

Reminder on weighted colimits

$$\begin{array}{c}
 X \times_{\mathcal{G}} Y \\
 (xg, yg) = (x, yg) \\
 X \times_{\mathcal{G}} Y \\
 \downarrow \text{act}_x \times 1 \quad \downarrow 1 \times \text{act}_y \\
 X \times Y
 \end{array}$$

$$\begin{array}{c}
 M \otimes_{\mathcal{R}} N \\
 x \otimes y = x \otimes y \\
 M \otimes_{\mathcal{R}} N \\
 \downarrow \downarrow \\
 M \otimes N \\
 M: \mathcal{R}^{\text{op}} \rightarrow \text{Ab} \\
 N: \mathcal{R} \rightarrow \text{Ab}
 \end{array}$$

$$\begin{array}{c}
 W \otimes_A D \\
 \sum_{b,c} W_c \otimes A(b,c) \otimes D_b \\
 \downarrow \downarrow \\
 \sum_a W_a \otimes D_a \\
 W: A^{\text{op}} \rightarrow \mathcal{V} \\
 D: A \rightarrow \mathcal{M}
 \end{array}$$

$$\begin{array}{c}
 \Delta^* \times_{\Delta} X \\
 (\theta^* t, x) \sim (t, \theta_* x) \\
 \Delta^* \times_{\Delta} X \\
 \downarrow \downarrow \\
 \sum \Delta^i \times X_i \\
 \Delta: \Delta \rightarrow \text{sSet} \\
 X: \Delta^{\text{op}} \rightarrow \text{Top} \dots
 \end{array}$$

$$W \otimes_A D = \int^{A \in \mathcal{A}} W_A \otimes D_A$$

$$\begin{array}{l}
 \otimes: \mathcal{M} \times \mathcal{N} \rightarrow \mathcal{P} \\
 \otimes_A: [\mathcal{A}^{\text{op}}, \mathcal{M}] \times [\mathcal{A}, \mathcal{N}] \rightarrow \mathcal{P}
 \end{array}$$

$$\sum_{b,c} W_c \otimes A(b,c) \otimes D_b \implies \sum_a W_a \otimes D_a \xrightarrow{\text{coend}} W \otimes_A D$$

$$/ \mathcal{M}(-, N)$$

$$\prod_{b,c} \{A(b,c), \mathcal{M}(W_c \otimes D_b, N)\} \leftarrow \prod_a \mathcal{M}(W_a \otimes D_a, N) \leftarrow \mathcal{M}(W \otimes_A D, N)$$

$$\prod_{b,c} \{A(b,c), \mathcal{V}(W_c, \mathcal{M}(D_b, N))\} \leftarrow \prod_a \mathcal{V}(W_a, \mathcal{M}(D_a, N)) \leftarrow [\mathcal{A}^{\text{op}}, \mathcal{V}](W, \mathcal{M}(D, N))$$

$$\mathcal{M}(D_A, N) \stackrel{\text{Yoneda}}{=} [\mathcal{A}^{\text{op}}, \mathcal{V}](A(-, A), \mathcal{M}(D, N)) \stackrel{\text{def}}{=} \mathcal{M}(A(-, A) \otimes_A D, N)$$

Example

$$A(-, A) \otimes_A D = DA$$

dually

$$\{A(A, -), D\}_A = DA$$

$$A(-, -) \otimes_A D = D$$

$$\mathcal{R} \otimes_{\mathcal{R}} M = M$$

$$\text{Hom}_{\mathcal{R}}(\mathcal{R}, M) = M$$

$$\left. \begin{array}{l}
 \mathcal{A}^{\text{op}} \rightarrow [\mathcal{B}, \mathcal{V}] \\
 W: \mathcal{A}^{\text{op}} \times \mathcal{B} \rightarrow \mathcal{V} \\
 D: \mathcal{A} \rightarrow \mathcal{M}
 \end{array} \right\} W \otimes_A D: \mathcal{B} \rightarrow \mathcal{M}$$

$$\subseteq \mathcal{M}_{\mathcal{R}} \otimes_{\mathcal{R}} \mathcal{R}^{\mathcal{X}}$$

Restriction of scalars: $F: \mathcal{B} \rightarrow \mathcal{A}$

$$A(-, F-) : \mathcal{A}^{\text{op}} \times \mathcal{B} \rightarrow \mathcal{V}$$

$$D: \mathcal{A} \rightarrow \mathcal{M}$$

$$A(-, F-) \otimes_A D = DF$$

dually

$$\{A(F-, -), D\}_A = DF$$

$$\mathcal{R} \rightarrow \mathcal{S}$$

$$\begin{array}{l}
 \mathcal{R} \otimes_{\mathcal{R}} \mathcal{V} = \mathcal{V}^{\mathcal{R}} \\
 \mathcal{S} \otimes_{\mathcal{S}} \mathcal{M} = \mathcal{M}
 \end{array}$$

$$F^* D$$

Adjunction:

$$[B, \mathcal{V}](C, \{A(F-, -), D\}_A) = [A, \mathcal{V}](A(F-, -) \otimes_{\mathcal{B}} C, D)$$

$$\text{Hom}_{\mathcal{R}}(M, \text{Hom}_{\mathcal{S}}(S, N)) \underset{\mathcal{R}}{=} \text{Hom}_{\mathcal{S}}(S \otimes_{\mathcal{R}} M, N) \quad \text{F.C. extension of scalars}$$

