

L6 Part 1 - Last categorical titbits

① The Triangle equations

Theorem

An adjunction $A \xrightleftharpoons[u]{F} B$

is specified by a pair of functors as above + natural transformations

$$\begin{array}{ccc} B & \xrightarrow{i} & B \\ F \downarrow \pi \Downarrow & \nearrow u & \\ A & & \end{array}$$

(unit)

$$\begin{array}{ccc} & u \nearrow & B \\ A & \xrightarrow{\varepsilon \Downarrow} & F \downarrow \\ & \searrow & \end{array} \quad \begin{array}{c} F \\ \Downarrow \end{array} \quad \begin{array}{c} \rightarrow \\ A \end{array}$$

(counit)

such that

$$\begin{array}{ccc} B & \xrightarrow{i} & B \\ F \downarrow \pi \Downarrow & \nearrow u & \\ A & \xrightarrow{\varepsilon \Downarrow} & A \end{array}$$

$$\begin{array}{ccc} & u \nearrow & B \\ A & \xrightarrow{\varepsilon \Downarrow} & F \downarrow \\ & \searrow & \end{array} \quad \begin{array}{c} i \\ \Downarrow \end{array} \quad \begin{array}{c} \rightarrow \\ A \end{array} = id_F$$

$$\begin{array}{ccc} F \pi_b & \rightarrow & FUFB \\ Fb \xrightarrow{\pi_b} & \Downarrow & \downarrow \varepsilon_{FB} \\ Fb & \xrightarrow{\text{?}} & FB \end{array} \quad \&$$

$$\begin{array}{ccc} \pi_{ua} & \rightarrow & UFua \\ ua \xrightarrow{\text{?}} & \Downarrow & \downarrow \varepsilon_a \\ ua & \xrightarrow{i} & ua \end{array}$$

~~Proof~~ (Sketch)

- Consider an adjunction $A(Fb, a) \xrightarrow{\epsilon} B(b, Ua)$.
 - We defined the unit $1 \Rightarrow UF$ last time.
 - Similarly $B(Ua, Ua) \xrightarrow{\epsilon^{-1}} A(FUa, a)$
- $\begin{array}{ccc} 1 & \xrightarrow{\epsilon^{-1}} & A(FUa, a) \\ & \longleftarrow & FUa \xrightarrow{\epsilon a} a \end{array}$
- determines the counit.

By naturality of ϵ^{-1} , have

$$\begin{array}{c} @ b \xrightarrow{\alpha} Ua \\ Fb \xrightarrow{\epsilon^{-1}\alpha} a \\ F\alpha \downarrow \quad \nearrow \epsilon a \\ FUa \end{array} + \begin{array}{c} @ \beta : Fa \rightarrow b \\ a \xrightarrow{\epsilon\beta} Ub \\ \nearrow " \quad \searrow UF \\ UFa \xrightarrow{\epsilon a} Uf \end{array} .$$

(last week)

- Now consider $B(Ua, Ua) \xrightarrow{\epsilon^{-1}} A(FUa, a) \xrightarrow{\epsilon} B(Ua, Ua)$
sends
 $1_{Ua} \mapsto \epsilon a : FUa \rightarrow a \mapsto Ua \xrightarrow{\pi_{Ua}} UFUa \xrightarrow{U\epsilon a} Ua$
so this equals the identity.
- Similarly the other triangle equation holds.

- Conversely consider ϵ & η satisfying triangle equations.
- We have $A(Fb, a) \xrightarrow{\epsilon} B(b, ua)$
- $Fb \xrightarrow{\alpha} a \xrightarrow{\eta_a} a \xrightarrow{ua} ua \xrightarrow{u\alpha} ub$
natural in a, b .
- So enough to show φ a bijection.
- Inverse given by

$$B(b, ua) \xrightarrow{\Theta} A(Fb, a)$$

$$b \xrightarrow{\alpha} ua \xrightarrow{\quad} Fb \xrightarrow{F\alpha} Fua \xrightarrow{\epsilon_a} a$$

& $\Theta\varphi = 1$, $\varphi\Theta = 1$ follow from triangle equations.

□

Remark) This defⁿ of adjunction via unit & counit + triangle equations captures sense in which adjunction is a weaker form of iso / equivalence.

