_0 F X_0 = FX_0 / \text{im } Fd_1$$

$$\cong FX_0 / \ker Fd_0$$

$$\cong FX \text{ as } Fd_0 \text{ epi. } \square$$

Further properties of left derived functors

① Horseshoe lemma (without proof)

- If $0 \rightarrow X \xrightarrow{f} Y \xrightarrow{g} Z \rightarrow 0 \in \mathcal{A}$ is ses we can lift it to a ses of chain complexes

$$\begin{array}{ccccccc} 0 & \rightarrow & X & \xrightarrow{f} & Y & \xrightarrow{g} & Z \rightarrow 0 \\ & & \downarrow & & \downarrow & & \downarrow \\ 0 & \rightarrow & X[0] & \xrightarrow{f[0]} & Y[0] & \xrightarrow{g[0]} & Z[0] \rightarrow 0 \end{array}$$

as on top row,

which is (moreover) componentwise

split exact: for each $n \geq 0$, of form

$$0 \rightarrow A \xrightarrow{i_A} A \oplus B \xrightarrow{p_B} B \rightarrow 0$$

for a biproduct diagram (such is always short exact).

② If $0 \rightarrow X \xrightarrow{f} Y \xrightarrow{g} Z \rightarrow 0 \in \mathcal{A}$ is ses & $F: \mathcal{A} \rightarrow \mathcal{B}$ we obtain a les

$$\dots \rightarrow L_n F Z \rightarrow L_n F X \rightarrow L_n F Y \rightarrow L_n F Z \rightarrow F X \rightarrow F Y \rightarrow F Z \rightarrow 0$$

of left derived functors:

- Indeed, F takes split exact sequences to split exact sequences (as it preserves biproducts)

- Then apply les of homology to

$$0 \rightarrow F X \xrightarrow{Ff} F Y \xrightarrow{Fg} F Z \rightarrow 0$$

③ The right exact F is exact $\Leftrightarrow L_n F = 0 \ \forall n \geq 1$
 $\Leftrightarrow L_1 F = 0$:

- If F is exact, FX_\bullet is exact sequence, so
 $H_n FX_\bullet = 0$ all $n \geq 1$, so $L_n F = 0$ all $n \geq 1$.
 - Therefore $L_1 F = 0$.

Ass. $L_1 F = 0$, at seq $0 \rightarrow X \xrightarrow{F} Y \xrightarrow{g} Z \rightarrow 0 \in \mathcal{A}$
 obtain les

$$\begin{array}{ccccccc} \dots & \rightarrow & L_1 F Z & \rightarrow & FX & \xrightarrow{FP} & FY & \xrightarrow{Fg} & FZ & \rightarrow & 0 \\ & & \parallel & & & & & & & & \\ & & 0 & & & & & & & & \end{array}$$

so $0 \rightarrow FX \xrightarrow{FP} FY \xrightarrow{Fg} FZ \rightarrow 0$
 is a seq.

Conclusion

The higher $L_n F$ for $n \geq 1$ measure the failure of F is to be exact.

Right derived functors

- An inj resolution of $X \in A$ is an exact sequence
 $0 \rightarrow X \rightarrow X^0 \rightarrow X^1 \rightarrow \dots$ with each X^i injective.

This gives a positive cochain comp. X^\bullet & morphism $X \in \mathcal{O} \rightarrow X^\bullet$ which induces an iso. on cohomology.

- If A has enough injectives, each $X \in A$ has an injective resolution.
- If $F: A \rightarrow B$ is lex functor & A has enough injectives, we can form its n 'th right derived functor $R^n F: A \rightarrow B$, which has value $R^n F(X) = H^n(FX^\bullet)$.

We have dual props to those from before:

- $R^0 F \cong F$
- For each seq $0 \rightarrow X \rightarrow Y \rightarrow Z \rightarrow 0$ obtain les
 $0 \rightarrow FX \rightarrow FY \rightarrow FZ \rightarrow R^1 FX \rightarrow R^2 FY \rightarrow \dots$

- F is exact \Leftrightarrow

$$R^n F = 0 \quad \text{all } n \geq 1 \Leftrightarrow$$

$$R^1 F = 0.$$

Ext & Tor (no proofs!)

- Begin with Tor, but focus on Ext.
- For A right R -mod, B left R -module, can form their t.p. $A \otimes_R B \in \text{Ab}$ w' universal property:

$A \otimes_R B \longrightarrow C \in \text{Ab}$ is in bijⁿ w' function $f: A \times B \longrightarrow C$ which is Ab-gp hom. in each var.