(like $S \otimes_{\mathcal{R}} L$)

$$[B, \mathcal{V}](C, F^* D) = [A, \mathcal{V}](F_* C, D)$$

$$\leftarrow \text{Lan}_F C$$

left Kan extension

eg. $F: \mathcal{B} \hookrightarrow \mathcal{A}$ full embedding

$$\begin{array}{ccc}
 B & \xrightarrow{c} & M \\
 \text{fully} & \downarrow F & \parallel \\
 \text{faithful} & A & \xrightarrow{\text{Lan}_F C} M
 \end{array}$$

$$\begin{aligned}
 F_! C(B) &= A(F-, B) \otimes_B C \\
 &= B(-, B) \otimes_B C \\
 &= C_B \quad \text{real extension}
 \end{aligned}$$

$$B \xrightarrow{F} A \xrightarrow{D} M$$

$$A \xrightarrow{D} M \xrightarrow{1} M$$

$$W \otimes_A D = \dots \otimes_A 1$$

$$W \otimes_B (D \circ F) = \bar{W} \otimes_A D$$

$$\begin{aligned}
 W \otimes_B (A(-, F-) \otimes_A D) &= (W \otimes_B A(-, F-)) \otimes_A D \\
 \bar{W} &= \text{Lan}_F W
 \end{aligned}$$

$$\begin{aligned}
 M \otimes_R N &= M \otimes_R (S \otimes_S N) \\
 &= (M \otimes_R S) \otimes_S N
 \end{aligned}$$

restricted from S to R

weighted colimit

$$\text{Hom}(W \otimes_A D, X) \cong \text{Hom}_A(W, \text{Hom}(D, X))$$

$$\begin{array}{c}
 M \quad A(-, A) \\
 \uparrow \\
 [A^{\text{op}}, V]
 \end{array}$$

$$\bullet \quad A^A \otimes_A D \cong D_A \quad \text{Yoneda}$$

$$(R \otimes_R M \cong M)$$

$$\bullet \quad (\text{colim } W_i) \otimes_A D \cong \text{colim } (W_i \otimes_A D)$$

$$\bullet \quad K \otimes_* X = K \otimes X$$

Reedy model categories, framings

$$M(M, N) \quad \downarrow \cong$$

Motivation. Any model category M is in some weak sense enriched over $sSet$ and, as a result, $Ho(M)$ will be enriched over $Ho(sSet)$.

Start with an honest $sSet$ -model category = simplicial model category

— since $K \in sSet$ is a colimit $K = \text{colim}_{(n, \Delta^n \rightarrow K)} \Delta^n = K \cdot_{\Delta} \Delta^*$

$$\begin{aligned} \Delta^* &: \Delta \rightarrow sSet \\ \leftarrow &: \Delta^{op} \rightarrow sSet \end{aligned}$$

we have $K \otimes M = \text{colim} (\Delta^n \otimes M)$ and it is enough to give $\Delta^* \otimes M$

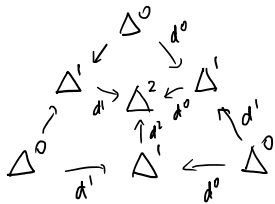
— clearly $\Delta^0 \otimes M \cong M$, since Δ^0 is the monoidal unit

— what is the essential property of $\Delta^1 \otimes M$?

it is a cylinder! ... at least for $M \in M_c$:

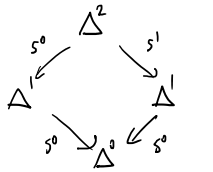
$$\Delta^0 + \Delta^0 \xrightarrow{[d^0, d^1]} \Delta^1 \xrightarrow{s^0} \Delta^0 \quad \rightsquigarrow \quad M + M \xrightarrow{\quad} \Delta^1 \otimes M \xrightarrow{\sim} M$$

— what about $\Delta^2 \otimes M$?