Fully & Faithful Functors

A Functor $U: A \rightarrow B$ has functions $A(a, b) \xrightarrow{U_{a,b}} B(Ua, Ub)$.

Def) U is

- ① Faithful if each $U_{a,b}$ is injective
- ② Full - - - - - surjective
- ③ Fully Faithful - - - - - bijective.

- In elementary terms, Faithful means given $a \xrightarrow{\alpha} b$ st. $U\alpha = U\beta$ then $\alpha = \beta$.
- Full means given $Ua \xrightarrow{\alpha} Ub$ $\exists a \xrightarrow{\beta} b$ st $U\beta = \alpha$.

Defⁿ - A subcategory A of B is a subcollection of objects and arrows in B closed under identities and composition.

- Then $i: A \hookrightarrow B$ a Functor. If i is Fully Faithful it is a full subcategory.

Examples

- $U: \text{Grp}, \text{Mon}, \text{Ring} \longrightarrow \text{Set}$
are all faithful.
- $i: \text{Grp} \longrightarrow \text{Mon}$ is a
Full subcategory.

Lecture 6 Part 2 - Universal Algebra

- Universal : study of "sets with operations satisfy. equations"

E.g. groups, monoids, rings ...
operations : $x+y, 0, \dots$
equations : $x+0=x, \dots$

But not things like Fields :

$0 \neq 1$ not an equation

- "Universal algebra" also called "general algebra" ...
- Two excellent books I like :
 - Algebra Chapter 0 by Aluffi
 - An invitation to General Alg. by Bergman
(I am using this book a little)

Def") A signature $\Sigma = (\Sigma_n)_{n \in \mathbb{N}}$ consists of a set Σ_n for each $n \in \mathbb{N}$.

Remark) The elements of Σ_n are thought of as " n -ary operation symbol".
 0-ary operations = "constants".

Examples

① Signature Σ for monoids, groups... has
 $\Sigma = \{e, m\}$. i.e. $\Sigma(n) = \emptyset$ for $n \neq 0, 2$.

$e \in \Sigma(0)$ $m \in \Sigma(2)$

e nullary m binary

② Signature for rings

$\Sigma = \{0, 1, +, \times\}$

0-ary 2-ary

i.e.

- $\Sigma(0) = \{0, 1\}$
- $\Sigma(2) = \{+, \times\}$
- $\Sigma(n) = \emptyset$ otherwise.

Defⁿ let Σ be a signature.

An Σ -algebra is a set X together with, for each $n \in \mathbb{N}$ & $m \in \Sigma(n)$, a function

$$m_x : X^n \longrightarrow X : \\ (a_1, \dots, a_n) \mapsto m_x(a_1, \dots, a_n).$$

Remark what is X^0 ? well $X^n \cong \text{Set}(n, X)$, the set of functions from the n elt set $\rightarrow X$.

- In partic., $X^0 = \text{Set}(0, X) \cong \underline{\underline{I}}$ empty set which is initial.
- So given $e \in \Sigma(0)$ & X an Σ -algebra, we have $e_x : I \longrightarrow X$, so this is just an element e_x of X .
- This is why elts of $\Sigma(0)$ are thought of as nullary op. symbols = constants.

Notation) For an Σ -alg.

$$(X, (m_x)_{m \in \Sigma_n, n \in \mathbb{N}})$$

I will denote it by \overline{X} , or perhaps just X .

Example

For $\mathcal{S} = \{e, m\}$
 nullary binary

an \mathcal{S} -algebra is a magma:

- we have:

- a binary function $m_x : X^2 \rightarrow X : (a, b) \mapsto m_x(a, b)$
 $\text{ex: } i \rightarrow X : - \mapsto \text{lex}_X$

Magnas do not need to satisfy any equations.

Exercise) Find a signature \mathcal{S} st.
 IR-vector spaces are certain
 \mathcal{S} -algebras.

Def") A homomorphism $f: \bar{X} \rightarrow \bar{Y}$
 of \mathcal{S} -algebras is a function $f: X \rightarrow Y$
 such that $\forall n \in \mathbb{N}, m \in \mathcal{S}_n$

$$\begin{array}{ccc} x^n & \xrightarrow{f^n} & y^n \\ m_x \downarrow & " & \downarrow m_y \\ x & \xrightarrow{f} & y \end{array}$$

This says $\forall a_1, \dots, a_n \in X$, we have
 $f(m_x(a_1, \dots, a_n)) = m_y(fa_1, \dots, fa_n)$.

