

Reminder on weighted colimits

$$\begin{array}{c}
 X \times_{\mathcal{G}} Y \\
 (xg, yg) = (x, yg) \\
 X \times_{\mathcal{G}} Y \\
 \downarrow \text{act}_x \times 1 \quad \downarrow 1 \times \text{act}_y \\
 X \times Y
 \end{array}$$

$$\begin{array}{c}
 M \otimes_{\mathcal{R}} N \\
 x \otimes y = x \otimes y \\
 M \otimes_{\mathcal{R}} N \\
 \downarrow \downarrow \\
 M \otimes N \\
 M: \mathcal{R}^{\text{op}} \rightarrow \mathcal{A}b \\
 N: \mathcal{R} \rightarrow \mathcal{A}b
 \end{array}$$

$$\begin{array}{c}
 W \otimes_{\mathcal{A}} D \\
 \sum_{b,c} W_c \otimes \mathcal{A}(b,c) \otimes D_b \\
 \downarrow \downarrow \\
 \sum_a W_a \otimes D_a \\
 W: \mathcal{A}^{\text{op}} \rightarrow \mathcal{V} \\
 D: \mathcal{A} \rightarrow \mathcal{M}
 \end{array}$$

$$\begin{array}{c}
 \Delta^i \times_{\Delta} X \\
 (\theta^i t, x) \sim (t, \theta_* x) \\
 \Delta^i \times_{\Delta} X \\
 \downarrow \downarrow \\
 \sum \Delta^i \times X_i \\
 \Delta: \Delta \rightarrow \text{sSet} \\
 X: \Delta^{\text{op}} \rightarrow \text{Top} \dots
 \end{array}$$

$$W \otimes_{\mathcal{A}} D = \int^{A \in \mathcal{A}} W_{\mathcal{A} \otimes D A}$$

$$\begin{array}{l}
 \otimes: \mathcal{M} \times \mathcal{N} \rightarrow \mathcal{P} \\
 \otimes_{\mathcal{A}}: [\mathcal{A}^{\text{op}}, \mathcal{M}] \times [\mathcal{A}, \mathcal{N}] \rightarrow \mathcal{P}
 \end{array}$$

$$\sum_{b,c} W_c \otimes \mathcal{A}(b,c) \otimes D_b \implies \sum_a W_a \otimes D_a \xrightarrow{\text{coend}} W \otimes_{\mathcal{A}} D$$

$$/ \mathcal{M}(-, N)$$

$$\prod_{b,c} \{ \mathcal{A}(b,c), \mathcal{M}(W_c \otimes D_b, N) \} \leftarrow \prod_a \mathcal{M}(W_a \otimes D_a, N) \leftarrow \mathcal{M}(W \otimes_{\mathcal{A}} D, N)$$

$$\prod_{b,c} \{ \mathcal{A}(b,c), \mathcal{V}(W_c, \mathcal{M}(D_b, N)) \} \leftarrow \prod_a \mathcal{V}(W_a, \mathcal{M}(D_a, N)) \leftarrow [\mathcal{A}^{\text{op}}, \mathcal{V}](W, \mathcal{M}(D, N))$$

$$\mathcal{M}(D, N) \stackrel{\text{Yoneda}}{=} [\mathcal{A}^{\text{op}}, \mathcal{V}](\mathcal{A}(-, A), \mathcal{M}(D, N)) \stackrel{\text{adj}}{=} \mathcal{M}(\mathcal{A}(-, A) \otimes_{\mathcal{A}} D, N)$$

Example:

$$\mathcal{A}(-, A) \otimes_{\mathcal{A}} D = DA$$

dually

$$\{ \mathcal{A}(A, -), D \}_{\mathcal{A}} = DA$$

$$\mathcal{A}(-, -) \otimes_{\mathcal{A}} D = D$$

$$\mathcal{R} \otimes_{\mathcal{R}} M = M$$

$$\text{Hom}_{\mathcal{R}}(\mathcal{R}, M) = M$$

$$\left. \begin{array}{l}
 \mathcal{A}^{\text{op}} \rightarrow [\mathcal{B}, \mathcal{V}] \\
 W: \mathcal{A}^{\text{op}} \times \mathcal{B} \rightarrow \mathcal{V} \\
 D: \mathcal{A} \rightarrow \mathcal{M}
 \end{array} \right\} W \otimes_{\mathcal{A}} D: \mathcal{B} \rightarrow \mathcal{M}$$

$$\leq M_{\mathcal{R}} \otimes_{\mathcal{R}} \mathcal{R}^{\mathcal{N}}$$

Restriction of scalars: $F: \mathcal{B} \rightarrow \mathcal{A}$

$$\mathcal{A}(-, F-) : \mathcal{A}^{\text{op}} \times \mathcal{B} \rightarrow \mathcal{V}$$