, more compactly $\partial \Delta^2 \rightarrow \Delta^2$
 \parallel
 $L_2 \Delta^*$

$$\rightsquigarrow L_2 \Delta^* \otimes M \rightarrow \Delta^2 \otimes M \xrightarrow{\sim} M \otimes M$$



, more compactly $\Delta^2 \xrightarrow{\sim} M_2 \Delta^*$

$\bigvee_0 \Delta^i: \Delta \rightarrow sSet$
 a cosimplicial object in $sSet$

— need some calculus of such diagrams $\Delta \rightarrow sSet \rightsquigarrow \Delta^* \otimes M: \Delta \rightarrow M$
 \rightarrow Reedy categories, Reedy model structures $\Delta \rightarrow M$ a cosimplicial object in M

A Reedy category has two kinds of maps — direct and inverse (either d^i and s^i in Δ)

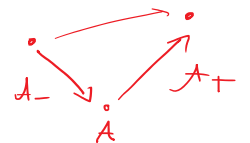
Definition. A **direct category** is a category \mathcal{A} together with a functor $\text{deg}: \mathcal{A} \rightarrow \lambda$ that satisfies $f: A \rightarrow B \rightsquigarrow f = 1 \Leftrightarrow \text{deg } A = \text{deg } B$
 \uparrow ordinal $\Rightarrow \text{deg } A < \text{deg } B$ $f \neq 1 \Rightarrow \text{deg } A < \text{deg } B$

An **inverse category** is a dual notion, i.e. a category \mathcal{A} together with a functor $\text{deg}: \mathcal{A}^{op} \rightarrow \lambda$ satisfying $f = -1 \Leftrightarrow \text{deg } A = \text{deg } B$.

A **Reedy category** is a category \mathcal{A} together with two subcategories $\mathcal{A}^+, \mathcal{A}_-$ and a function $\text{deg}: \text{ob } \mathcal{A} \rightarrow \lambda$ that

- makes \mathcal{A}^+ into a direct category
- makes \mathcal{A}_- into an inverse category
- any morphism has a unique decomposition:

$$\sum_{\mathcal{A} \neq \pm} \mathcal{A}^+(A, C) \times \mathcal{A}_-(B, A) \xrightarrow{\cong} \mathcal{A}(B, C)$$



any morphism ...

$$\sum_{A \in \mathcal{A}} A^+(A, C) \times A_-(B, A) \xrightarrow{\cong} A(B, C)$$

$\sigma = \downarrow \circ \uparrow$
A

as a composition of an inverse and a direct morphism

The Yoneda lemma gives $X \cong A(-, -) \cdot X$ and we will describe a way of building X by decomposing $A(-, -)$. The axioms actually give

$$A(-, -) = \sum_{A \in \mathcal{A}} A^+(A, -) \times A_-(-, A) = \sum_{A \in \mathcal{A}} A_A^+ \times A_A^- \quad (\text{but only of functors } A^{op} \times A^+ \rightarrow \text{Set } \nabla)$$

but in order to understand $A(-, -)$ we factor it into $st_n A(-, -)$... maps that factor through some A of degree $\leq n$ and clearly we have

$$A(-, -) = \text{colim}_{n < \infty} st_n A(-, -)$$

so that it remains to study the difference between $st_n A$ and $st_{n-1} A$ or, better for n limit, $st_{\infty} A = \text{colim}_{i < \infty} st_i A$.

Quite clearly, we have a pushout square

$$\begin{array}{ccc} 0 & \longrightarrow & st_{\infty} A \\ \downarrow & & \downarrow \\ \sum_{A \in \mathcal{A}_n} A_A^+ \times A_A^- & \longrightarrow & st_n A \end{array}$$

with A ranging over all objects of degree n

However, it will be crucial to express this in terms of representable functors on \mathcal{A} , rather than on A^+, A_- .

Examples.

- any ordinal is a direct category
 - $\begin{array}{c} \rightarrow \\ \leftarrow \\ \rightarrow \end{array}$ is a direct category
 - $\begin{array}{c} \leftarrow \\ \rightarrow \end{array}$ is a Reedy category
 - Δ^+ is a direct category (non-empty finite ordinals + monos)
 - Δ_- is an inverse category (non-empty finite ordinals + epis)
 - Δ is a Reedy category $\Rightarrow \Delta^{op}$ Reedy
- A direct $\Leftrightarrow A^{op}$ inverse
 A Reedy $\Leftrightarrow A^{op}$ Reedy

Notation. We denote

- $A_A = A(A, -) \in [A, \text{Set}]$ the covariant representable
- $A^A = A(-, A) \in [A^{op}, \text{Set}]$ the contravariant representable ... think Δ^n

We further define two subfunctors

- $i_A: \circ A_A \subseteq A_A$ of maps that factor through an object of lower degree.
- $i^A: \circ A^A \subseteq A^A$ of maps that factor through an object of lower degree.