Remark

This captures the general notion of homomorphism in monoids, groups, rings etc when we choose the appropriate signature.

Example

IF $\Sigma = \{ \circ, \downarrow \}$
 $\begin{matrix} 2\text{-ary} & 0\text{-ary} \end{matrix}$

A homomorphism of Σ -algs is a fn

$f: X \rightarrow Y$ satisfying

$$f(a \cdot_x b) = f(a) \cdot_y f(b) \quad \& \quad f(c_x) = c_y .$$

(Here I write $a \cdot_x b = {}_x^*(a, b)$ for convenience.)

- There is a category $\Sigma\text{-Alg}$ of Σ -algebras & Σ -algebra homomorphisms,
 & it has a forgetful functor

$$U: \Sigma\text{-Alg} \longrightarrow \text{Set} .$$

- Our goal now is to describe its left adjoint.

Terms & term algebras

- Let Σ be a signature & X a set.

We define a set

$T_\Sigma(X)$ of " Σ -Terms" in X
as follows:

- $T_\Sigma^0(X) = X$
- $T_\Sigma^{i+1}(X) =$ formal expression
 $T_\Sigma^i(X) \cup \{m(t_1, \dots, t_n) : n \in \mathbb{N}, m \in \Sigma(n), \{t_1, \dots, t_n\} \subseteq T_\Sigma^i(X)\}$
- $T_\Sigma(X) = \bigcup_{i=0}^{\infty} T_\Sigma^i(X).$

The depth of a term $t \in T_\Sigma(X)$,
 $d(t)$ = $\min \{n \in \mathbb{N} : t \in T_\Sigma^n(X)\}$.

Example) $\mathcal{R} = \left\{ \begin{array}{l} m, e \\ \text{2-ary} \quad \text{0-ary} \end{array} \right\}$, $X = \{a, b, c\}$

$\text{Tr}(\mathcal{R})$

$= \{ a, b, c, \dots \text{ depth 0} \}$
 $e, m(a, b), m(a, a), m(b, b), m(b, a) \dots \text{depth 1}$
 $m(a, e), m(a, m(b, c)), \dots \text{depth 2} \dots$
 $m(m(a, m(b, e)), c) \text{etc} \dots \}$

- Remark: $\text{Tr}(\mathcal{R})$ consists of all expressions one can build from "variables" X and operation symbols in \mathcal{R} .

- In fact, $\text{Tr}(\mathcal{R})$ is an \mathcal{R} -algebra:

given $m \in \mathcal{R}_n$, we
define $\text{Tr}(\mathcal{R}) \xrightarrow{m} \text{Tr}(\mathcal{R})$:

$$(t_1, \dots, t_n) \xrightarrow{m} m(t_1, \dots, t_n)$$

- We also have a function

$$n_x: X \longrightarrow \text{Tr}(\mathcal{R}): x \mapsto x.$$

Theorem

$U: \mathcal{R}\text{-Alg} \rightarrow \text{Set}$
has a left adjoint,
whose value at X is the
 \mathcal{R} -algebra $\text{Tr}(X)$ of \mathcal{R} -terms.

~~Proof~~

Given $X \xrightarrow{f} UY$ we must
show that $\exists!$ \mathcal{R} -algebra map
 $\text{Tr}(X) \xrightarrow{\bar{F}} Y$ such that

$$\begin{array}{ccc} & \text{UTr}(X) & \\ \pi_X \nearrow & \parallel & \searrow u\bar{F} \\ X & \xrightarrow{F} & UY \end{array} .$$

- Since the triangle must commute, we must define $\bar{F}x = fx$ for $x \in X$.
- We define \bar{f} inductively on depth:
 - depth 0 terms are elements of X , where $\bar{f}x = fx$ as above.
 - Suppose we have defined it on terms of depth $\leq i$. Then consider a term $m(t_1, \dots, t_n)$ of depth $i+1$.

- Then t_1, \dots, t_n are of depth $\leq i$.