$$D: \mathcal{A} \rightarrow \mathcal{M}$$

$$\mathcal{A}(-, F-) \otimes_{\mathcal{A}} D = DF$$

dually

$$\{ \mathcal{A}(F-, -), D \}_{\mathcal{A}} = DF$$

$$\mathcal{R} \rightarrow \mathcal{S}$$

$$\begin{array}{l}
 \mathcal{R} \otimes_{\mathcal{R}} \mathcal{V} = \mathcal{V}^{\mathcal{R}} \\
 \mathcal{S} \otimes_{\mathcal{S}} \mathcal{M} = \mathcal{M}
 \end{array}$$

$$F^* D$$

Adjunction:

$$[\mathcal{B}, \mathcal{V}](C, \{ \mathcal{A}(F-, -), D \}_{\mathcal{A}}) = [\mathcal{A}, \mathcal{V}](\mathcal{A}(F-, -) \otimes_{\mathcal{B}} C, D)$$

$$\text{Hom}_{\mathcal{R}}(M, \text{Hom}_{\mathcal{S}}(S, N)) = \text{Hom}_{\mathcal{S}}(S \otimes_{\mathcal{R}} M, N) \quad \text{F.i.C extension of scalars}$$

$$(\text{like } S \otimes_{\mathcal{R}} L)$$

$$[\mathcal{B}, \mathcal{V}](C, F^* D) = [\mathcal{A}, \mathcal{V}](F_! C, D)$$

$$\leftarrow \text{Lan}_F C$$

left Kan extension

eg. $F: \mathcal{B} \hookrightarrow \mathcal{A}$ full embedding

$$F_! C = \mathcal{A}(F-, B) \otimes C$$

eg. $F: B \hookrightarrow A$ full embedding

$$F_* C(B) = A(F-, B) \otimes_B C$$

$$= B(-, B) \otimes_B C$$

$$= C_B \quad \text{real extension}$$

$$\begin{array}{ccc} B & \xrightarrow{c} & M \\ F \downarrow & \parallel & \nearrow \\ A & & \text{Lan}_F C \end{array}$$

fully faithful

$$B \xrightarrow{F} A \xrightarrow{D} M$$

$$A \xrightarrow{D} M \xrightarrow{1} M$$

$$W \otimes_A D = \dots \otimes_A 1$$

$$W \otimes_B (D \circ F) = \bar{W} \otimes_A D$$

$$W \otimes_B \left(\underbrace{A(-, F-)}_{D \circ F} \otimes_A D \right) = \left(\underbrace{W \otimes_B A(-, F-)}_{\bar{W} = \text{Lan}_F W} \right) \otimes_A D$$

restricted from S to R

$$M \otimes_R N = M \otimes_R (S \otimes_S N)$$

$$= (M \otimes_R S) \otimes_S N$$

Reedy model categories, framings

$$M(M, N) \quad \downarrow \cong$$

Motivation. Any model category M is in some weak sense enriched over $sSet$ and, as a result, $Ho(M)$ will be enriched over $Ho(sSet)$.

Start with an honest $sSet$ -model category = simplicial model category

— since $K \in sSet$ is a colimit $K = \text{colim}_{(n, \Delta^n \rightarrow K)} \Delta^n = K \cdot_{\Delta} \Delta^*$

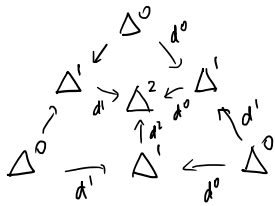
we have $K \otimes M = \text{colim} (\Delta^n \otimes M)$ and it is enough to give $\Delta^* \otimes M$

— clearly $\Delta^0 \otimes M \cong M$, since Δ^0 is the monoidal unit

— what is the essential property of $\Delta^1 \otimes M$?
it is a cylinder! ... at least for $M \in M_c$:

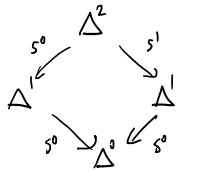
$$\Delta^0 + \Delta^0 \xrightarrow{[d^0, d^1]} \Delta^1 \xrightarrow{s^0} \Delta^0 \quad \rightsquigarrow \quad M + M \xrightarrow{\quad} \Delta^1 \otimes M \xrightarrow{\sim} M$$

— what about $\Delta^2 \otimes M$?