In the decomposition

$$\begin{aligned} A_A &= A(A, -) = \sum A^+(B, -) \times A_-(A, B) \\ \circ A_A &= \sum A^+(B, -) \times \circ A_-(A, B) \end{aligned}$$

↑ only 1 is excluded

This means that there is a pushout (coproduct)

$$\begin{array}{ccc} 0 & \longrightarrow & \partial A_A \\ \downarrow & & \downarrow i_A \\ A_A^+ & \longrightarrow & A_A \end{array}$$

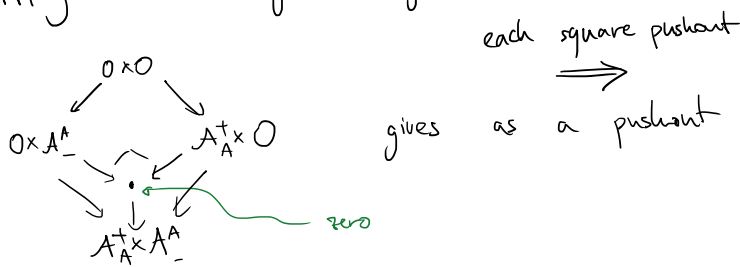
Dually, we get

$$\begin{aligned} A^+ &= A(-, A) = \sum A^+(B, A) \times A_-(-, B) \\ \partial A^+ &= \sum \partial A^+(B, A) \times A_-(-, B) \end{aligned}$$

and a pushout (coproduct)

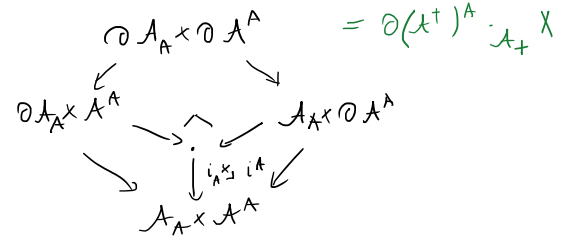
$$\begin{array}{ccc} 0 & \longrightarrow & \partial A^+ \\ \downarrow & & \downarrow i^+ \\ A_-^+ & \longrightarrow & A^+ \end{array}$$

Putting these together yields that



$\partial A^+(B, A)$ by Yoneda

$$\begin{aligned} &= \sum \partial(A^+)^A \times_{A^+} A^+(B, -) \times A_-(-, B) \\ &= \partial(A^+)^A \times_{A^+} \sum A^+(B, -) \times A_-(-, B) \\ &= \partial(A^+)^A \cdot i^+ A(-, -) \quad A(-, -) \\ &= \text{extension of } \partial(A^+)^A \\ &\implies \partial A^+ \cdot X = (\partial(A^+)^A \times_{A^+} A(-, -)) \cdot X \end{aligned}$$



Now we get a diagram

$$\begin{array}{ccc} 0 & \longrightarrow & \sum_A \partial A_A \times A^+ + \sum_A A_A \times \partial A^+ \\ \downarrow & & \downarrow \sum_A i_A \times i^+ \\ \sum_A A_A^+ \times A^+ & \longrightarrow & \sum_A A_A \times A^+ \end{array} \longrightarrow \text{sk}_n A(-, -)$$

sums range over $A \in A$ with $\text{deg } A = n$

in which the outer square is a pushout \implies so is the one on the right but now in $[A^{\text{op}}, \text{Set}]^{\triangleright}$

Upon applying $- \cdot X$, we denote:

$$\begin{array}{ccc} \partial A^+ \cdot X & \xrightarrow{i^+ \cdot X} & A^+ \cdot X \\ \parallel \text{def} & & \parallel \\ L_A X & \xrightarrow{\partial A X} & X_A \end{array}$$

Example. $A = \Delta^{\text{op}}, M = \text{Set}$
 $\implies L_n X \in X_n$ the subset of deg. simplices (needs a bit of work, see above).
 $0 = \text{sk}_{-1} X$
 $X = \text{colim sk}_n X$ and

Theorem. For any $X \in [A, M]$ we get

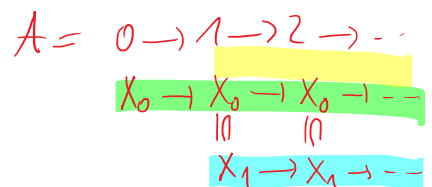
$$\begin{array}{ccc} \sum_A \partial A_A \cdot X_A + \sum_A A_A \cdot L_A X & \longrightarrow & \text{sk}_n X \\ \downarrow i_A \cdot \partial A X & & \downarrow \\ \sum_A A_A \cdot X_A & \longrightarrow & \text{sk}_n X \end{array}$$

Important special case.

When A is direct, we have $\partial A_A = 0$ and, consequently,

$$\sum_A A_A \cdot L_A X \longrightarrow \text{sk}_n X$$

proj. model structure generated by



$$\sum_A A_A \cdot L_A X \longrightarrow \text{str}_n X$$

$$\downarrow \quad \quad \quad \downarrow$$

$$\sum_A A_A \cdot X_A \longrightarrow \text{str}_n X$$

proj. model structure generated by $A_A \cdot K \rightarrow A_A \cdot L$
 $A \in \mathcal{A}, K \rightarrow L$

so that: $\forall A \in \mathcal{A} : \partial_A X : L_A X \rightarrow X_A$ cofibration