In order that \bar{f} is a homomorphism, we must set

$$\textcircled{3} \quad \bar{f}(m(t_1, \dots, t_n)) = m'(\bar{f}t_1, \dots, \bar{f}t_n)$$

already defined

- This completes the definition of \bar{f} &

\textcircled{4} shows that it is a homomorphism.

- Since we were forced to define \bar{f}

in this way, it is the unique extension
of f to a homomorphism as
required. \square

Example

If $\mathcal{R} = \{m, e\}$,
 $X = \{a, b, c\}$ &

$X \xrightarrow{f} \cup Y$, we obtain

$T_{\mathcal{R}} X \xrightarrow{\bar{f}} Y$

$a, b, c \xrightarrow{} f_a, f_b, f_c$

$e \xrightarrow{} e_Y$

$m(a, b) \xrightarrow{} m_Y(f_a, f_b)$

$m(a, m(b, c)) \xrightarrow{} m_Y(f_a, m_Y(f_b, f_c))$

.....

Equations

- Let \mathcal{S} be a signature.
- Then given a set X , we have the set $\text{Tr}(X)$ of \mathcal{S} -terms in X .

Def") An \mathcal{S} -equation in variables X is a pair $(s, t) \in \text{Tr}(X)^2$.

Remark) We often informally write an \mathcal{S} -equation as " $s = t$ ".

Example

For $\mathcal{S} = (\text{e}, \cdot)$ multary bin. The

signature for magmas, the following are equations

- $x \cdot e = x$
 - $e \cdot x = x$
 - $(x \cdot y) \cdot z = x \cdot (y \cdot z)$
- in variables $\{x, y, z\}$.

Satisfying equations

- Consider an X -tuple $a = (a_x)_{x \in X}$ of an \mathcal{R} -algebra A .

- In other words
a function

$$a: X \rightarrow UA$$
$$x \mapsto a_x$$

- We have $UTr(X) \xrightarrow{\bar{a}} U\bar{a}$

$$\begin{matrix} \pi_X \uparrow & \searrow \\ X & \xrightarrow{a} UA \end{matrix}$$

Defⁿ let $(s, t) \in Tr(X)$ be an equation & A an \mathcal{R} -algebra
Then the X -tuple a satisfies the equation $s = t$ if $\bar{a}(s) = \bar{a}(t)$.

- The \mathcal{R} -algebra A satisfies $s = t$ if each X -tuple a of A does so.

One writes $A \models s = t$.

Example

Given $\{x, y, z\} \xrightarrow{a} A$ a magma,
 $x, y, z \longmapsto a_1, a_2, a_3$
the extension

$$\begin{aligned} T_{\mathcal{R}}\{x, y, z\} &\longrightarrow A \\ (x \cdot y) \cdot z &\longmapsto (a_1 \cdot a_2) \cdot a_3 \\ x \cdot (y \cdot z) &\longmapsto a_1 \cdot (a_2 \cdot a_3) \end{aligned}$$

so A satisfies the equation

$$(x \cdot y) \cdot z = x \cdot (y \cdot z) \text{ iff}$$

A is associative.

Def) A universal algebraic Theory (\mathcal{R}, E) is a signature \mathcal{R} and a set E of \mathcal{R} -equations.

- By an (\mathcal{R}, E) -algebra we mean an \mathcal{R} -algebra A such that $A \models s = t$ for each equation $(s, t) \in E$.

Examples

• For $\mathcal{R} = \{\cdot, e\}$ as before, and

$$E = \left\{ \begin{array}{l} (x, y).z = x.(y.z) \\ x.e = x \\ e.x = x \end{array} \right\} \quad \begin{array}{l} \text{an } (\mathcal{R}, E)\text{-algebra} \\ \text{is a } \underline{\text{monoid}}. \end{array}$$

• If we take $\mathcal{R} = \{\cdot, e, (-)^{-1}\}$ and add the equations

$$\begin{array}{l} x.x^{-1} = e \\ x^{-1}.x = e \end{array} \quad \text{to } E,$$

then an (\mathcal{R}, E) -algebra is a group.

Defⁿ) For (\mathcal{R}, E) as above, we write

$(\mathcal{R}, E)\text{-Alg} \longleftrightarrow \mathcal{R}\text{-Alg}$
for the full subcategory of
 (\mathcal{R}, E) -algebras.