, more compactly $\partial \Delta^2 \rightarrow \Delta^2$
 \parallel
 $L_2 \Delta^*$

$\rightsquigarrow L_2 \Delta^* \otimes M \rightarrow \Delta^2 \otimes M \xrightarrow{\sim} M \otimes M$



, more compactly $\Delta^2 \xrightarrow{\sim} M_2 \Delta^*$

$\bigvee_0 \Delta^i: \Delta \rightarrow sSet$
a cosimplicial object in $sSet$

— need some calculus of such diagrams $\Delta \rightarrow sSet \rightsquigarrow \Delta^* \otimes M: \Delta \rightarrow M$
 \rightarrow Reedy categories, Reedy model structures $\Delta \rightarrow M$ a cosimplicial object in M

A Reedy category has two kinds of maps — direct and inverse (either d^i and s^i in Δ)

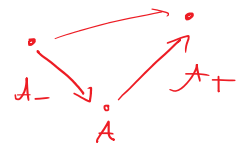
Definition. A **direct category** is a category \mathcal{A} together with a functor $\text{deg}: \mathcal{A} \rightarrow \lambda$ that satisfies $f: A \rightarrow B \rightsquigarrow f = 1 \Leftrightarrow \text{deg } A = \text{deg } B$
 \uparrow ordinal $\Rightarrow \text{deg } A < \text{deg } B$

An **inverse category** is a dual notion, i.e. a category \mathcal{A} together with a functor $\text{deg}: \mathcal{A}^{op} \rightarrow \lambda$ satisfying $f = 1 \Leftrightarrow \text{deg } A = \text{deg } B$.

A **Reedy category** is a category \mathcal{A} together with two subcategories $\mathcal{A}^+, \mathcal{A}_-$ and a function $\text{deg}: \text{ob } \mathcal{A} \rightarrow \lambda$ that

- makes \mathcal{A}^+ into a direct category
- makes \mathcal{A}_- into an inverse category
- any morphism has a unique decomposition:

$$\sum_{\mathcal{A} \leftarrow \mathcal{A}^+} \mathcal{A}^+(A, C) \times \mathcal{A}_-(B, A) \xrightarrow{\cong} \mathcal{A}(B, C)$$



any morphism ...

$$\sum_{A \in \mathcal{A}} A^+(A, C) \times A_-(B, A) \xrightarrow{\cong} A(B, C)$$

$\begin{matrix} \circ - \downarrow \circ \\ A \end{matrix}$

as a composition of an inverse and a direct morphism

The Yoneda lemma gives $X \cong A(-, -) \cdot X$ and we will describe a way of building X by decomposing $A(-, -)$. The axioms actually give

$$A(-, -) = \sum_{A \in \mathcal{A}} A^+(A, -) \times A_-(-, A) = \sum_{A \in \mathcal{A}} A^+_A \times A^-_A \quad (\text{but only of functors } A^{op} \times A^+ \rightarrow \text{Set } \nabla)$$

but in order to understand $A(-, -)$ we factor it into $st_n A(-, -)$... maps that factor through some A of degree $\leq n$ and clearly we have

$$A(-, -) = \text{colim}_{n < \infty} st_n A(-, -)$$

so that it remains to study the difference between $st_n A$ and $st_{n-1} A$ or, better for n limit, $st_{\infty} A = \text{colim}_{i < \infty} st_i A$.

Quite clearly, we have a pushout square

$$\begin{array}{ccc} 0 & \longrightarrow & st_{\infty} A \\ \downarrow & & \downarrow \\ \sum_{A \in \mathcal{A}_n} A^+_A \times A^-_A & \longrightarrow & st_n A \end{array}$$

with A ranging over all objects of degree n

However, it will be crucial to express this in terms of representable functors on \mathcal{A} , rather than on A^+, A_- .

Examples.

- any ordinal is a direct category
 - $\begin{matrix} + \rightarrow \\ \leftarrow \\ + \rightarrow \end{matrix}$ is a direct category
 - $\begin{matrix} \leftarrow \\ \circ \\ \rightarrow \end{matrix}$ is a Reedy category
 - Δ^+ is a direct category (non-empty finite ordinals + monos)
 - Δ_- is an inverse category (non-empty finite ordinals + epis)
 - Δ is a Reedy category $\Rightarrow \Delta^{op}$ Reedy
- $A \text{ direct} \Leftrightarrow A^{op} \text{ inverse}$
 $A \text{ Reedy} \Leftrightarrow A^{op} \text{ Reedy}$

Notation. We denote

- $A_A = A(A, -) \in [A, \text{Set}]$ the covariant representable
- $A^A = A(-, A) \in [A^{op}, \text{Set}]$ the contravariant representable ... think Δ^n

We further define two subfunctors

- $i_A: \circ A_A \subseteq A_A$ of maps that factor through an object of lower degree.
- $i^A: \circ A^A \subseteq A^A$ of maps that factor through an object of lower degree.