next theorem

(\Rightarrow) X is cofibrant in the projective model structure

Eg. $X = X_0 \rightarrow X_1 \rightarrow X_2 \rightarrow \dots$

$$L_0 X \rightarrow X_0 : 0 \rightarrow X_0$$

$$L_n X \rightarrow X_n : X_{n-1} \rightarrow X_n$$

More generally, for $f: X \rightarrow Y$ we denote pushout corner map in

$$\begin{array}{ccc} L_A X & \longrightarrow & L_A Y \\ \partial_A X \downarrow & \searrow & \downarrow \partial_A Y \\ X_A & \longrightarrow & Y_A \end{array} \quad \partial_A f = \begin{array}{ccc} \partial_A^A & & X \\ \downarrow i^A & i_A & \downarrow f \\ A^A & & Y \end{array}$$

$$X = \text{str}_n^X Y$$

$$Y = \text{colim}_{n \in \mathbb{N}} \text{str}_n^X Y \text{ and}$$

Theorem. For any map $f: X \rightarrow Y$ of $[A, M]$ we get

$$\begin{array}{ccc} \longrightarrow & \text{str}_n^X Y & \\ \downarrow \Sigma i_A \cdot i_A \cdot \partial_A f & \searrow & \downarrow \\ \longrightarrow & \text{str}_n^X Y & \end{array}$$

$$\begin{array}{ccc} X & & X \\ \downarrow f & = & \downarrow f \\ Y & & Y \end{array} \quad \text{str}_n^X Y = \text{str}_n Y + \text{str}_n X$$

Theorem. For a map $f: X \rightarrow Y$ of $[A, M]$ TFAE

- $\forall A : \partial_A f$ is a (trivial) cofibration
- $f \in \text{cell} \{ i_A \cdot i : A \in \mathcal{A}, i \in M \}$ a (trivial) cofibration

$$\begin{array}{ccc} \partial_A \cdot L + \partial_A \cdot K & & A_A \cdot K \\ \downarrow & & \downarrow \\ A_A \cdot L & & \end{array}$$

These maps are called **Reedy (trivial) cofibrations**.

Proof. The implication " \Downarrow " is the previous theorem.

For the implication " \Uparrow " observe that ∂_A commutes with all cellular constructions so that it is enough to study $\partial_A (i_B \cdot i)$:

$$\partial_A^A \cdot i_A \left(\begin{array}{ccc} \partial_A B & & K \\ \downarrow i_B & \cdot & \downarrow i \\ A_B & & L \end{array} \right) = \left(\begin{array}{ccc} \partial_A^A & & \partial_A B \\ \downarrow i^A & i_A & \downarrow i_B \\ A^A & & A_B \end{array} \right) \cdot \downarrow i$$

either iso for $A \neq B$

$$A_{\neq 1}(B, A) \hookrightarrow A(B, A) \text{ for } A = B$$

pushout of $0 \rightarrow 1$

pushout of i

□

Dually, we denote

$$M_A X = \{ \partial_A \cdot i_A \cdot X \}_A$$

$$\delta_A X \uparrow \quad \quad \uparrow i_A \cdot X \cdot \beta_A$$

$$X_A = \{ A_A \cdot X \}_A$$

and more generally for $f: X \rightarrow Y$:

$S_{A,f} = \{i_A, f\}_{A,r}$ = pullback corner map in

$$\begin{array}{ccc} X_A & \longrightarrow & M_A X \\ \downarrow & & \downarrow \\ Y_A & \longrightarrow & M_A Y \end{array}$$

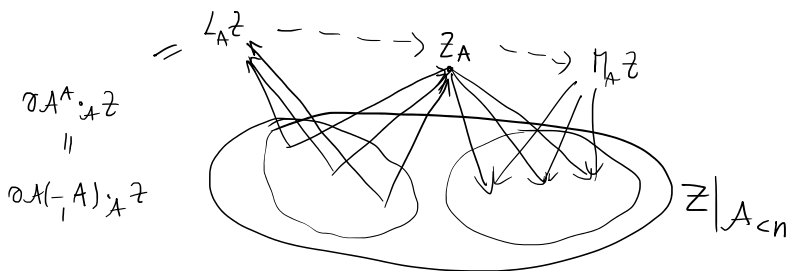
Theorem TFAE

- $\forall A \in \mathcal{A}$: $S_{A,f}$ is a (trivial) fibration
 - $f \in \text{cocell} \{ \{i_A, p\}_r \mid A \in \mathcal{A}, p \in \mathcal{M} \text{ a (trivial) fibration} \}$
- \rightsquigarrow **Reedy (trivial) fibrations.**

Theorem. There is a model structure on $[\mathcal{A}, \mathcal{M}]$, called the **Reedy model structure** with $\mathcal{C} =$ Reedy cofibrations, $\mathcal{F} =$ Reedy fibrations, $\mathcal{W} =$ pointwise weak equivalences.

Proof. We need to show that $\mathcal{W} \cap \mathcal{C} =$ exactly the **Reedy** inductively fibrations. Easily $i_A \circ j \in \mathcal{C} \Leftrightarrow j \in \underbrace{\{i_A, f\}_r}_{\text{fibrations by definition}}$.