This means that the objects

are (\mathcal{R}, E) -algebras and

the morphisms simply the

\mathcal{R} -algebra homomorphisms
between them.

- The inclusion $(\mathcal{R}, E)\text{-Alg} \hookrightarrow \mathcal{R}\text{-Alg}$ is a (Fully Faithful) Functor.
- We obtain a composite forgetful functor to Set, as in the diagram below:

$$\begin{array}{ccc} (\mathcal{R}, E)\text{-Alg} & \xrightarrow{i} & \mathcal{R}\text{-Alg} \\ U \searrow & & \swarrow U_{\mathcal{R}} \\ & \text{Set} & \end{array}$$

- Later we will show that i and U have left adjoints - in particular, there exist free (\mathcal{R}, E) -algebras.

Examples

- All of the algebraic categories earlier considered -

e.g. Vect, Grp, Rng, Mon etc -
are of the form $(\mathcal{R}, E)\text{-Alg}$
for suitable \mathcal{R} and E .

- This is the framework and scope of what people typically call universal algebra.

- Our goal now is to study some of the good properties of categories of the form $\mathcal{R}\text{-Alg}$ and, more generally, $(\mathcal{R}, E)\text{-Alg}$.
- In fact, $(\mathcal{R}, E)\text{-Alg}$ is closed in $\mathcal{R}\text{-Alg}$ under all of the constructions we will care about - limits, coequalisers, so we can concentrate on $\mathcal{R}\text{-Alg}$ firstly.

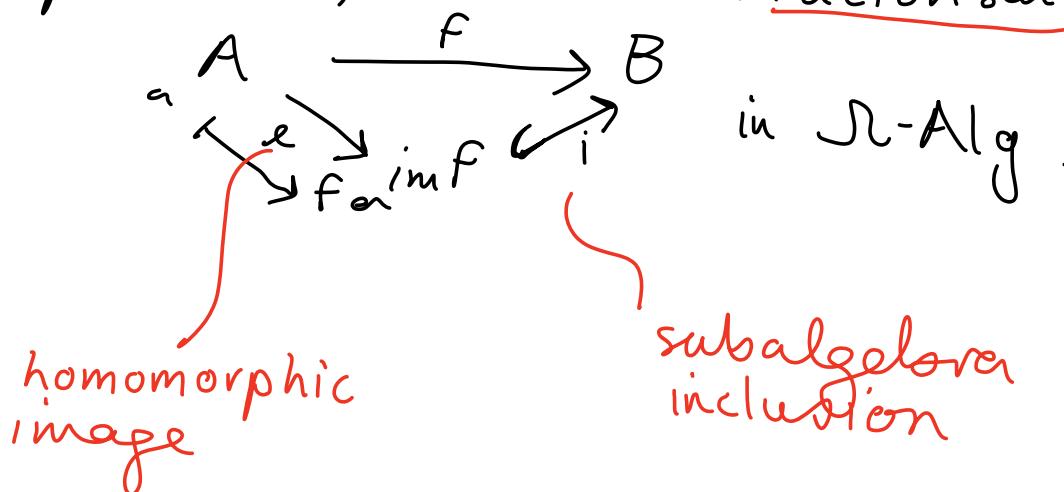
First, a definition :

Defⁿ) Let $A \in \mathcal{R}\text{-Alg}$. A subalgebra
 $B \hookrightarrow A$ is a subset B of A
such that if $s \in \mathcal{R}$ & $b_1, \dots, b_n \in B$
then $s(b_1, \dots, b_n) \in B$.

• In particular a subalgebra B is itself
an \mathcal{R} -algebra & the inclusion
 $B \hookrightarrow A$ a homomorphism.

Defⁿ) A homomorphic image is a
homomorphism $f: A \rightarrow B$ which
is surjective.

- Let $f: A \rightarrow B$ be a homomorphism.
 Let $\text{im } f = \{b \in B : \exists a \text{ with } fa = b\}$.
 Then $\text{im } f \hookrightarrow B$ is a subalgebra of B ,
 since if $b_1 = fa_1, \dots, b_n = fa_n$ & $s \in \mathcal{S}_n$,
 then $s(b_1, \dots, b_n) = s(fa_1, \dots, fa_n)$
 $= f(s(a_1, \dots, a_n)) \in \text{im } f$.
- In particular, we obtain a Factorisation



Proposition

$\mathcal{R}\text{-Alg}$ has all limits & $U: \mathcal{R}\text{-Alg} \rightarrow \text{Set}$ preserves them.