In the decomposition

$$\begin{aligned} A_X &= A(A, -) = \sum A^+(B, -) \times A_-(A, B) \\ \circ A_A &= \sum A^+(B, -) \times \circ A_-(A, B) \end{aligned}$$

\uparrow only 1 is excluded

This means that there is a pushout (coproduct)

$$\begin{array}{ccc} 0 & \longrightarrow & \partial A_A \\ \downarrow & & \downarrow i_A \\ A_A^+ & \longrightarrow & A_A \end{array}$$

Dually, we get

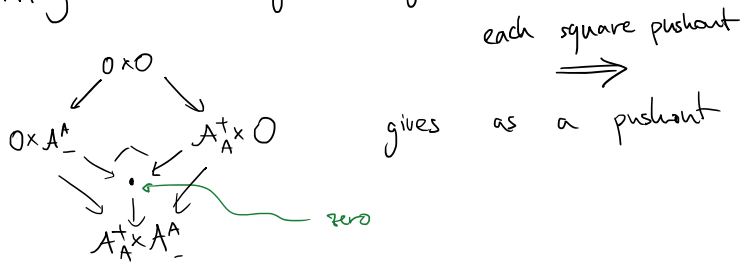
$$\partial A^A = A(-, A) = \sum A^+(B, A) \times A_-(-, B)$$

$$\partial A^A = \sum \partial A^+(B, A) \times A_-(-, B)$$

and a pushout (coproduct)

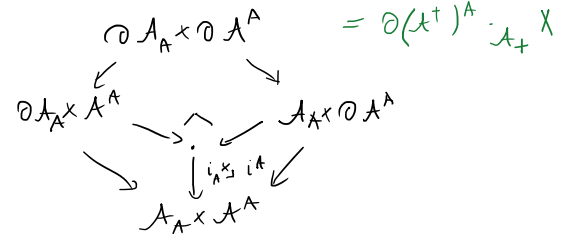
$$\begin{array}{ccc} 0 & \longrightarrow & \partial A^A \\ \downarrow & & \downarrow i^A \\ A_-^+ & \longrightarrow & A^A \end{array}$$

Putting these together yields that



$\partial A^+(B, A)$ by Yoneda

$$\begin{aligned} &= \sum \partial(A^+)^A \times_{A^+} A^+(B, -) \times A_-(-, B) \\ &= \partial(A^+)^A \times_{A^+} \sum A^+(B, -) \times A_-(-, B) \\ &= \partial(A^+)^A \cdot i^+ A(-, -) \cdot A(-, -) \\ &= \text{extension of } \partial(A^+)^A \\ &\implies \partial A^A \cdot X = (\partial(A^+)^A \times_{A^+} A(-, -)) \cdot X \end{aligned}$$



Now we get a diagram

$$\begin{array}{ccc} 0 & \longrightarrow & \sum_A \partial A_A \times A^A + \frac{A_A \times \partial A^A}{\partial A_A \times \partial A^A} \longrightarrow \text{sk}_{<n} A(-, -) \\ \downarrow & & \downarrow \sum_A i_A^+ \cdot i^A \\ \sum_A A_A^+ \times A_-^+ & \longrightarrow & \sum_A A_A \times A^A \longrightarrow \text{sk}_n A(-, -) \end{array}$$

sums range over $A \in A$ with $\text{deg } A = n$

in which the outer square is a pushout \implies so is the one on the right but now in $[A^{\text{op}} \times A, \text{Set}]_0^\triangleright$

Upon applying $- \cdot X$, we denote:

$$\begin{array}{ccc} \partial A^A \cdot X & \xrightarrow{i^+ \cdot X} & A^A \cdot X \\ \parallel \text{def} & & \parallel \\ L_A X & \xrightarrow{\partial_A X} & X_A \end{array}$$

latching object

Example. $A = \Delta^{\text{op}}$, $M = \text{Set}$

$\implies L_n X \in X_n$ the subset of deg. simplices (needs a bit of work, see above).

$$0 = \text{sk}_{-1} X$$

$$X = \text{colim sk}_n X \text{ and}$$

Theorem. For any $X \in [A, M]$ we get

$$\begin{array}{ccc} \sum_A \partial A_A \cdot X_A + \frac{A_A \cdot L_A X}{\partial A_A \cdot L_A X} & \longrightarrow & \text{sk}_{<n} X \\ \downarrow i_A^+ \cdot \partial_A X & & \downarrow \\ \sum_A A_A \cdot X_A & \longrightarrow & \text{sk}_n X \end{array}$$

Important special case.