The factorizations are produced inductively using the following idea:



to give an extension of z from $\mathcal{A}_{<n}$ to $\mathcal{A}_{\leq n}$ we need $\forall A$ of degree n to factor

$$\begin{array}{ccccc} L_A z & \longrightarrow & Z_A & \longrightarrow & M_A z \\ & \searrow & & \nearrow & \\ & & & & \text{obtained from } z|_{\mathcal{A}_{<n}} \end{array}$$

$$\begin{array}{ccccc} \text{Now} & L_A X & \longrightarrow & L_A z & \longrightarrow & L_A Y \\ & \downarrow & & \swarrow & & \downarrow \\ & X_A & \longrightarrow & Z_A & \longrightarrow & Y_A \\ & \downarrow & & \swarrow & & \downarrow \\ & M_A X & \longrightarrow & M_A z & \longrightarrow & M_A Y \end{array}$$

Application. • Properness in $\mathcal{M}_c / \mathcal{M}_f$ differently.

• Given $A \in \mathcal{M}_c$ consider $\text{cst } A \in [\Delta, \mathcal{M}]$ and consider its Reedy cofibrant replacement $A^* \in [\Delta, \mathcal{M}]$. \mathcal{M} simplicial $A^* = \Delta^* \circ A$ does the job $X_* = \{ \Delta^*, X \} = X \Delta^*$ cofibrant when A is (= $\Delta_0 \cdot A$)

$O \mapsto A^* \xrightarrow{\sim} \text{cst } A$

– we may achieve that $A^0 = A$ and then A^* is called a **cosimplicial frame** on A . Any two are related by a zig-zag

of cosimplicial frames. Thus for $X \in M_f$

$$M(A^*, X) \in sSet$$

Dually $M(A, X_*) \in sSet$

$$A^*: \Delta \rightarrow M \in M \quad \} \quad M(A^*, X): \Delta^{op} \rightarrow sSet$$

via a simplicial frame on X . $M(A^*, X)_0 = M(A, X)$
 $M(A^*, X)_1 = M(Cyl A, X)$

$$A^*: \Delta \rightarrow M$$

$$K: \Delta^{op} \rightarrow sSet$$

get $K \cdot_{\Delta} A^* \in M$, an action of $sSet$ on M_c
 $K * A \stackrel{def}{=} K \cdot_{\Delta} A^*$, but not associative

The bifunctor $sSet \times cM \xrightarrow{- \cdot_{\Delta} -} M$ is close to being left Quillen

cofibrations on the left:

$$i^n \cdot_{\Delta} C_R \in C$$

$$i^n \cdot_{\Delta} W \cap C \subseteq W \cap C$$

action of $Ho(sSet)$ on $Ho(M_c)$ is assoc.
 Reedy cof. between frames

trivial cofibrations on the left more subtle:

$$(\Delta^{n-1} \hookrightarrow \Delta^n) \cdot_{\Delta} (X \rightarrow Y): X_{n-1} \rightarrow Y_{n-1}$$

$$\downarrow \quad \downarrow$$

$$X_n \rightarrow Y_n$$

these "generate" analogue extensions in some sense

$$\Delta^+ + \Delta^0 = \circ \Delta^1 \xrightarrow{\sim} \Delta^1 \xrightarrow{\sim} \Delta^0 \quad \text{cylinder on } \Delta^0$$

$$A + A = \circ \underbrace{\Delta^1 \cdot_{\Delta} A^*}_{L \cdot A} \xrightarrow{\sim} \underbrace{\Delta^1 \cdot_{\Delta} A^*}_{A^*} \xrightarrow{\sim} A \quad \text{cylinder on } A$$

Balancing. $M(A^*, X_*) \in s^2Set = [\Delta^{op} \times \Delta^{op}, sSet]$

$$\text{comparison maps: } M(A^*, X_*) \leftarrow M(\text{cst } A, X_*)$$

$$\uparrow$$

$$M(A^*, X)$$

ptwise w.e. between Reedy cof. obj's

$$M(A^n, X_*) \leftarrow M(\text{cst } A, X_*) \quad \text{since } A^n \xrightarrow{\sim} \text{cst } A$$

Take the geometric realizations:

$$\Delta \cdot_{\Delta} M(A^*, X_*) \xleftarrow{\sim} \Delta \cdot_{\Delta} M(\text{cst } A, X_*) = M(A, X_*)$$



$$\text{diag } M(A^*, X_*)$$

then this + the dual statement proves the balancing

this is about bisimplicial sets $K_{..} = (K.)$.

$$\Delta \cdot_{\Delta} K. = \int^{n \in \Delta} \Delta^n \times_{sSet} K_n = \int^{n \in \Delta} \Delta^n \times \left(\int^{k \in \Delta} \Delta^k \cdot K_{nk} \right)$$

$$= \int^{n, k} (\Delta^n \times \Delta^k) \cdot K_{nk} = \int^{(n, k)} \Delta \times \Delta (\text{diag } -, (n, k)) \cdot K_{nk}$$

$$\stackrel{\text{Yoneda}}{=} \text{diag } K$$

Application. $A = (0 \leftarrow 1 \xrightarrow{+} 2)$.