- In this way, we obtain a functor

$-\otimes_R - : R\text{-Mod} \times \text{Mod}_R \longrightarrow \text{Ab}$
 $(A, B) \longmapsto A \otimes_R B$, & so
functors $A \otimes_R - : \text{Mod}_R \longrightarrow \text{Ab}$ & $-\otimes_R B : R\text{-Mod} \longrightarrow \text{Ab}$,
and these are rex.

Then $\text{Tor}_n^A(A, B)$:= $L_n(- \otimes_R B)(A)$,
which is given by n th homology of

$$\dots \longrightarrow A_2 \otimes_R B \longrightarrow A_1 \otimes_R B \longrightarrow A_0 \otimes_R B$$

where A_\bullet is proj. resolution of $A \in R\text{-Mod}$,
so $H_n(A_\bullet \otimes_R B)$.

Equiv., it can be calc as

$L_n(A \otimes_R -)(B)$, so take proj.
resolution of B , tensor by A &
calc. homology,

then $H_n(A_\bullet \otimes_R B)$ \cong $H_n(A \otimes_R B_\bullet)$.

Ext

- Given $A \in \text{Mod}_R$, can form $\text{Mod}_R(A, -) := \text{Hom}(A, -)$ which is rex & as Mod_R has enough inj's, can form right der. functor

$\text{Ext}_n(A, -) := R^n \text{Hom}(A, -)$, so

$\text{Ext}_n(A, B) := H^n \text{Hom}(A, B^\bullet)$ for B^\bullet a inj. resolution $B_0 \rightarrow B_1 \rightarrow B_2 \dots$ of B .

- Note also have rex functor

$$\text{Hom}(-, B) : \text{Mod}_R^{\text{op}} \longrightarrow \text{Ab}$$

& so can calculate

$$R^n \text{Hom}(-, B) : \text{Mod}_R^{\text{op}} \longrightarrow \text{Ab}$$

$R^n \text{Hom}(-, B)(A)$ as H^n app to

$$\text{Hom}(A_0, B) \rightarrow \text{Hom}(A_1, B) \rightarrow \text{Hom}(A_2, B) \dots$$

for A a proj resolution of A
(i.e. inj. resolution in Mod_R^{op} .)

In fact, $\text{Ext}^n(A, B) =$

$$\underline{H^n \text{Hom}(A, B^\bullet)} \cong H^n \text{Hom}(A_\bullet, B)$$

so 2 ways of calc. $\text{Ext}^n(A, B)$.

• Ext can be understood more explicitly using extensions:
 an extension of A by B is a seq

$$0 \rightarrow B \rightarrow X \rightarrow A \rightarrow 0.$$

• Two extensions of A by B are equiv (\simeq) if \exists iso of seq of form

$$\begin{array}{ccccccc} 0 & \rightarrow & B & \rightarrow & X & \rightarrow & A \rightarrow 0 \\ & & \downarrow & & \downarrow & & \downarrow \\ 0 & \rightarrow & B & \rightarrow & Y & \rightarrow & A \rightarrow 0 \end{array}$$

• Then $\text{Ext}^1(A, B) = \{ \text{extensions of A by B} \} / \simeq$

• From earlier work,

A is proj. $(\Leftrightarrow) \text{hom}(A, -)$ is exact

$(\Leftrightarrow) \text{Ext}^1(A, B) = 0$ all B:

This says that each extension as above of A by B is iso to trivial ext.

$$0 \rightarrow B \rightarrow A \oplus B \rightarrow A \rightarrow 0,$$

or split exact.

• For higher n, elts of $\text{Ext}^n(A, B)$ are exact sequences of length n

$$0 \rightarrow B \rightarrow \dots \rightarrow A \rightarrow 0 \text{ modulo equivalence}$$

Group cohomology

• G a group, can form group (co)homology using above:

- consider category Mod_G of G -modules: such consists of an ab. group A with an assoc. action $G \times A \rightarrow A: (g, a) \mapsto g \cdot a$, equiv. a hom. $G \rightarrow \text{Ab}(A, A)$.

(Will look at these more closely in group representation theory!!)

- These are $\mathbb{Z}G$ -modules where $\mathbb{Z}G$ the group ring, whose elements are sums

$$n_1 g_1 + \dots + n_k g_k, \quad n_i \in \mathbb{Z}, g_i \in G.$$

• The inclusion

$\text{Ab} \xrightarrow{i} G\text{-Mod}$ sends A to same ab. group A w' trivial action $g \cdot a = a$.