When A is direct, we have $\partial A_A = 0$ and, consequently,

$$\begin{array}{ccc} \sum_A A_A \cdot L_A X & \longrightarrow & \text{sk}_{<n} X \\ \parallel & & \downarrow \end{array}$$

proj. model structure generated by

$$\begin{array}{ccc} \sum_A A_A \cdot L_A X & \longrightarrow & \text{st}_{<n} X \\ \downarrow & & \downarrow \\ \sum_A A_A \cdot X_A & \longrightarrow & \text{st}_n X \end{array}$$

proj. model structure generated by $A_A \cdot K \rightarrow A_A \cdot L$
 $A \in \mathcal{A}, K \rightarrow L$

so that: $\forall A \in \mathcal{A}: \partial_A X: L_A X \rightarrow X_A$ cofibration

$\Rightarrow X$ is cofibrant in the projective model structure

More generally, for $f: X \rightarrow Y$ we denote $\partial_A f = i_A \cdot j_A f$, i.e. the pushout corner map in

$$\begin{array}{ccc} L_A X & \longrightarrow & L_A Y \\ \partial_A X \downarrow & \searrow & \downarrow \partial_A Y \\ X_A & \longrightarrow & Y_A \end{array} \quad \partial_A f = \begin{array}{ccc} \partial_A^A & & X \\ \downarrow i_A & & \downarrow j_A \\ A^A & & Y \end{array}$$

$$X = \text{st}_{<n}^X Y \quad Y = \text{colim}_{n < \infty} \text{st}_n^X Y \quad \text{and}$$

Theorem. For any map $f: X \rightarrow Y$ of $[A, M]$ we get

$$\begin{array}{ccc} \longrightarrow & \text{st}_{<n}^X Y & \\ \downarrow \Sigma i_A \cdot j_A \cdot \partial_A f & \searrow & \downarrow \\ \longrightarrow & \text{st}_n^X Y & \end{array} \quad \begin{array}{ccc} X & & X \\ \downarrow f & & \downarrow \\ Y & & Y \end{array} = \begin{array}{ccc} X & & X \\ \downarrow & & \downarrow \\ \text{st}_n^X Y & = & \text{st}_n Y + \text{st}_n X \\ \downarrow & & \downarrow \\ Y & & Y \end{array}$$

Theorem. For a map $f: X \rightarrow Y$ of $[A, M]$ TFAE

- $\forall A: \partial_A f$ is a (trivial) cofibration
- $f \in \text{cell} \{ i_A \cdot j_A \mid A \in \mathcal{A}, i_A \in M \}$ a (trivial) cofibration

$$\begin{array}{c} \partial_A A \cdot L + \partial_A A \cdot K \quad A_A \cdot K \\ \downarrow \\ A_A \cdot L \end{array}$$

These maps are called **Reedy (trivial) cofibrations**.

Proof. The implication " \Downarrow " is the previous theorem.

For the implication " \Uparrow " observe that ∂_A commutes with all allular constructions so that it is enough to study $\partial_A (i_B \cdot j_A)$:

$$\partial_A^A \cdot j_A \left(\begin{array}{ccc} \partial_A B & & K \\ \downarrow i_B & \searrow & \downarrow i \\ A_B & & L \end{array} \right) = \left(\begin{array}{ccc} \partial_A^A & & \partial_A B \\ \downarrow i_A & \searrow & \downarrow i_B \\ A^A & & A_B \end{array} \right) \cdot j_A \cdot \downarrow i$$

either iso for $A \neq B$

$$A_{\neq 1}(B, A) \hookrightarrow A(B, A) \text{ for } A = B$$

pushout of $0 \rightarrow 1$

pushout of i

□

Dually, we denote

$$\begin{array}{ccc} M_A X = \{ \partial_A X_A \mid X \} & & \\ \partial_A X \uparrow & \uparrow i_A \cdot j_A & \\ X_A = \{ A_A \cdot X \} & & \end{array}$$

and more generally for $f: X \rightarrow Y$:

$$X_A \longrightarrow M_A X$$

$S_{A^*} f = \{i_A, f\}_{A^*} =$ pullback corner map in

$$\begin{array}{ccc} \wedge A & \cdots & \wedge A \\ \downarrow & & \downarrow \\ Y_A & \longrightarrow & M_A Y \end{array}$$

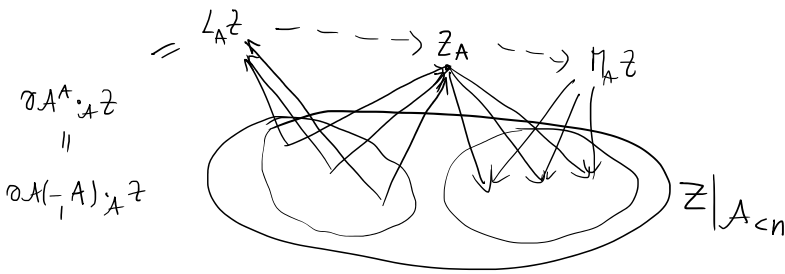
Theorem TFAE

- $\forall A \in A: S_{A^*} f$ is a (trivial) fibration
 - $f \in \text{cocell} \{ \{i_A, p\}_{A^*} \mid A \in A, p \in M \text{ a (trivial) fibration} \}$
- \leadsto **Reedy (trivial) fibrations.**

Theorem. There is a model structure on $[A, M]$, called the **Reedy model structure** with $\mathcal{C} =$ Reedy cofibrations, $\mathcal{F} =$ Reedy fibrations, $\mathcal{W} =$ pointwise weak equivalences.