$$[A, M] \xrightleftharpoons[\text{cst}]{\text{colim}} M$$

cst pres w.e. (always)

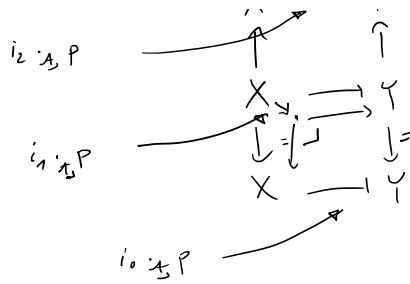
and fibrations: $p: X \twoheadrightarrow Y$

\Rightarrow colim takes w.e.'s between Reedy cofibrant to w.e.'s

$i_2 \cdot_{\Delta} p$

$$\begin{array}{ccc} X & \twoheadrightarrow & Y \\ \uparrow & & \uparrow \\ X_{..} & \twoheadrightarrow & Y \end{array}$$

between Reedy cofibrant diagrams to w.e.l.s
 (would not work for $A = (0 \leftarrow 1 \rightarrow 0)$)



Application.

M simplicial
 $[A^{\text{op}}, \text{sSet}] \times [A, M] \xrightarrow{- \otimes_A -} M$ is left Quillen
 proj. ptwise ptwise
 $Q^* = N(-/A)$ a particular cofibrant replacement of $*$

$\Rightarrow Q^* \otimes_A - : [A, M_c] \rightarrow M_c$ preserves w.e.l.s
 \downarrow is a w.e. on proj. cof. diagrams
 $* \otimes_A -$
 colim_A
 $\Rightarrow N(-/A) \otimes_A - = \text{hocolim}_A$

This can be extended to all model categories via frames:

$N(-/A) : \Delta^{\text{op}} \times A^{\text{op}} \rightarrow \text{Set}$
 $D^* : \Delta \times A \rightarrow M_c$ frame on $D : A \rightarrow M_c$

$\rightsquigarrow N(-/A) \underset{\Delta \times A}{\otimes} D^* =: \text{hocolim}_A D$
 $N(-/A) \otimes_A D$ if we denote $K \otimes X = K \underset{\Delta}{\cdot} X^*$ the "action" of sSet on M

Bousfield localizations

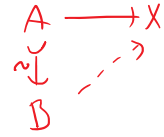
$$M \rightsquigarrow \text{Ho}(M) \sim \text{Mod}/\text{htpy}$$

What should a subcategory be?

We want to limit fibrant/cofibrant objects.

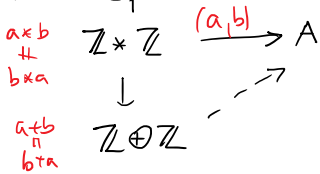
(homotopy) injectivity
 ≡ reflective subcat's

projectivity

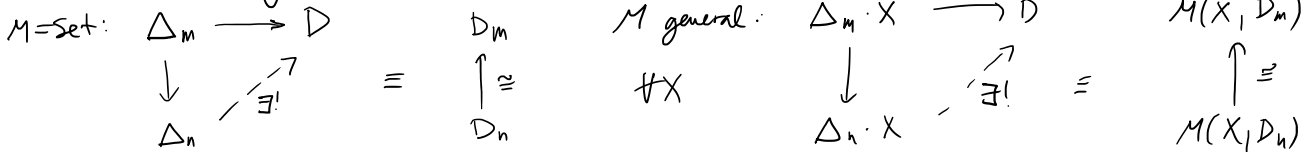


Example

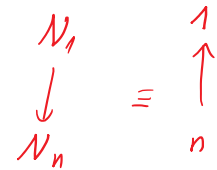
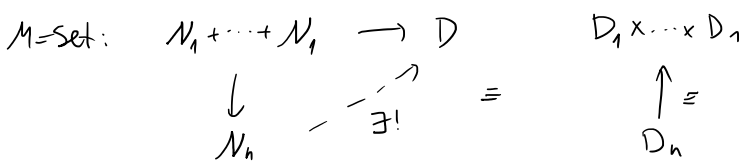
• $Ab \in \text{Gp}$



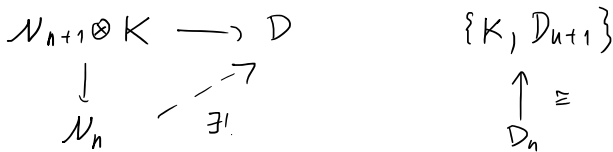
• constant diagrams $D: \Delta \rightarrow M$



• diagrams $D: \mathcal{U} \rightarrow M$ st. $D_n \cong D_1 \times \dots \times D_1$