• We have adjunctions $\text{Coinv} \dashv i \dashv \text{Invar}$.

• Here $\text{Invar}(M) = \{x \in M : gx = x \text{ all } g \in G\}$ or equally the limit of the diagram

$$\begin{array}{ccc} \Sigma G & \xrightarrow{M} & \text{Ab} \\ \downarrow \text{ob. cat.} & \cdot \mathcal{P}g & \hookrightarrow & M \mathcal{P}g \end{array}$$

ie.

$$\begin{array}{ccc} \text{Invar}(M) & \hookrightarrow & M \\ & \searrow & \downarrow \text{" } \mathcal{P}g \text{ } \cup g \\ & & M \end{array}$$

• Its colimit is $\text{Coinvar}(M) = M / \langle gx - x : x \in M, g \in G \rangle$.

• In partic.,

$\text{Invar} : \mathbb{Z}G\text{-Mod} \rightarrow \text{Ab}$ is lex,

& the group cohomology of G w' coefficients in M is defined as

$R^n(\text{Invar}(-))(M)$. Note, however that

$\text{Invar}(A) \cong \mathbb{Z}G\text{-Mod}(i\mathbb{Z}, A)$, so

$R^n(\text{Invar}(-))(M) \cong \text{Ext}_n(i\mathbb{Z}, M)$, which can be calc. using proj. resolution of $i\mathbb{Z}$ called Bar resolution.

Group cohomology described in this way captures extensions, connections with crossed homomorphisms, factor systems ... arising in group theory.

• When M is an ab. group, viewed as triv. G -mod. iM , then this coincides w' cohomology of top. space BG w' coefficients in M described earlier.

Projective dimension

Defⁿ) R a ring & $A \in \text{Mod } R$. A proj. resolution of form

$\dots \rightarrow 0 \rightarrow P_n \rightarrow \dots \rightarrow P_1 \rightarrow P_0 \rightarrow A \rightarrow 0$
with all $P_i \neq 0$ for $i \leq n$ is said to have length n .

- The projective dimension, $\text{pd}(A)$, is the smallest n for which A has such a projective resolution.
- If no such $n \in \mathbb{N}$ exists, we say $\text{pd}(A) = \infty$.

Notation: For a proj. res. P of A as above, $K_n = \ker(P_n \rightarrow P_{n-1})$ are called n -syzygies.

Examples

① For A an abelian group, consider

$$0 \rightarrow \ker(\varepsilon) \hookrightarrow \text{free ab. group} \xrightarrow{\varepsilon} A \rightarrow 0$$

- Then as each subgroup of free ab. group is free

$\ker(\varepsilon)$ is free, so A has proj. resolution of length 1.

② In the ex., we saw that

when $R = \mathbb{Z}/4$, $\mathbb{Z}/2$ has a resolution

$$\dots \rightarrow \mathbb{Z}/4 \xrightarrow{\cdot 2} \mathbb{Z}/4 \xrightarrow{\cdot 2} \mathbb{Z}/4 \rightarrow \mathbb{Z}/2 \rightarrow 0$$

of infinite length.

- The following shows close relationship between projective dimension & Ext.

Prop TFAE

- ① $\text{pd } A \leq n$
- ② $\text{Ext}^k(A, B) = 0$ all B & $k \geq n+1$.
- ③ $\text{Ext}^k(A, B) = 0$ all B & $k = n+1$
- ④ For each proj. res. P of A , we have P_{n-1} is projective.

Proof

- Assume ① & let P be proj. resolution.

Then $\text{Ext}^k(A, B)$ is k 'th cohomology of

$$\dots \rightarrow \text{Hom}(P_{k-1}, B) \rightarrow \text{Hom}(P_k, B) \rightarrow \text{Hom}(P_{k+1}, B) \rightarrow \dots$$

but as $P_k = 0$, $\text{Hom}(P_k, B) = 0$
so $\text{Ext}^k(A, B) = 0$, proving (2).

(2) \Rightarrow (3) trivially.

- For (3) \Rightarrow (4), observe

$P_{n+2} \rightarrow P_{n+1} \rightarrow P_n \rightarrow K_{n-1} \rightarrow P_{n-1}$
is proj. res. of K_{n-1} .

Hence $\text{Ext}^i(K_{n-1}, B)$ is cohomology of
 $\text{Hom}(P_n, B) \rightarrow \text{Hom}(P_{n+1}, B) \rightarrow \text{Hom}(P_{n+2}, B)$.

so $\text{Ext}^i(K_{n-1}, B) = \text{Ext}^{n+i}(A, B) = 0$;
hence K_{n-1} is projective.