Proof I will just check that $\mathcal{R}\text{-Alg}$ has (possibly infinite) products.

and equalisers, since these generally are limits (not proven in course) and are important for us.

- Consider a family $(A_i)_{i \in I}$ of \mathcal{R} -algebras.

- Their product as sets is the direct product

$$\prod_{i \in I} A_i = \{ \bar{a} = (a_i)_{i \in I} : a_i \in A_i \} \xrightarrow{\pi_i} A_i \\ (a_i)_{i \in I} \longmapsto a_i .$$

- We want to show $\prod_{i \in I} A_i$ has the structure of an \mathcal{R} -algebra such that each π_i is a homomorphism: this says: given $s \in \mathcal{R}^n$ and $\bar{a}', \dots, \bar{a}^n$ we

$$\text{have } s(\bar{a}', \dots, \bar{a}^n)_i = s(\bar{a}'_i, \dots, \bar{a}^n_i) \sim$$

i.e. the componentwise \mathcal{R} -algebra structure.

- Given an \mathcal{R} -algebra B and $(f_i : B \rightarrow A_i)_{i \in I}$ we have a unique function $B \xrightarrow{f} \prod_{i \in I} A_i$ such that $\pi_i \circ f = f_i$ for each $i \in I$: namely $(f(b))_i = f_i(b)$, and we must check that f is a homomorphism if each f_i is one: i.e. $f(s(b'), \dots, b^n) = s(f(b'), \dots, f(b^n))$

but the components at $i \in I$ are

$$f_i(s(b'), \dots, b^n) = s(f_i(b'), \dots, f_i(b^n))$$

which are equal since f_i a homomorphism.

- Given $A \xrightarrow{\begin{matrix} f \\ g \end{matrix}} B$ the equaliser

$$E = \{ x \in A : f(x) = g(x) \} \xrightarrow{i} A \text{ in Set}$$

is a subalgebra of A :

i.e. given x_1, \dots, x_n st $Fx_i = gx_i$ then

$$f(s(x_1, \dots, x_n)) = s(fx_1, \dots, fx_n) = s(gx_1, \dots, gx_n) = g(s(x_1, \dots, x_n))$$

so $s(x_1, \dots, x_n) \in E$

In particular, $i: E \rightarrow A$ is a homomorphism and it is easy to check it has the u.p. of the equaliser.

□

What about colimits? The key kind are the simple quotients by congruences.

- These generalise :

- quotients by equiv. rels (for sets)
- quotients by normal subgroups (for groups)
ideals (for rings)

these last two are less obvious.

Def") - let A be an \mathcal{R} -algebra. An equivalence relation $E \subseteq A$ is called a congruence if E is a subalgebra of A .

- In elementary terms, a cong. is an \mathcal{E} -rel E : (so $xEx, xEy \Rightarrow yEx, xEy \wedge yEqz \Rightarrow xEqz$) such that if $x, Eqy_1, \dots, x_n Eqy_n$ & $s \in \mathcal{R}_n$ then $s(x_1, \dots, x_n) Eq s(y_1, \dots, y_n)$.

Exercise : • For groups, show that there

is a bijⁿ bet. cong. on A & normal subgroups of A =

$$\begin{array}{ccc} E & \xrightarrow{\quad} & \bar{E} = \{x \in A \text{ s.t. } \\ & \downarrow & \downarrow x \in N\} \\ A \times A & \xrightarrow{\quad} & A \end{array} \quad \left| \quad \begin{array}{ccc} N & \xrightarrow{\quad} & E_N = \{(x, y) \\ & \downarrow & \downarrow xy^{-1} \in N\} \\ A & \xrightarrow{\quad} & A \times A \end{array} \right.$$

- What about rings?