• diagrams $D: \mathcal{U} \rightarrow M$ st. $D_n \cong \{K, D_{n+1}\}$

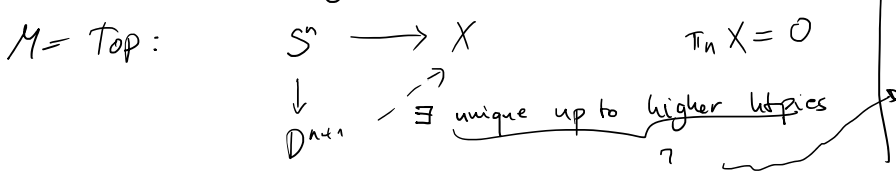


$\mathcal{V} = \text{Top}_* / \text{sSet}_*$
 $K = S^1$
 $\rightarrow D_n \rightarrow \Omega D_{n+1}$
 $\text{Hom}(N_{n+1} \otimes K, D) = \mathcal{V}(K, \text{Hom}(N_{n+1}, D))$
 $= \mathcal{V}(K, D_{n+1})$

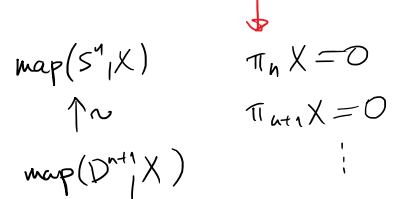
• sheaves $\text{colim } \mathcal{F}_{U_i} \rightarrow D$ $\lim D_{U_i}$
 \downarrow \nearrow $\exists!$ $\uparrow \cong$
 \mathcal{F}_U D_U

$\mathcal{F} = \text{Op}(X)$
 U_i open covering of U
 closed under intersections

Example where htpy version is necessary



X is a homotopy $(n-1)$ -type



more generally $A \rightarrow X$ $\text{map}(A, X)$
 \uparrow $\uparrow \dots$



more generally

$$\begin{array}{ccc}
 A & \longrightarrow & X \\
 \downarrow & \nearrow & \\
 B & \xrightarrow{\exists! h} &
 \end{array}
 \equiv
 \begin{array}{ccc}
 \text{map}(A, X) & & \\
 \uparrow \sim & & \\
 \text{map}(B, X) & &
 \end{array}$$



better:

$$\begin{array}{ccc}
 A & \longrightarrow & X \\
 \downarrow & \nearrow & \\
 B & \xrightarrow{\exists! h} &
 \end{array}
 \equiv
 \begin{array}{ccc}
 \text{map}(A, X) & \longleftarrow & L \\
 \uparrow \sim & \nearrow & \uparrow \\
 \text{map}(B, X) & \longleftarrow & K
 \end{array}
 \equiv
 \begin{array}{ccc}
 K \times B +_{K \times A} L \times A & \longrightarrow & X \\
 \downarrow & \nearrow & \\
 L \times B & \xrightarrow{\exists} &
 \end{array}$$

in Top $S^{n-1} \rightarrow D^n$

cofibration (between cofibrant)

Will assume \mathcal{M} left proper, cellular = cofibrantly generated with cofibrations effective mono's

Definition. Let $f: A \rightarrow B$ be a cofibration between cofibrant objects.

We say that W is f -local if it is fibrant and

$$f^*: \text{map}(B, W) \xrightarrow{\sim} \text{map}(A, W)$$

is a weak equivalence of simplicial sets.

$$\left. \begin{array}{l}
 \{ \text{Yon} \{ L^n B +_{L^n A} A^n \rightarrow B^n \} \}^{\square} \\
 n=0: A \rightarrow B
 \end{array} \right\}$$

Definition. A map $g: X \rightarrow Y$ is an f -local equivalence if its cofibrant replacement $\tilde{g}: \tilde{X} \rightarrow \tilde{Y}$

$$\begin{array}{ccc}
 0 \rightarrow \tilde{X} & \xrightarrow{\sim} & X \\
 \parallel & \downarrow \tilde{g} & \downarrow g \\
 0 \rightarrow \tilde{Y} & \xrightarrow{\sim} & Y
 \end{array}$$

$\Rightarrow f$ is an f -local equiv.

gives, for each f -local W , a w.e. $\tilde{g}^*: \text{map}(\tilde{Y}, W) \xrightarrow{\sim} \text{map}(\tilde{X}, W)$

(any two related by a zig-zag of w.e.'s of such \Rightarrow independent of choice)

Definition. An f -localization of X is an f -local equivalence $j: X \rightarrow \hat{X}$ with \hat{X} f -local.

Aim. Construct an " f -local" model structure in which:

- f -local = fibrant; fibrations are complicated BUT cofibrations of \mathcal{M}
- f -local equivalence = weak equivalence
- f -localization = fibrant replacement = "reflection"

\rightarrow better: $\text{Id}: \mathcal{M} \rightleftarrows \mathcal{M}^{f\text{-local}}: \text{Id}$

preserves cof. & w.e.

is interpreted as $\mathcal{M}_{\text{cf}} \xrightleftharpoons[\text{f-localization}]{} \mathcal{M}_{\text{cf}}^{f\text{-local}}$