Proof. We need to show that $\mathcal{W} \cap \mathcal{C}$ are exactly the Reedy \leftarrow *inductively* fibrations by definition. Easily $i_A \circ j \in \mathcal{C} \iff j \in \underbrace{\{i_A, f\}_{A^*}}_{\text{fibrations by definition}}$.

The factorizations are produced inductively using the following idea:



to give an extension of z from $A_{<n}$ to $A_{\le n}$ we need $\forall A$ of degree n to factor

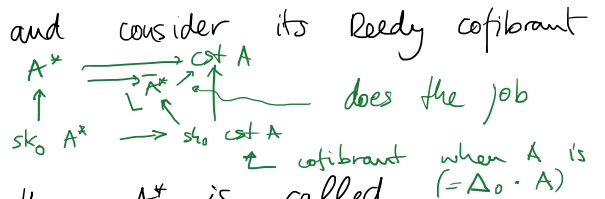
$$\begin{array}{ccccc} L_A Z & \longrightarrow & Z_A & \longrightarrow & M_A Z \\ & & \searrow & & \nearrow \\ & & & \text{obtained from } z|_{A_{<n}} & \end{array}$$

Now

$$\begin{array}{ccccc} L_A X & \longrightarrow & L_A Z & \longrightarrow & L_A Y \\ \downarrow & & \swarrow & & \downarrow \\ X_A & \longrightarrow & Z_A & \longrightarrow & Y_A \\ \downarrow & & \swarrow & & \downarrow \\ M_A X & \longrightarrow & M_A Z & \longrightarrow & M_A Y \end{array}$$

Application. • Properness in M_c / M_f differently.

• Given $A \in M_c$ consider $\text{cst } A \in [\Delta, M]$ and consider its Reedy cofibrant replacement $A^* \in [\Delta, M]$



$0 \rightarrow A^* \xrightarrow{\sim} \text{cst } A$

- we may achieve that $A^*_0 = A$ and then A^* is called a **cosimplicial frame** on A . Any two are related by a zig-zag of cosimplicial frames. Thus for $X \in M_f$

$M(A^*, X) \in \text{sSet}$

Dually $M(A, X_*) \in \text{sSet}$ via a simplicial frame on X .

$A^* : \Delta \rightarrow M$
 $K : \Delta^{op} \rightarrow \text{Set}$ } get $K : \Delta A^* \in M$, an action of sSet on M
 $K * A \stackrel{\text{def}}{=} K \cdot \Delta A^*$, but not associative

The bifunctor $\text{sSet} \times cM \xrightarrow{\Delta} M$ is close to being left Quillen

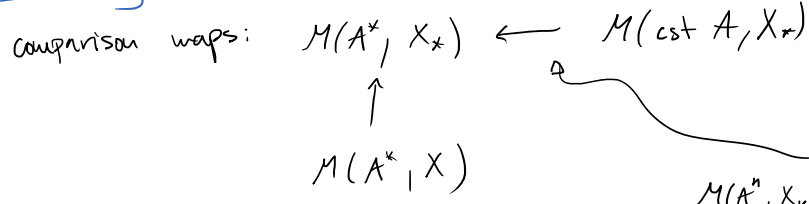
cofibrations on the left: $i^n : \Delta \hookrightarrow \mathcal{C}_R \in \mathcal{C}$
 $i^n : \Delta \hookrightarrow Wn \mathcal{C} \in Wn \mathcal{C}$

trivial cofibrations on the left more subtle: $(\Delta^{n-1} \xrightarrow{\sim} \Delta^n) : \Delta (X \rightarrow Y) : \begin{matrix} X_{n-1} \rightarrow Y_{n-1} \\ \downarrow \quad \downarrow \\ X_n \rightarrow Y_n \end{matrix}$
 these "generate" analoge extensions in some sense

$\Delta^+ \Delta^0 = \text{co} \Delta^1 \xrightarrow{\sim} \Delta^1 \xrightarrow{\sim} \Delta^0$ cylinder on Δ^0

$A + A = \text{co} \Delta \cdot \Delta A^* \xrightarrow{\sim} \Delta \cdot \Delta A^* \xrightarrow{\sim} A$ cylinder on A