For (4) \Rightarrow (3), take proj. resolution

$$0 \rightarrow K_n \hookrightarrow P_{n-1} \rightarrow P_{n-2} \rightarrow \dots \rightarrow P_1 \rightarrow A.$$

Cor $\text{pd}(A) = n$ if n is the largest number
st \exists non-zero $\text{Ext}^n(A, B)$.

Remark • In the resolution

$$\dots \rightarrow \mathbb{Z}/4 \xrightarrow{\cdot 2} \mathbb{Z}/4 \xrightarrow{\cdot 2} \mathbb{Z}/4 \rightarrow \mathbb{Z}/2 \rightarrow 0$$

$K_n = \mathbb{Z}/2$, which is not proj.

as $\mathbb{Z}/4$ -module (check.)

- Hence $\mathbb{Z}/2$ inf. proj d.

Calculation of Ext for abelian groups

- Of course $\text{Ext}_{\mathbb{Z}}^0(A, B) = \text{Ab}(A, B)$
- By the first, ex., since $\text{pd}(A) \leq 1$, $\text{Ext}_{\mathbb{Z}}^n(A, B) = 0$ all $n \geq 2$.
- $\text{Ext}_{\mathbb{Z}}^1(A, B)$ is more complicated but we can easily calculate it if A is f.g. abelian group, since these are of form $A = \mathbb{Z}^m \oplus \mathbb{Z}_{p_1} \oplus \dots \oplus \mathbb{Z}_{p_k}$.
- Not hard to see $\text{Ext}^1(-, B)$ pres direct. sums & since \mathbb{Z}^m is free (so proj) $\text{Ext}_{\mathbb{Z}}^1(A, B) = \text{Ext}^1(\mathbb{Z}_{p_1}, B) \oplus \dots \oplus \text{Ext}^1(\mathbb{Z}_{p_k}, B)$.
- Now $0 \rightarrow \mathbb{Z} \xrightarrow{-p} \mathbb{Z} \rightarrow \mathbb{Z}/p \rightarrow 0$ is ses giving proj. res of \mathbb{Z}/p , so $\text{Ext}^1(\mathbb{Z}/p, B)$ is cohomology of $\begin{array}{ccc} \text{hom}(\mathbb{Z}, B) & \xrightarrow{(-p)^*} & \text{hom}(\mathbb{Z}, B) \rightarrow 0 \\ \text{||} & & \text{||} \\ B & \xrightarrow{-p} & B \rightarrow 0 \end{array}$
ie. $\underline{B/pB}$.
- Hence $\text{Ext}_{\mathbb{Z}}^1(A, B) = B/p_1B \oplus \dots \oplus B/p_kB$.
- If A is not fin. gen, more complex.

Defⁿ The left projective dimension of a ring R is defined by
$$\text{lpd}(R) = \sup \{ \text{pd}(A) : A \text{ a left } R\text{-mod} \}$$

Remark • Dually, one can consider injective dimension,

$\text{id}(A)$, the shortest length of an injective resolution of A .

• Since $\text{Ext}^k(A, B)$ can be calc. using a proj. res. of A or an inj. res. of B , the above result has a dual version for inj. dimension.

• Again, one can consider left inj. dimension $\text{lid}(R)$. In fact,

Prop $\text{lpd}(R) = \text{lid}(R)$

Proof $\text{lpd}(R) \leq n \iff$

$\exists A \in \text{Mod}_R, \text{pd}(A) \leq n \iff$

$\forall A, B \text{ Ext}_R^k(A, B) = 0 \text{ all } A, B, k > n$

$\Leftrightarrow \forall B \in \text{Mod}_R, \text{id}(B) \leq n$

$\Leftrightarrow \text{lid}(R) \leq n$

The common value

$\text{lgl} := \text{lpd}(R) = \text{lid}(R)$ is called the left global dimension of R .

• Of course, there is also right global dimension $\text{rgd}(R)$, concerning right R -modules.

Of course, if R is commutative, they coincide, & we just speak of global dimension $\text{gd}(R)$.

Hilbert's Syzygy Theorem

IF R is a field, $\text{gd } R[x_1, \dots, x_n] = n$.

Exercise

- Consider the poly. ring $R = \mathbb{C}(x, y)$ of length ∞ & let M be the max. ideal of polys with no zero term, and consider R/M as an R -module
- Consider the projection $R \rightarrow R/M$. Extend this to a proj. resolution of length 2.