- If $E \hookrightarrow A^2$ is a congruence on A , then we can form the diagram

$$E \xrightarrow{i} A^2 \xrightarrow{\begin{matrix} p_1 \\ p_2 \end{matrix}} A \text{ in } R\text{-Alg},$$

and so obtain a pair of R -alg. homomorphisms

$$E \xrightarrow{\begin{matrix} a \\ c \end{matrix}} A : (x, y) \in E \xrightarrow{\begin{matrix} x \\ y \end{matrix}} y .$$

- We are interested in their coequalisers

$$A \xrightarrow{p} A/E$$

defined as follows :

- elements of A/E are equiv. classes $[a]$, with $p(a) = [a] = \{x : x \in a\}$
- observe that p is surjective. Therefore if p is to be a homomorphism, we are forced to define $s^{[a]}([a_1], \dots, [a_n]) = [s(a_1, \dots, a_n)]$.
- Is this well defined?

Suppose $[b_1] = [a_1], \dots, [b_n] = [a_n]$, so that $b_1 \in a_1, \dots, b_n \in a_n$. Then $s(b_1, \dots, b_n) \in s(a_1, \dots, a_n)$ as we have a congruence, so that

$[s(b_1, \dots, b_n)] = [s(a_1, \dots, a_n)]$ as required.

In particular, A/E is an \mathbb{R} -algebra and $p: A \rightarrow A/E$, a surjective homomorphism.

Proposition

$$E \xrightarrow{d} A \xrightarrow{p} A/E$$

is a wequaliser in $\mathbb{R}\text{-Alg}$.

Proof - Firstly, if $(x, y) \in E$, then

$$pd(x, y) : [x] = [y] = pc(x, y) \text{ so}$$

$$pd = pc.$$

- Given $A \xrightarrow{f} B$ with $f|_E = f|_C$,
this means if $(x, y) \in E$ then $fx = fy$,
so that $[x] = [y] \Rightarrow fx = fy$.

- Define $A/E \xrightarrow{\bar{f}} B$ $[a] \mapsto f(a)$ for some choice

- The def. of \bar{f} representatives.
of \bar{f} is indep. of the choice as f is.

- The triangle

$$\begin{array}{ccc} A & \xrightarrow{p} & A/E \\ & \downarrow & \downarrow \bar{f} \\ & & B \end{array}$$

commutes for this reason too.

- Clearly \bar{f} is a homomorphism, since f is.

Def') Let $f: A \rightarrow B$ be an \mathbb{R} -alg homomorphism. The kernel of f is the congruence

$$K_f = \{(x, y) : fx = fy\} \hookrightarrow A^2.$$

- It is easy to see that this is a congruence. Check it!

Categorically, K_f is the pullback

$$\begin{array}{ccc} (x,y) & \xrightarrow{\quad} & x \\ \downarrow c \quad \lrcorner \quad \lrcorner \quad \downarrow f & \xrightarrow{d} & A \\ y & \xrightarrow{\quad} & A \xrightarrow{f} B \end{array}$$

so, in particular,
we have
 $f d = f c$.

Therefore we get a unique
factorisation of f through the

coequaliser $A \xrightarrow{f} B$

$$\begin{array}{ccc} a & \swarrow p & \nearrow \pi_f \\ A/K_f & & [a] \end{array}$$

f_a

Fist isomorphism theorem

- Given a homomorphism $f: A \rightarrow B$, we have an isomorphism $t: A/K_f \rightarrow \text{im } f$ making the diagram

$$\begin{array}{ccccc}
 & & A/K_F & & \\
 \rho \nearrow & & \downarrow s/t & \searrow \bar{f} & \\
 A & = & \text{im } f & = & B \\
 f_1 \searrow & & \swarrow f_2 & & \\
 & & & &
 \end{array}$$

Proof - The formula is
 $[a] \xrightarrow{t} fa$

It is clearly a homomorph as \bar{f} is.

- Let us show t is injective.
 Suppose $t[a] = t[b]$.
- That is, $fa = fb$.
- Then $(a, b) \in K_F$ so $[a] = [b]$.
- For surjectivity, if $b \in B$, we have $b = fd$ so
 $t[a] = fa = b$. □

Corollary

If $f: A \rightarrow B$ is a surjective homom.,

then $A/K_F \cong B$. In particular

$K_F \xrightarrow[c]{d} A \xrightarrow{f} B$ is a coequaliser diagram.

Proof / In this case, $B = \inf$.