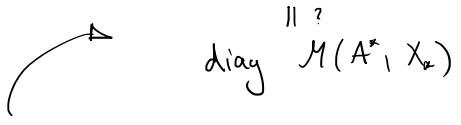
Balancing $M(A^*, X_*) \in \text{s}^2\text{Set} = [\Delta^{op} \times \Delta^{op}, \text{Set}]$



ptwise w.e. between Reedy cof. obj's
 $M(A^*, X_*) \xleftarrow{\sim} M(\text{cst } A, X_*)$ since $A^n \xrightarrow{\sim} \text{cst } A$

Take the geometric realizations:

$$\Delta \cdot \Delta M(A^*, X_*) \xleftarrow{\sim} \Delta \cdot \Delta M(\text{cst } A, X_*) = M(A, X_*)$$



then this + the dual statement proves the balancing

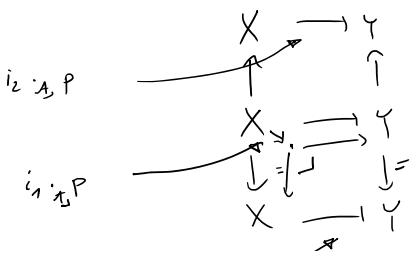
this is about bisimplicial sets $K_{..} = (K.)$.

$$\begin{aligned} \Delta \cdot \Delta K. &= \int^{n \in \Delta} \Delta^n \times_{\text{sSet}} K_n = \int^{n \in \Delta} \Delta^n \times \left(\int^{k \in \Delta} \Delta^k \cdot K_{nk} \right) \\ &= \int^{n, k} (\Delta^n \times \Delta^k) \cdot K_{nk} = \int^{(n, k)} \Delta \times \Delta (\text{diag } -) (n, k) \cdot K_{nk} \\ &\stackrel{\text{Yoneda}}{=} \text{diag } K \end{aligned}$$

Application $A = (0 \xleftarrow{-} 1 \xrightarrow{+} 2)$.

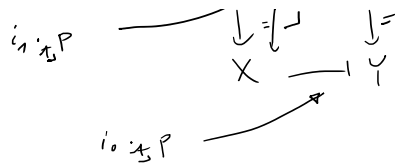
$[A, M] \xrightarrow[\text{cst}]{\text{colim}} M$ cst pres w.e. (always) and fibrations: $p : X \twoheadrightarrow Y$

\Rightarrow colim takes w.e.'s between Reedy cofibrant diagrams to w.e.'s



(would not work for

(would not work for
 $A = (0 \leftarrow 1 \rightarrow 0)$)



Application.

M simplicial

$[A^{\text{op}}, \text{sSet}] \times [A, M] \xrightarrow{-\otimes_A^-} M$ is left Quillen

proj. ptwise
 ptwise proj.

$Q_* = N(-/A)$ a particular cofibrant replacement of $*$

$\Rightarrow Q_* \otimes_A^- : [A, M_c] \rightarrow M_c$ preserves w.e.'s

\downarrow is a w.e. on proj. cof. diagrams
 $* \otimes_A^- \Rightarrow N(-/A) \otimes_A^- = \text{hocolim}_A$
 "colim"
 \downarrow

This can be extended to all model categories via frames:

$N(-/A) : \Delta^{\text{op}} \times A^{\text{op}} \rightarrow \text{Set}$

$D^* : \Delta \times A \rightarrow M_c$

frame on $D : A \rightarrow M_c$

$\rightsquigarrow N(-/A) \cdot_{\Delta \times A} D^* =: \text{hocolim}_A D$

"
 $N(-/A) \otimes_A D$ if we denote $K \otimes X = K \cdot_{\Delta} X^*$ the "action" of sSet on M

Bousfield localizations

$$M \rightsquigarrow \text{Ho}(M) \sim M_{\text{cf}} / \text{htpy}$$

What should a subcategory be?

We want to limit fibrant/cofibrant objects.

(homotopy) injectivity projectivity
 $\hat{=}$ reflective subcat's

Example.

• $Ab \in Gp$

$$\begin{array}{ccc} \mathbb{Z} * \mathbb{Z} & \longrightarrow & A \\ \downarrow & \nearrow & \\ \mathbb{Z} \oplus \mathbb{Z} & & \end{array}$$

• constant diagrams $D: \Delta \rightarrow M$

$M = \text{Set}$: $\Delta_m \rightarrow D \quad D_m$
 $\downarrow \nearrow \exists!$ $\hat{=}$ $\uparrow \hat{=}$
 $\Delta_n \quad D_n$

M general: $\Delta_m \cdot X \rightarrow D \quad M(X, D_m)$
 $\downarrow \nearrow \exists!$ $\hat{=}$ $\uparrow \hat{=}$
 $\Delta_n \cdot X \quad M(X, D_n)$

• diagrams $D: \mathcal{U} \rightarrow M$ st. $D_n \hat{=} D_1 \times \dots \times D_1$

$M = \text{Set}$: $N_1 + \dots + N_1 \rightarrow D \quad D_1 \times \dots \times D_1$
 $\downarrow \nearrow \exists!$ $\hat{=}$ $\uparrow \hat{=}$
 $N_n \quad D_n$

• diagrams $D: \mathcal{U} \rightarrow M$ st. $D_n \hat{=} \{K, D_{n+1}\}$

$N_{n+1} \otimes K \rightarrow D \quad \{K, D_{n+1}\}$
 $\downarrow \nearrow \exists!$ $\hat{=}$ $\uparrow \hat{=}$
 $N_n \quad D_n$

$\mathcal{V} = \text{Top}_* / \text{Set}_*$
 $K = S^1$
 $\rightsquigarrow D_n \rightarrow \Omega D_{n+1}$

• sheaves $\text{colim } \mathcal{F}_{U_i} \rightarrow D \quad \text{lim } D_{U_i}$
 $\downarrow \nearrow \exists!$ $\hat{=}$ $\uparrow \hat{=}$
 $\mathcal{F}_U \quad D_U$

$\mathcal{F} = \text{Op}(X)$
 U_i open covering of U
 closed under intersections

Example. where htpy version is necessary

$M = \text{Top}$: $S^n \rightarrow X \quad \pi_n X = 0$
 $\downarrow \nearrow \exists$ unique up to higher htpys
 $D_{n+1} \quad ?$

$\text{map}(S^1, X) \quad \pi_n X = 0$
 $\uparrow \cong$ $\pi_{n+1} X = 0$
 $\text{map}(D^{n+1}, X) \quad \vdots$

more generally $A \rightarrow X \quad \text{map}(A, X)_n$

more generally

$$\begin{array}{ccc} A & \longrightarrow & X \\ \downarrow & \nearrow \exists! h & \\ B & & \end{array} \cong \begin{array}{ccc} & & \text{map}(A, X) \\ & \uparrow \sim & \\ & & \text{map}(B, X) \end{array}$$

better:

$$\begin{array}{ccc} A & \longrightarrow & X \\ \downarrow & \nearrow \exists! h & \\ B & & \end{array} \cong \begin{array}{ccc} \text{map}(A, X) & \longleftarrow & L \\ \uparrow \sim & \nearrow & \uparrow \\ \text{map}(B, X) & \longleftarrow & K \end{array} \cong \begin{array}{ccc} K \times B +_{K \times A} L \times A & \longrightarrow & X \\ \downarrow & \nearrow \exists & \\ L \times B & & \end{array}$$

Will assume \mathcal{M} left proper, cellular = cofibrantly generated with cofibrations effective mono's

Definition. Let $f: A \rightarrow B$ be a cofibration between cofibrant objects. We say that W is **f -local** if it is fibrant and

$$f^*: \text{map}(B, W) \xrightarrow{\sim} \text{map}(A, W) \quad \left. \vphantom{f^*} \right\} (\mathcal{Y} \cup \{L^n B +_{L^n A} A^n \rightarrow B^n\})^{\square}$$

is a weak equivalence of simplicial sets.

Definition. A map $g: X \rightarrow Y$ is an **f -local equivalence** if its cofibrant replacement $\tilde{g}: \tilde{X} \rightarrow \tilde{Y}$

$$\begin{array}{ccc} 0 \rightarrow \tilde{X} & \xrightarrow{\sim} & X \\ \parallel & \downarrow \tilde{g} & \downarrow g \\ 0 \rightarrow \tilde{Y} & \xrightarrow{\sim} & Y \end{array}$$

gives, for each f -local W , a w.e. $\tilde{g}^*: \text{map}(\tilde{Y}, W) \xrightarrow{\sim} \text{map}(\tilde{X}, W)$

(any two related by a zig-zag of w.e.'s of such \Rightarrow independent of choice)

Definition. An **f -localization** of X is an f -local equivalence $j: X \rightarrow \hat{X}$ with \hat{X} f -local.

Aim. Construct an " f -local" model structure in which:

- f -local = fibrant; fibrations are complicated BUT cofibrations of \mathcal{M}
- f -local equivalence = weak equivalence
- f -localization = fibrant replacement = "reflection"

\rightarrow better: $\text{Id}: \mathcal{M} \rightleftarrows \mathcal{M}^{f\text{-local}}: \text{Id}$

preserves cof. & w.e.

is interpreted as $\mathcal{M}_{\text{cf}} \xrightleftharpoons[f\text{-localization}]{\text{Id}} \mathcal{M}_{\text{cf}}^{f\text{-local}}$